

Participatory Networking

Rodrigo Fonseca

Join work with Andrew Ferguson, Arjun Guha, Jordan Place



BROWN

Networking in the Cloud



Networking in the Cloud

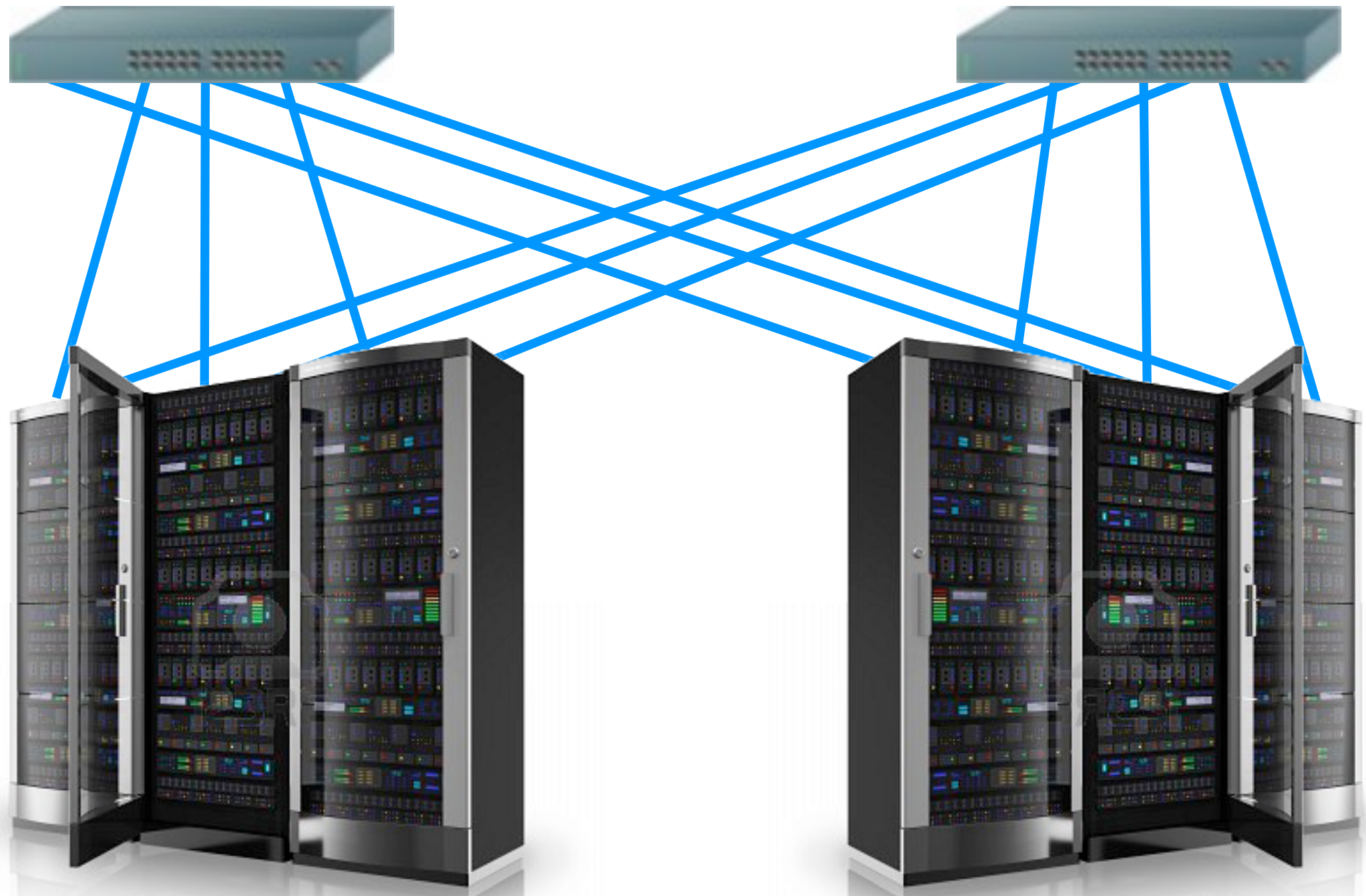


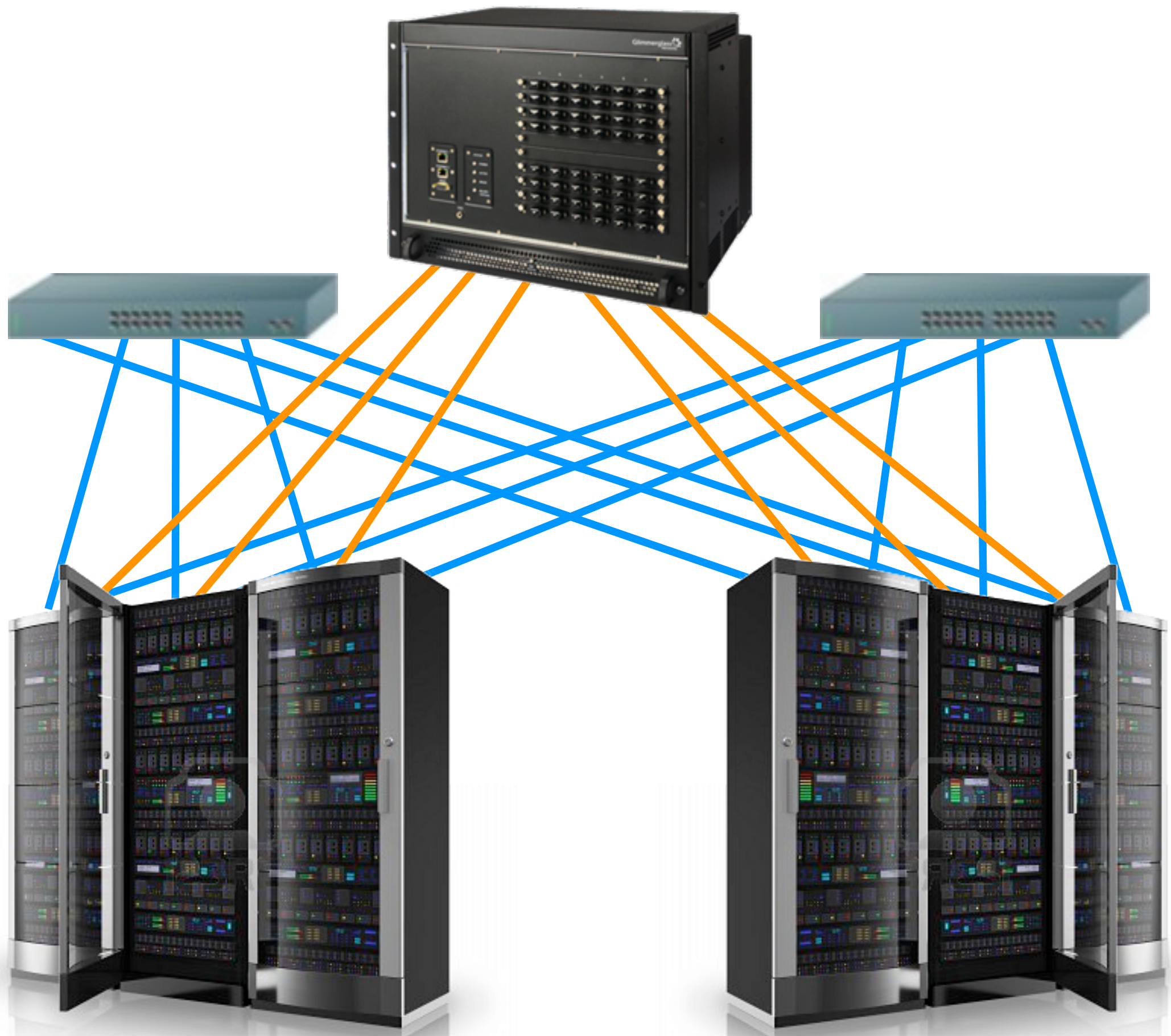
