

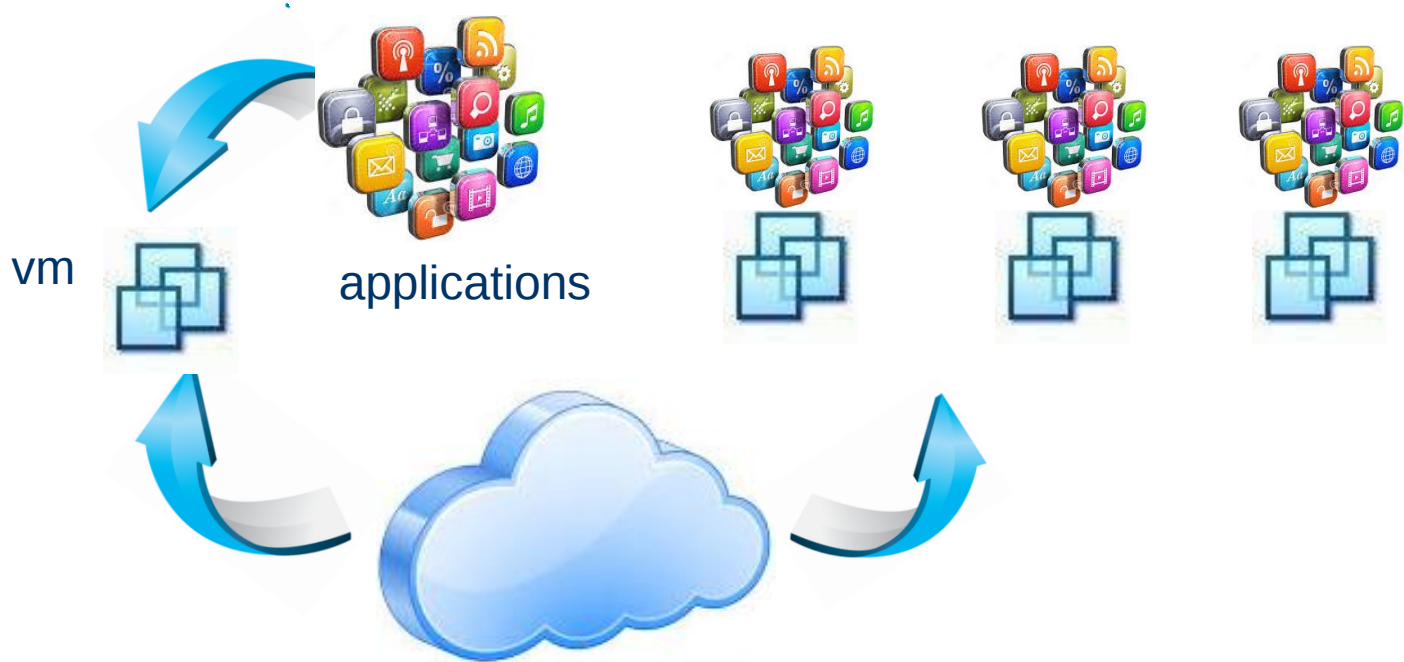
# Modelling Adaptation Policies as Domain-Specific Constraints