Can the cloud provider get **help** in configuring the network?

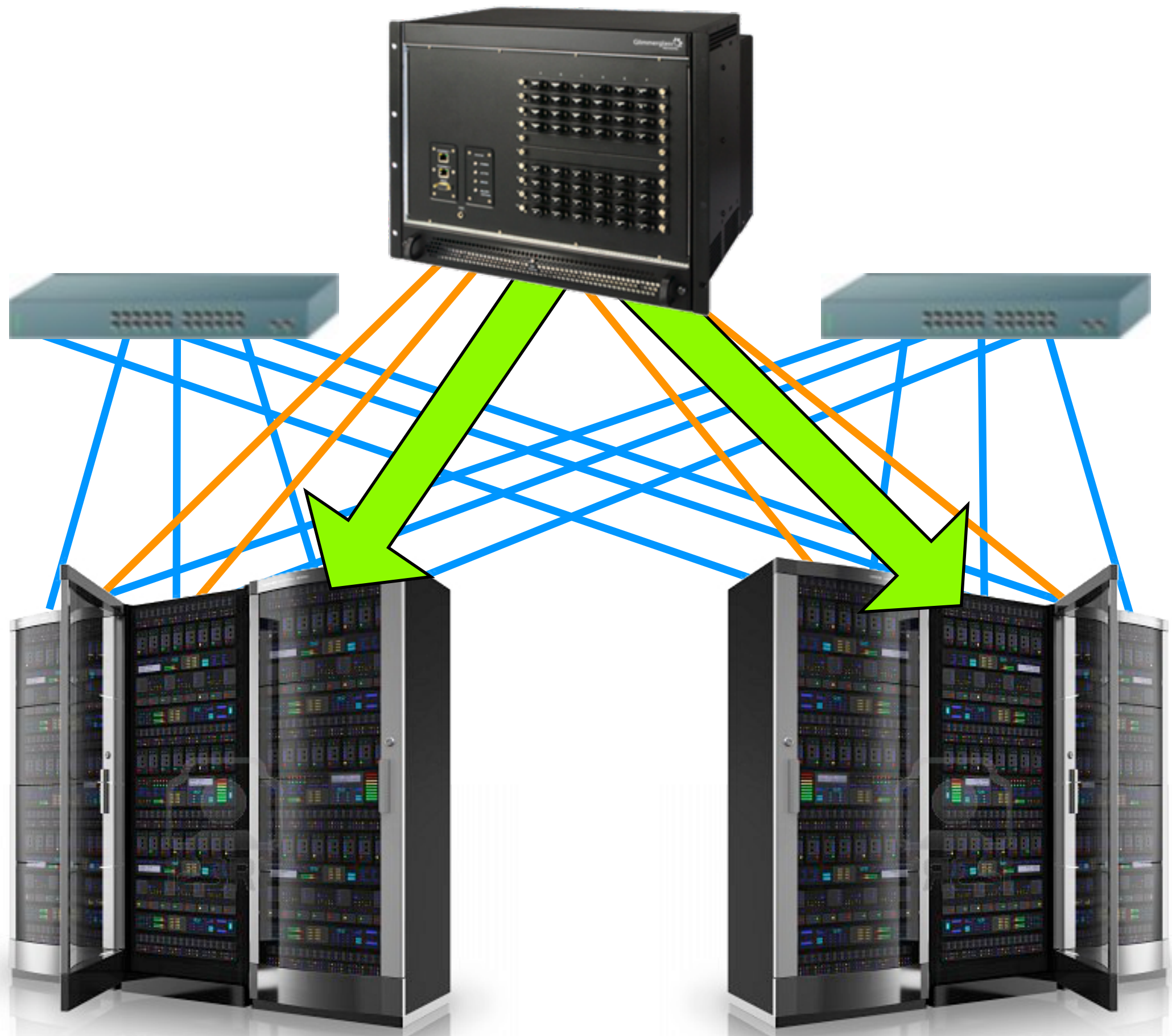
A few motivating examples

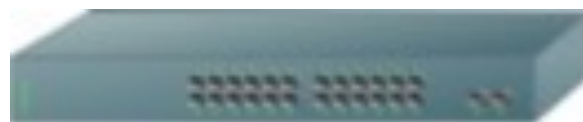
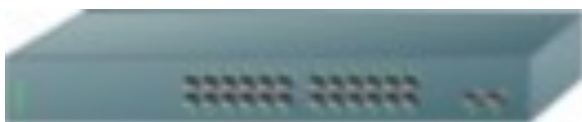
Large-scale processing