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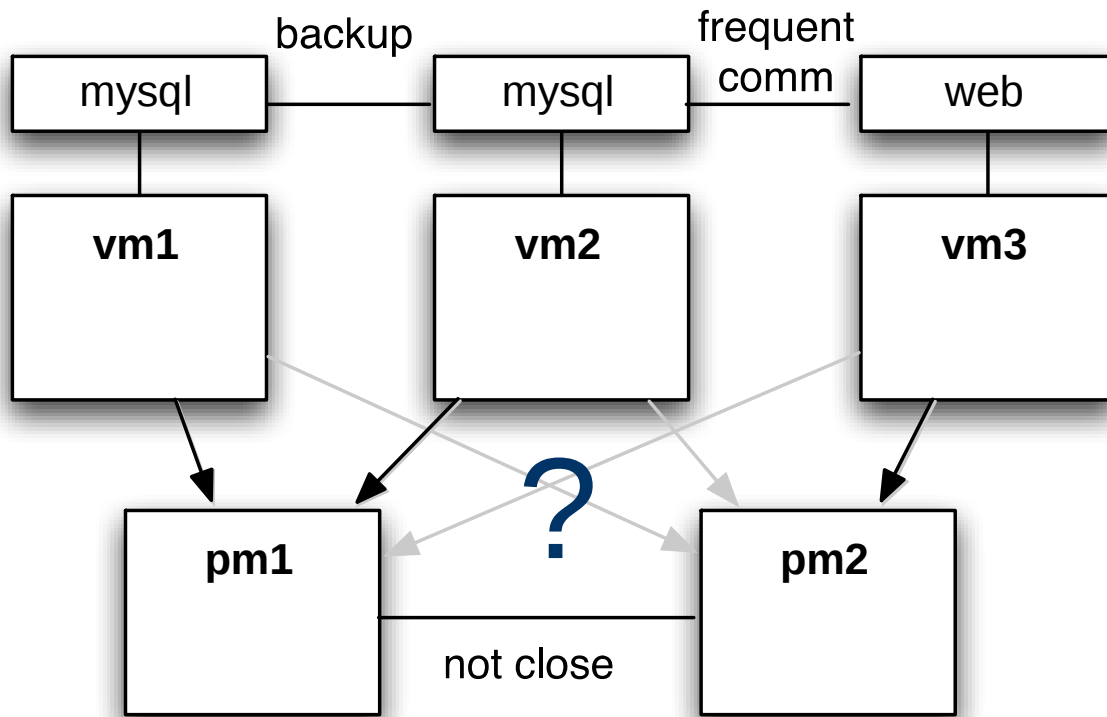
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# VM Placement in Cloud



# Thinking in Models



Resource limitation

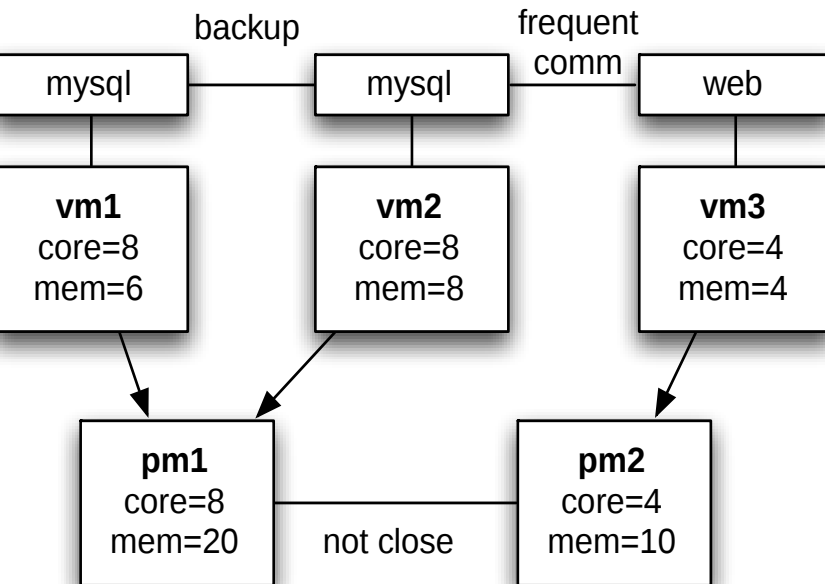
Consolidation

Backup split

Frequent close

Migration cost

# Developing the Adaptation Behaviour



- Concentrate VMs to fewer pms to save energy
  - but only when the sum of VMs' memory does not exceed the pms' memory
  - Don't move very big VMs
- Separate backup VMs to different PMs
  - But only when the sum of VMs' memory...
  - Don't move very big...
  - Concentrate VMs..., when possible
- Put frequently communicating VMs closer
  - Separate backup VMs...
    - But only when the sum...
    - Don't move very big...
  - Concentrate...
  - ...

```

if backup(vm1,vm2) and vm1.host = vm2.host,
  if vm1.mem < vm2.mem or ( frqt(vm1, vm2) and not vm1.mem >> vm2.mem )
    if vm2.mem < pm2.available and vm1.core <= pm2.core
      move vm1 to pm2
    else
      ...
  
```

If written in Action-based adaptation policy

# Modelling Adaptation Policies

## ■ Challenges

- Many interrelating concerns

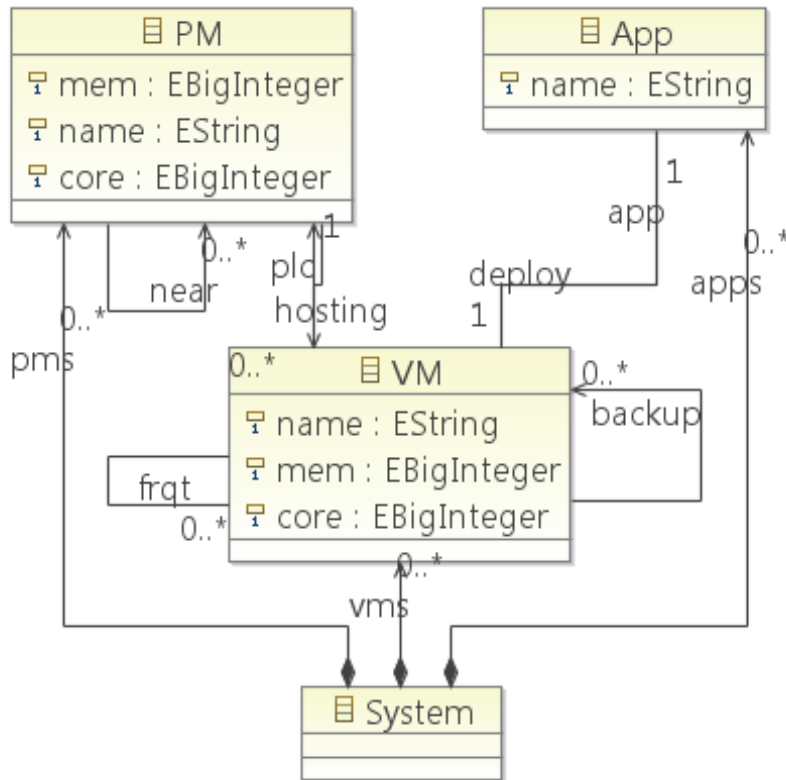
## ■ Actions policies (if-then-else or event-condition-action)?

- Explosion of branches
- Hard to introduce new concerns
- Abstraction gap between concerns and actions

## ■ New way of modeling!

- Just write down the constraints themselves
  - "what the system should be like" rather than "how to achieve that"
- Potential conflicts?
  - Soft constraints with different weights

# Modeling Adaptation Policies as Constraints



## Memory Limitation

### context PM inv:

hosting->collect(mem)->sum() <= mem  
(priority: mandatory)

## Consolidation

context PM inv: self.hosting->size() = 0  
(priority: low)

## Backup split:

### context VM inv:

backup->forall(e|e.plc != self.plc)  
(priority: high)

## Migration cost:

vm1.plc = pm1 (priority: 8)