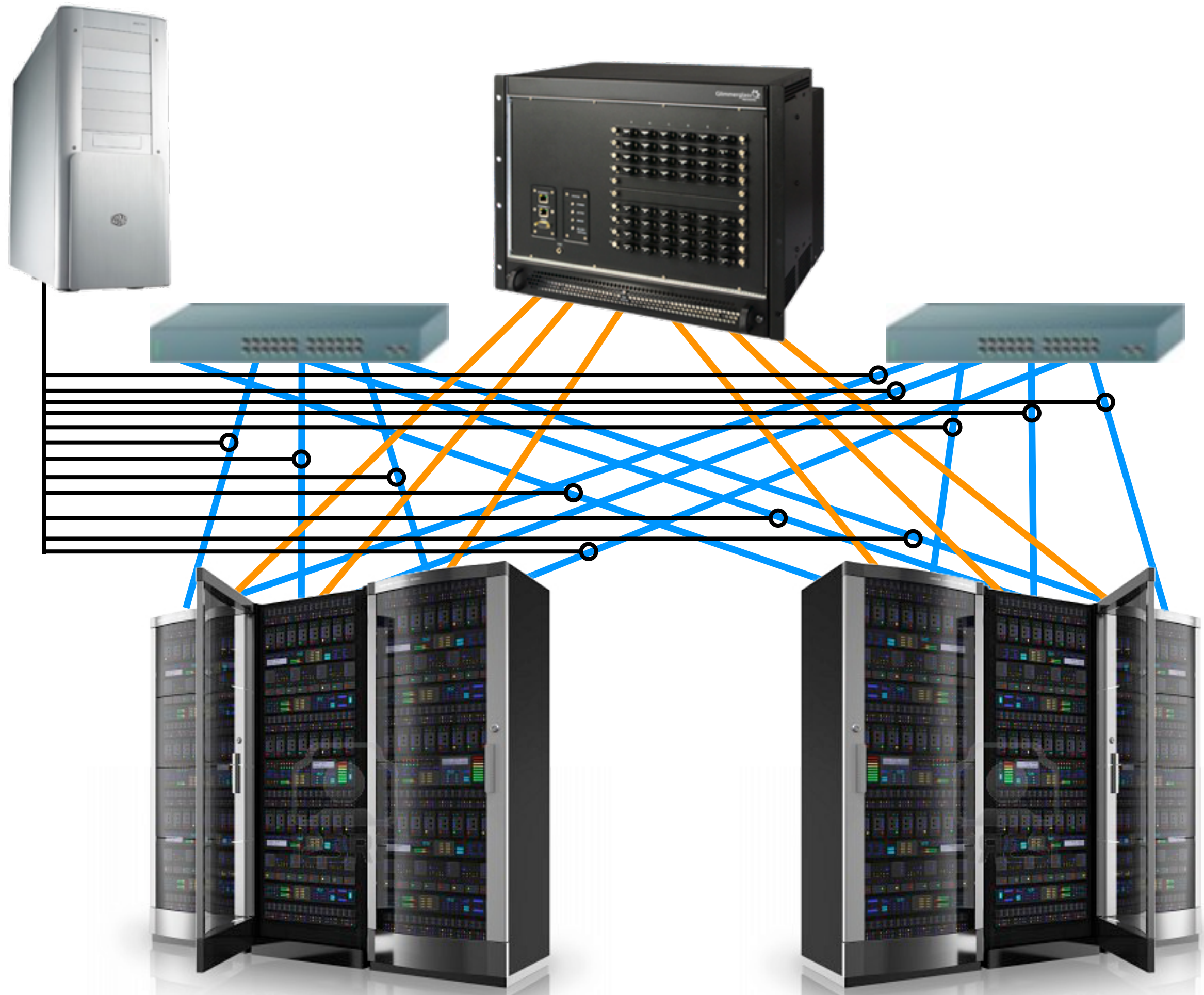


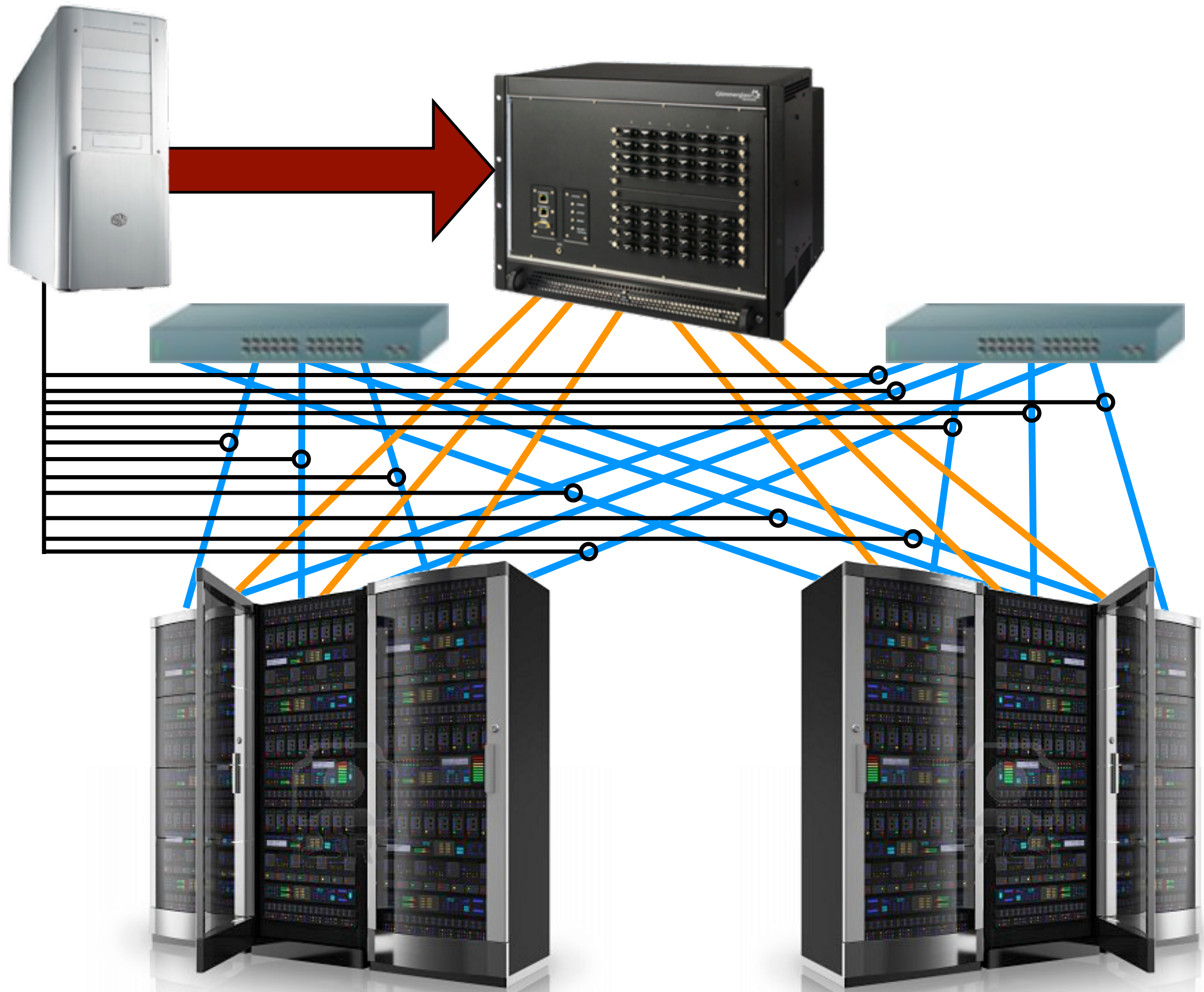