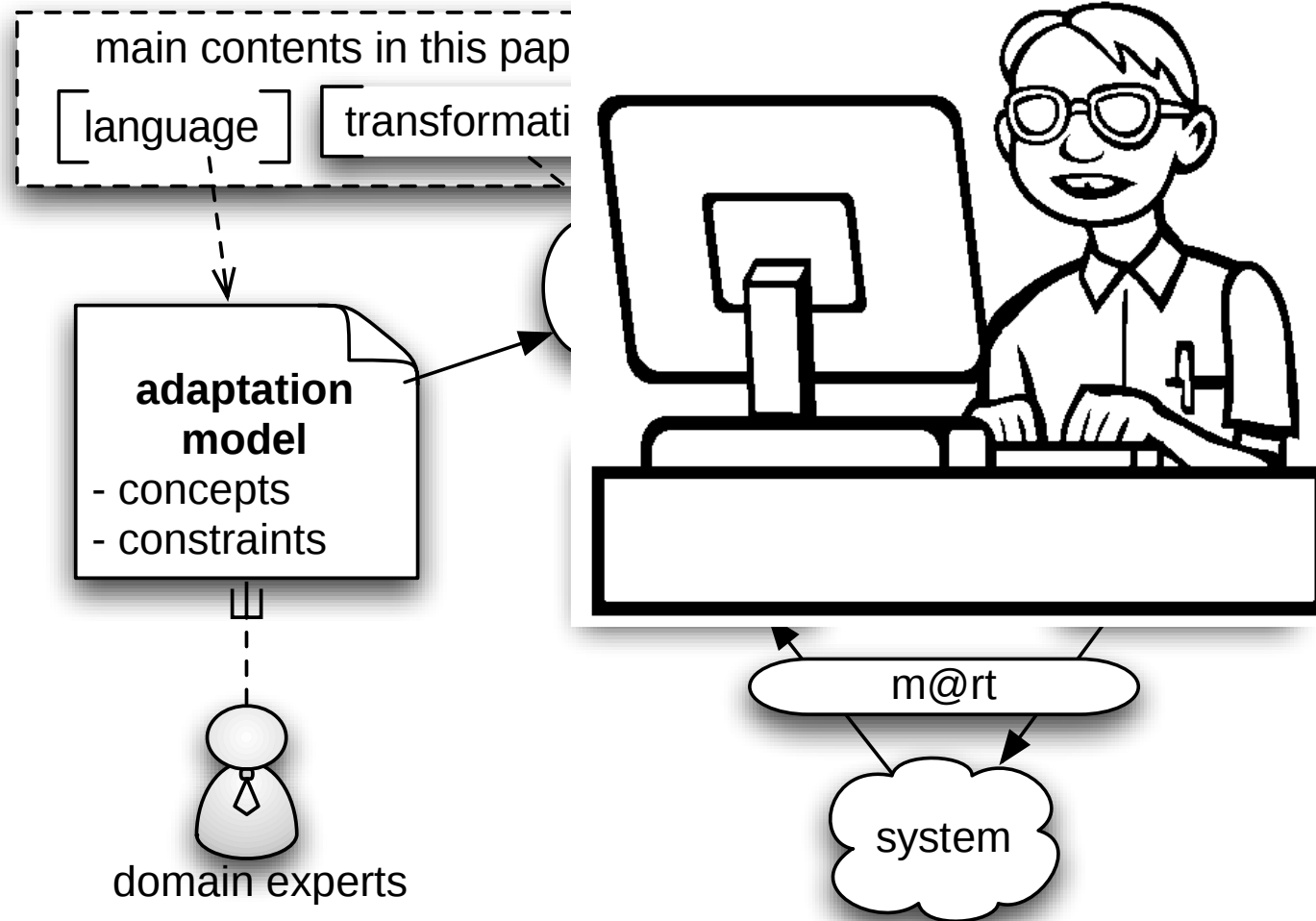
Vm3.plc = pm2 (priority: 4)

# Constraint Modelling Language and Editor

```
9  class VM{
10      id String name      Integer mem
11      config Integer core : {domain = Set{1,2,4,8} resistance = 20}
12      config PM plc : { resistance=(mem*10) }
13      VM[*] backup : {derived
14          VM.allInstances()->select(v|v<>self and v.app.name = self.app.name) }
15      ref VM[*] frqt opposite frqt
16      ref App app opposite deploy
17      constraint (hard) CoreLimit: self.core <= self.plc.core
18      constraint (priority = 80) BackupSplit: backup->forAll(ele.plc <> self.plc)
19      constraint (priority = 30) FrequentNear :
20          frqt->forAll(v|v.plc.near->includes(self.plc)) }
21  class PM{
22      id String name
23      Integer mem      Integer core      ref PM[*] near
24      ref VM[*] hosting:{derived VM.allInstances()->select(plc=self)}
25      constraint(hard) MemLimit: hosting->collect(ele.mem)->sum() < mem }
26  class App{
27      String name
```