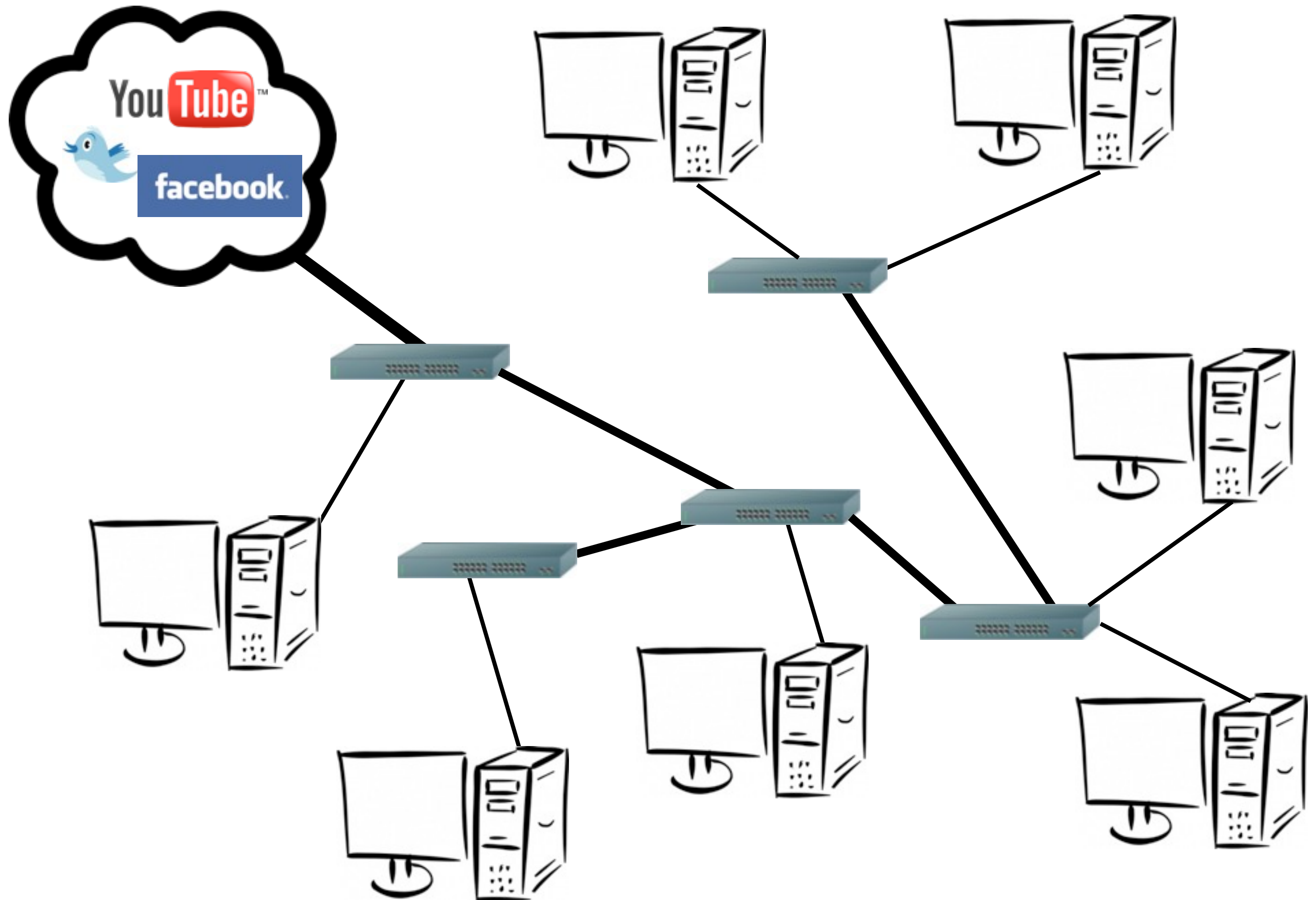


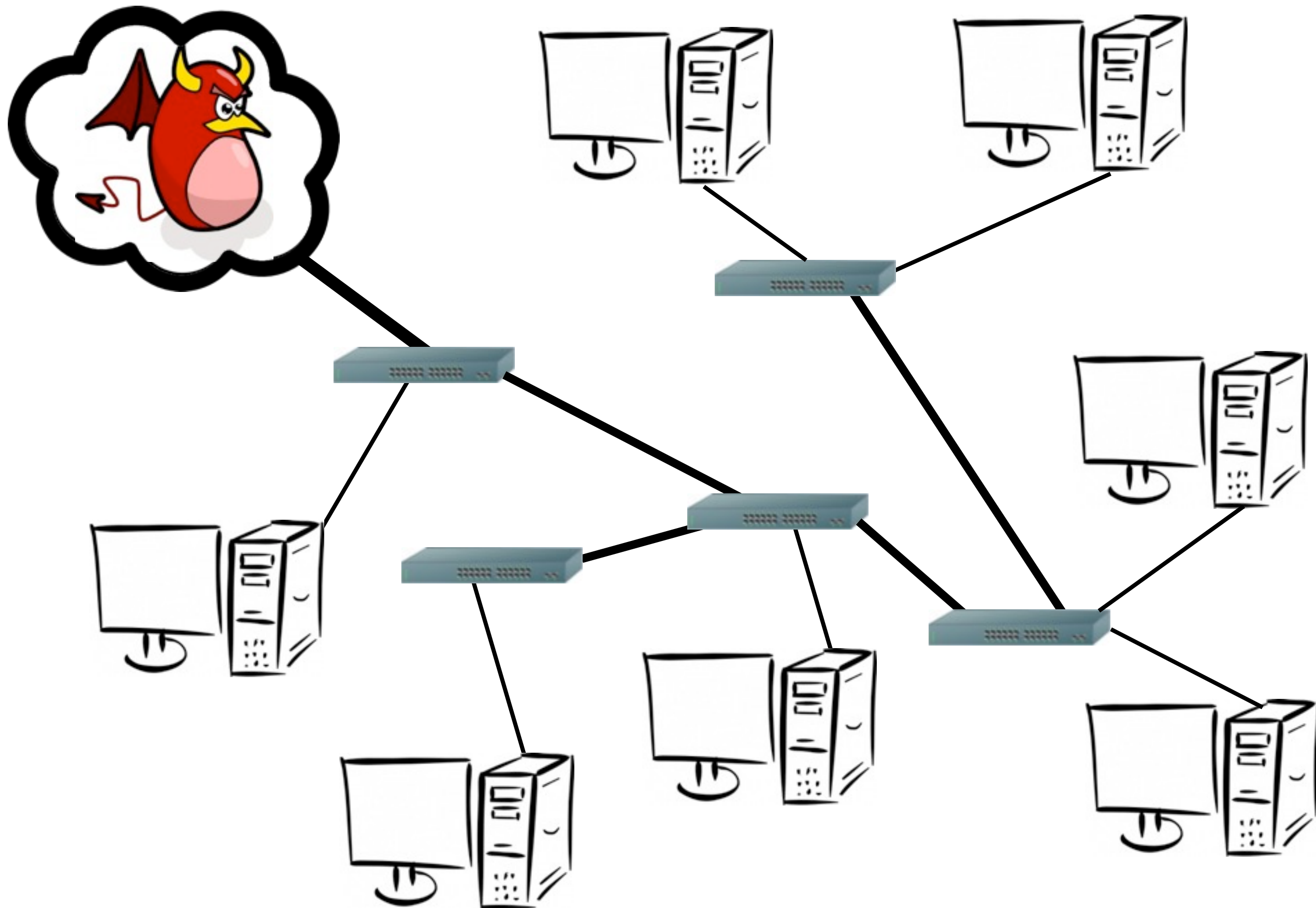


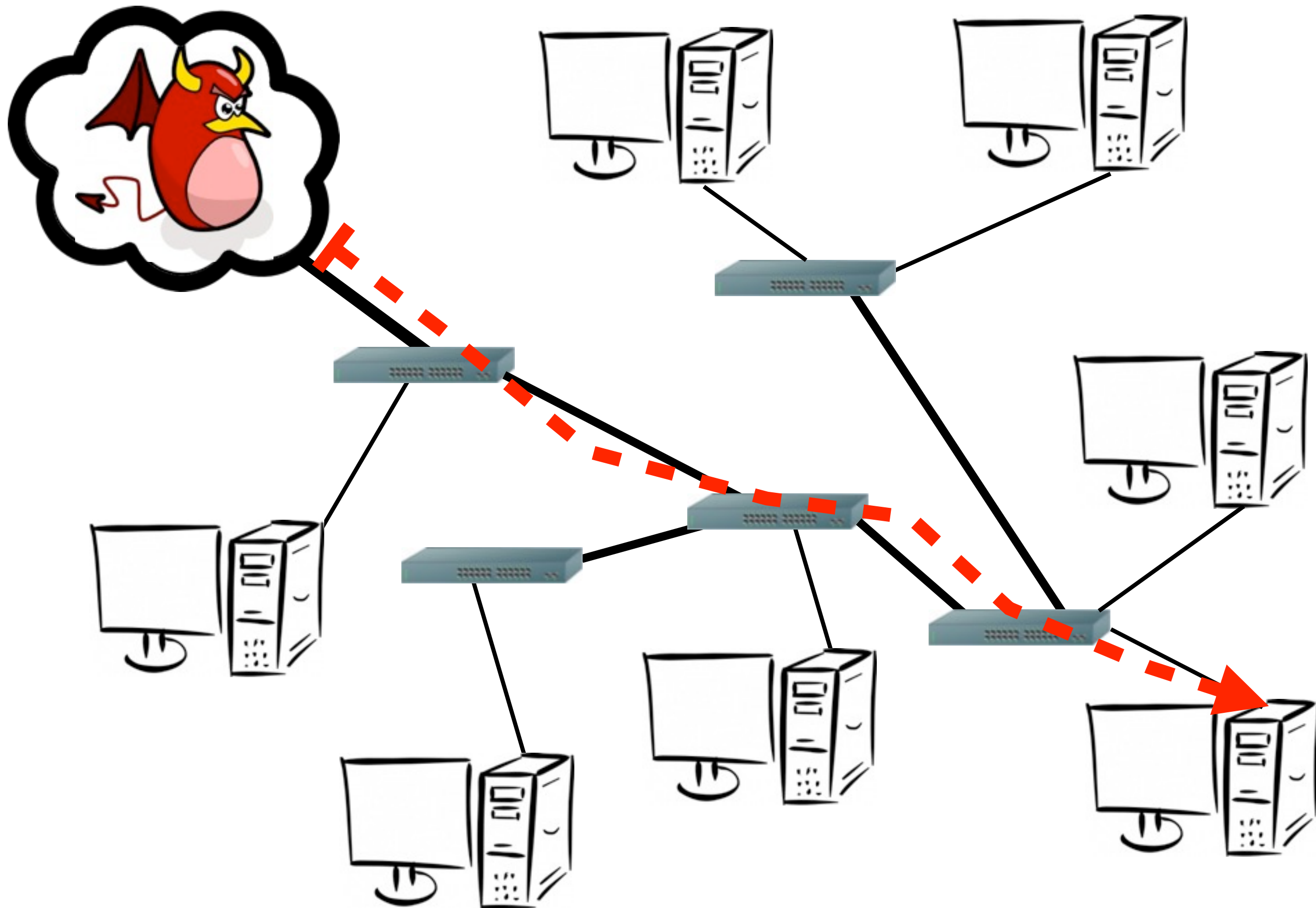


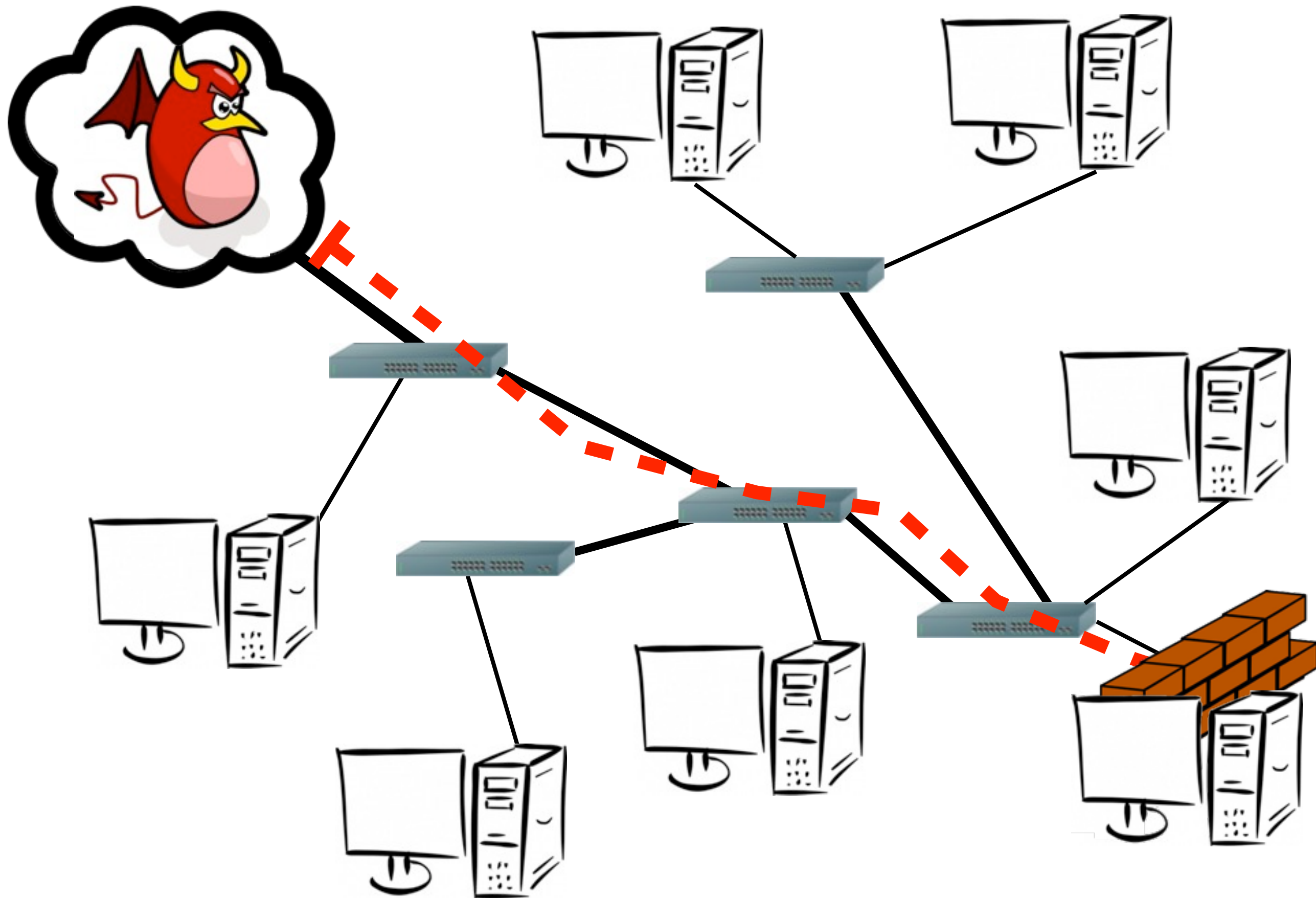