<https://bitbucket.org/huis/constraintml>

# Adaptation Based on Constraints

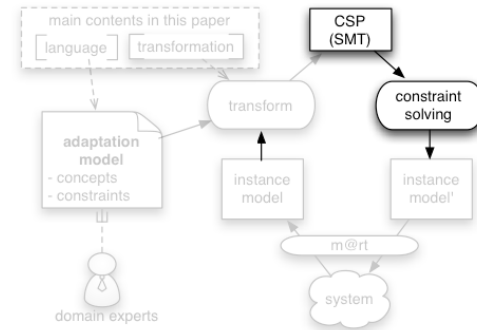
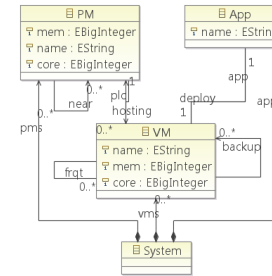


# Satisfactory Modulo Theory

## Memory Limitation

context PM inv:

hosting- $\rightarrow$ collect(mem)- $\rightarrow$ sum()  $\leq$  mem



$VM : \{vm_1, vm_2, vm_3\}, PM : \{pm_1, pm_2\}, int, boolean$

$vc : VM \rightarrow int, pc : PM \rightarrow int, vmem : VM \rightarrow int, pmem : PM \rightarrow int$

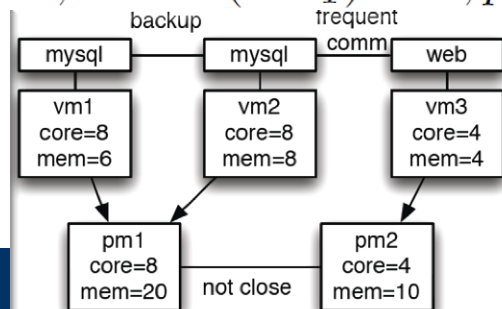
$plc : VM \rightarrow PM, frqt : VM \times VM \rightarrow bool, near : PM \times PM \rightarrow bool$

$\forall vm, pm. (plc(vm) = pm \Rightarrow vc(vm) \leq pc(pm))$

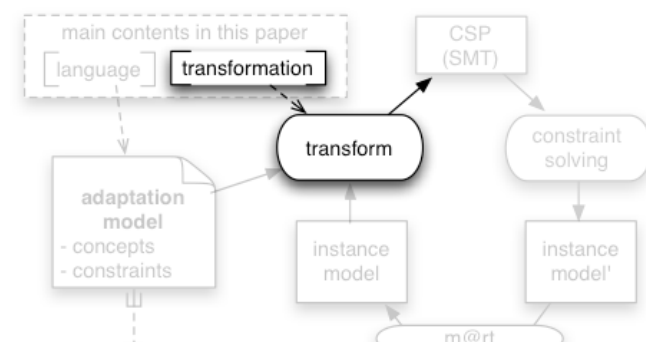
$\forall pm. \left( \sum_{vm \in VM} ite(plc(vm) = pm, vmem(vm), 0) \leq pmem(pm) \right)$

$\forall vm, vm'. (frqt(vm, vm') \Rightarrow near(plc(vm), plc(vm'))), \forall pm. (near(pm, pm))$

$vc(vm_1) = 8, vmem(vm_1) = 6, plc(vm_2) = pm_1, frqt(vm_2, vm_3)...$



# Transformation



## ■ OCL expressions

■ Principle #1: values, elements and most operations keep the same

■ Principle #2: replace property reference by function call

■  $vm1.plc = pm1 \Rightarrow plc(vm_1) = pm_1$

■  $vm1.frqt = \{frqt(vm_1, vm_1)? vm_1: \perp, frqt(vm_2, vm_1)? vm_2: \perp \dots\}$

■ Principle #3: properties that do not change will be resolved

■  $vm1.backup \rightarrow forall(e|e.plc != vm1.plc) \Rightarrow plc(vm_1) != plc(vm_2)$

■ Principle #4: complex operations have simple alternatives

■  $vm1.frqt.forall(e.x) \Rightarrow$

$(frqt(vm_1, vm_1)? x(vm_1): True) \& (frqt(vm_2, vm_1)? x(vm_2): True \& \dots$

## ■ A new set of partial-evaluation semantic rules on OCL

■ *"which this slide is too narrow to contain"*

# Constraint Solving

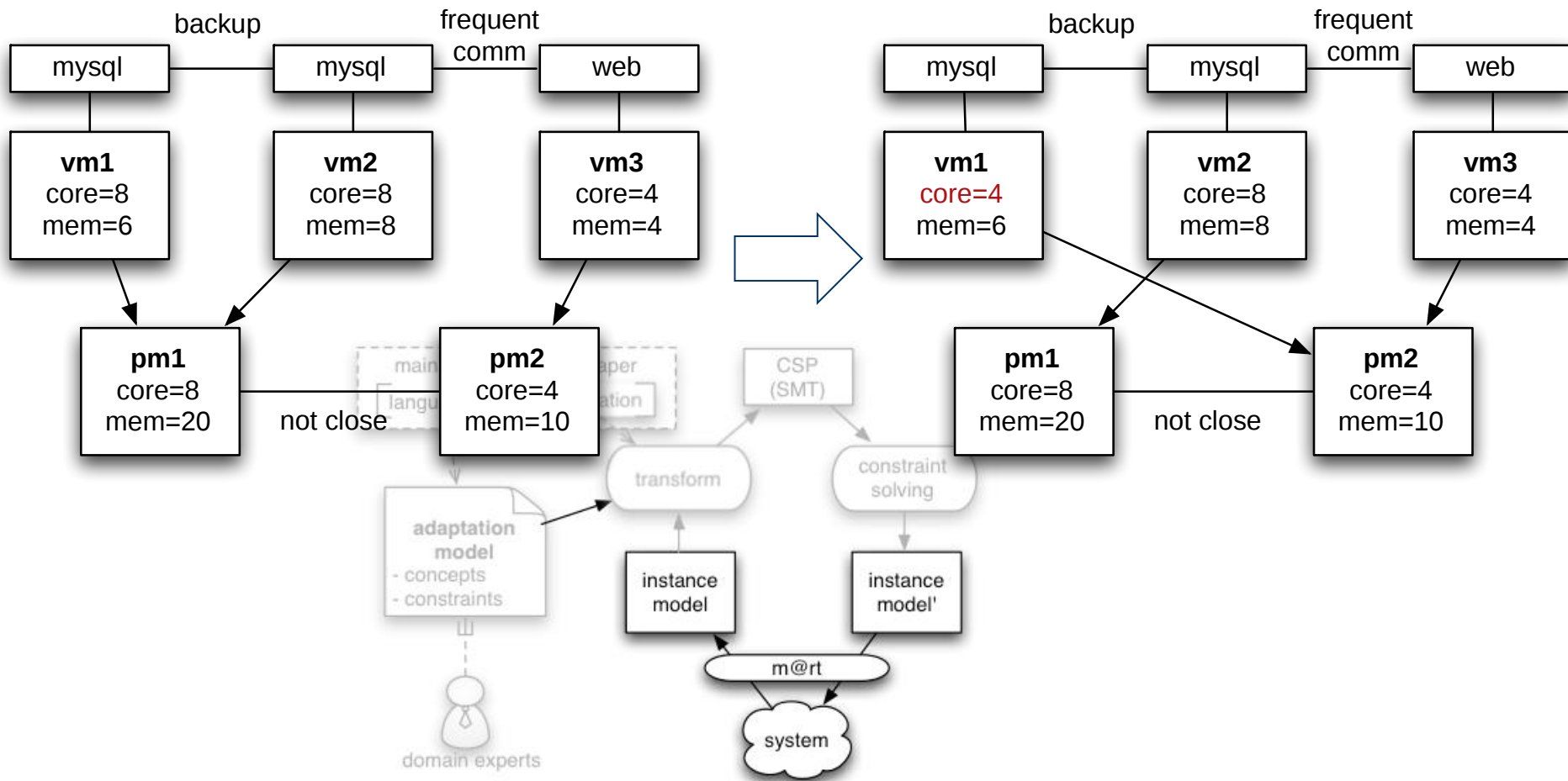
$$\begin{array}{c}
 \text{VM} : \{vm_1, vm_2, vm_3\}, \text{PM} : \{pm_1, pm_2\}, \text{int}, \text{boolean} \\
 vc : \text{VM} \rightarrow \text{int}, pc : \text{PM} \rightarrow \text{int}, vmem : \text{VM} \rightarrow \text{int}, pmem : \text{PM} \rightarrow \text{int} \\
 plc : \text{VM} \rightarrow \text{PM}, frqt : \text{VM} \times \text{VM} \rightarrow \text{bool}, near : \text{PM} \times \text{PM} \rightarrow \text{bool} \\
 \hline
 \forall vm, pm. (plc(vm) = pm \Rightarrow vc(vm) \leq pc(pm)) \\
 \\
 \forall pm. \left( \sum_{vm \in \text{VM}} \text{ite}(plc(vm) = pm, vmem(vm), 0) \right) \leq pmem(pm) \\
 \\
 \forall vm, vm'. (frqt(vm, vm') \Rightarrow near(plc(vm), plc(vm'))), \forall pm. (near(pm, pm)) \\
 \hline
 vc(vm_1) = 8, vmem(vm_1) = 6, plc(vm_2) = pm_1, frqt(vm_2, vm_3) \dots
 \end{array}$$