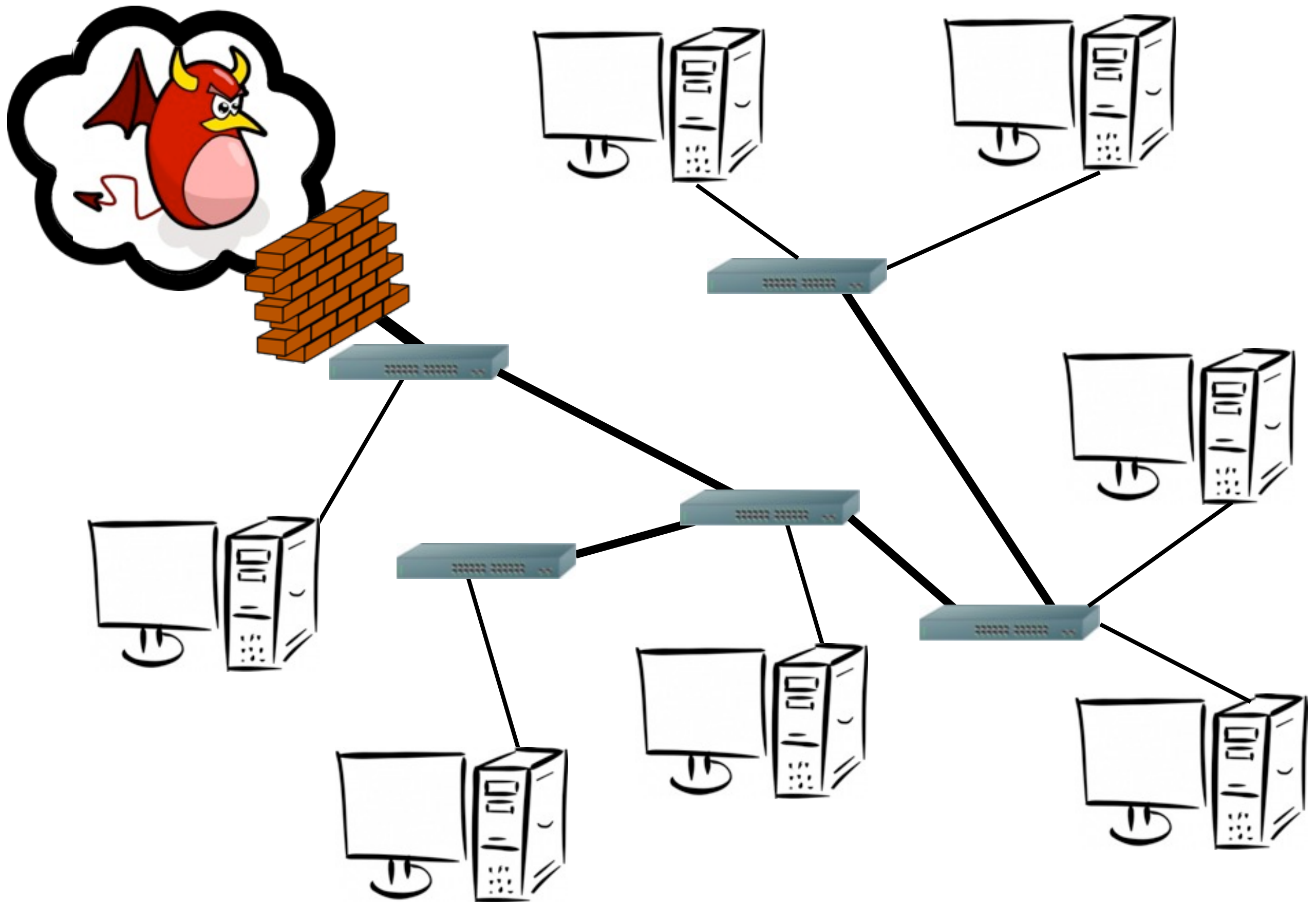
Defending from attacks