- Constraint Solver: Is there an interpretation to each function, to satisfy all the constraints?
  - Yes: Return the interpretation,
  - No: Ignore some "weakest" constraints, and return an interpretation to satisfy all the others – An optimisation problem
- Solver: Z3 by Microsoft research



Song, H., S. Barrett, A. Clarke, and S. Clarke (2013). Self-adaptation with End-User Preferences: Using Run-Time Models and Constraint Solving. In: Model-Driven Engineering Languages and Systems. pp.555–571.

# models@runtime



Nicolas Ferry, Hui Song, Alessandro Rossini, Franck Chauvel and Arnor Solberg, CloudMF: Applying MDE to Tame the Complexity of Managing Multi-Cloud Applications, UCC 2014, to appear

# A Demo

- Focus on the adaptation effect
- Ignore SMT generation and models@runtime

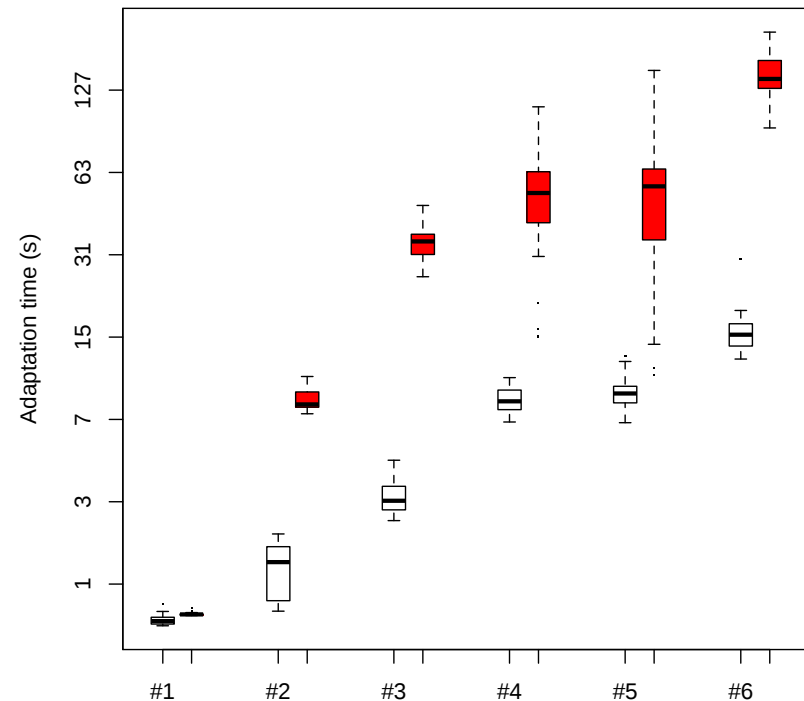
# Performance

## ■ Acceptable for medium sized private clouds

- 60s: 100 vm, 10 pm, 600 properties, 10 changes (as in paper)
- 60s: 500 vm, 50pm, 3000 properties, 10 changes (now)

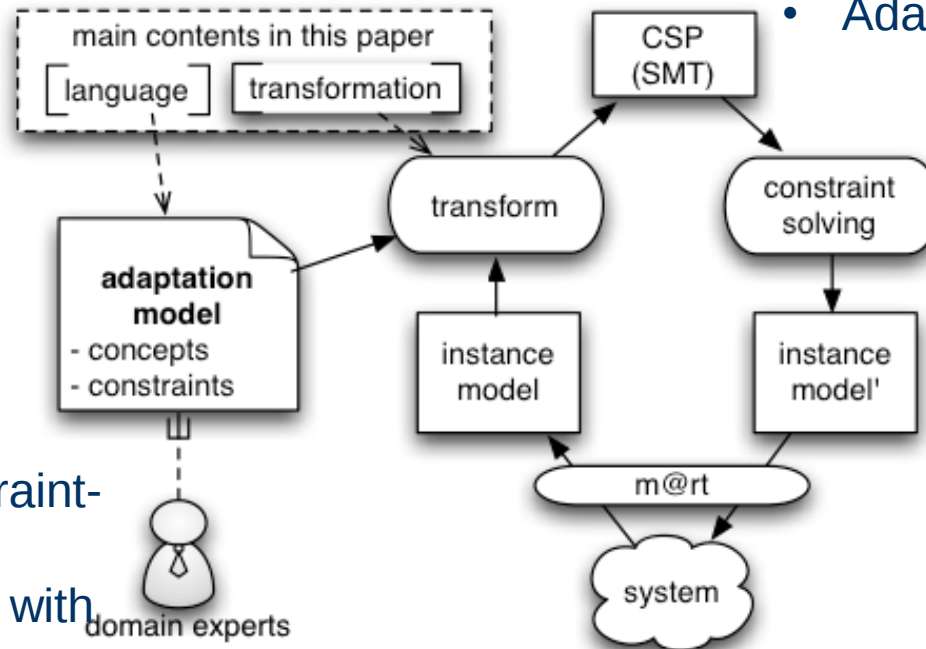
## ■ Why

- Powerful new constraint solver
- Usually simple constraints
- A big portion of fixed properties
- Partial evaluation!



# A Short Summary

- SMT representation of architectural models
- OCL to SMT transformation, with **partial evaluation**
- SMT solving with soft constraints
  - Conflicting constraints
  - Adaptation costs



- Declarative constraint-based modelling
- A text-based DSL with a powerful editor

# Application

- Directly used for adaptation
  - Constraints from cloud domain experts
  - Searching for better deployment that fits the context better
- Assessment of adaptation cost
  - *"Is a diverse system easier to be adapted for changing contexts"?*
  - "Dry-run" the solving process on controlled models
  - The total weight of broken constraints is the cost of performing the adaptation

# Thank You!

## Questions, Comments, Suggestions?

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<https://github.com/songhui/cspadapt>



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