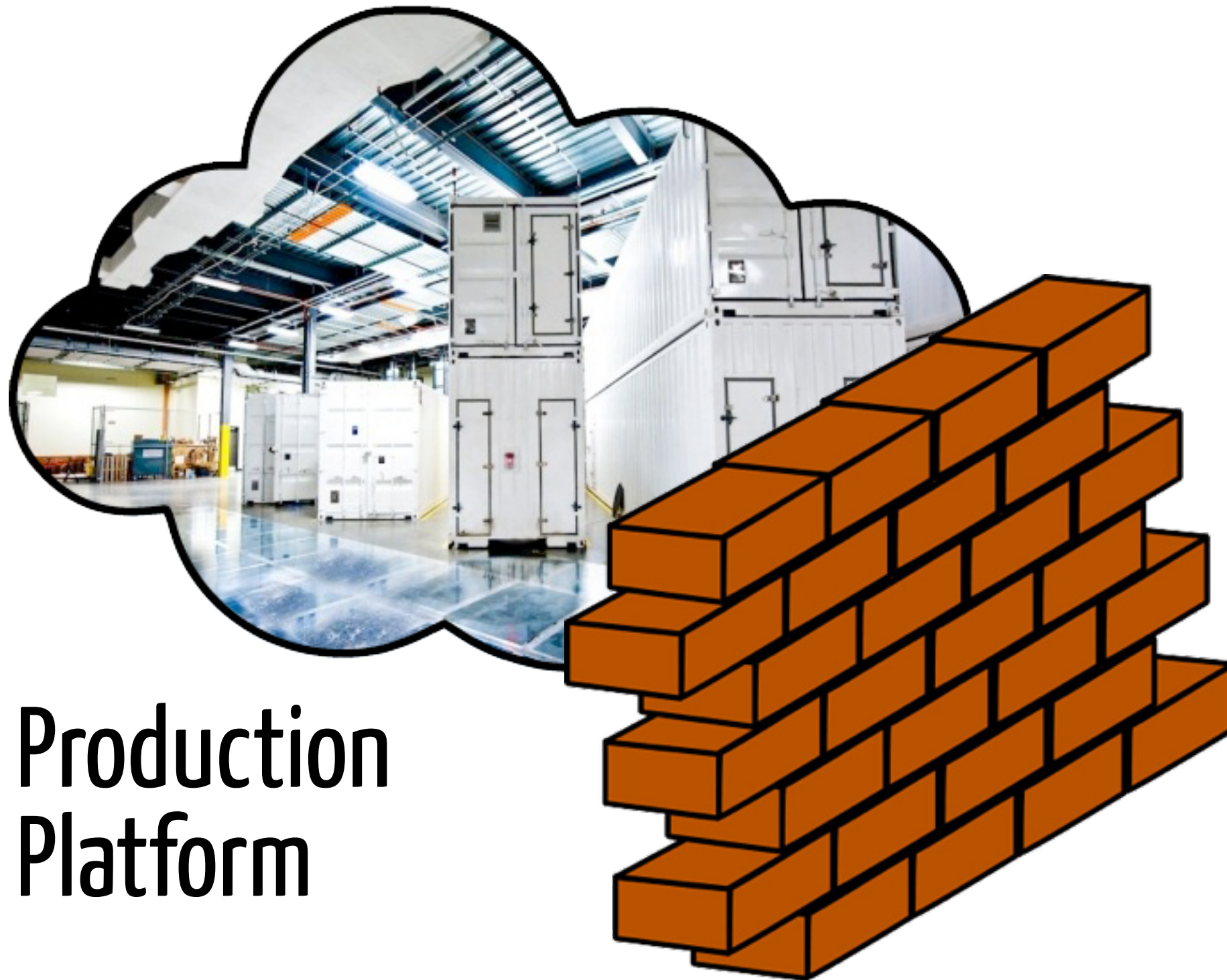




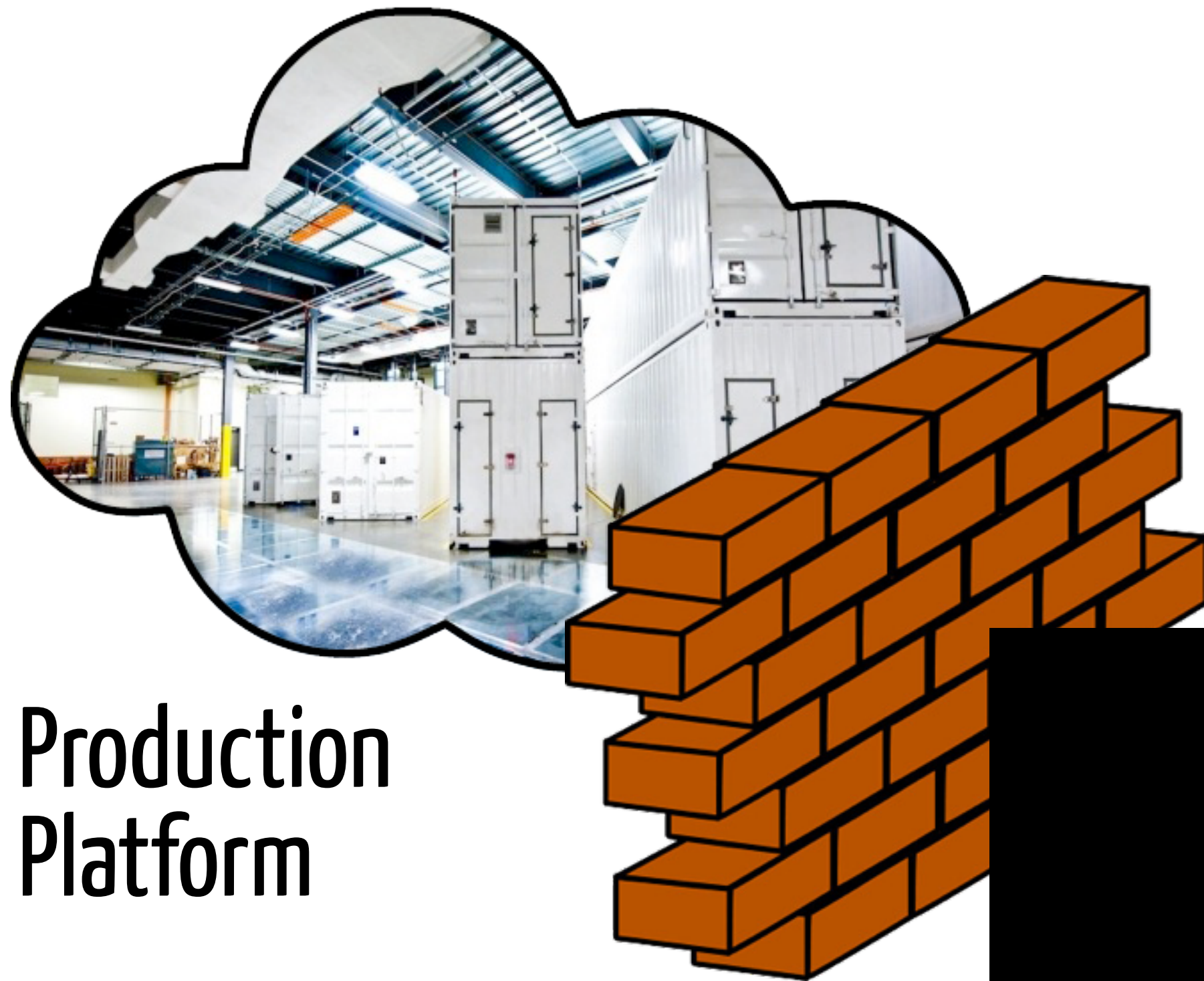
Untrusted VMs



Production Platform



Production Platform



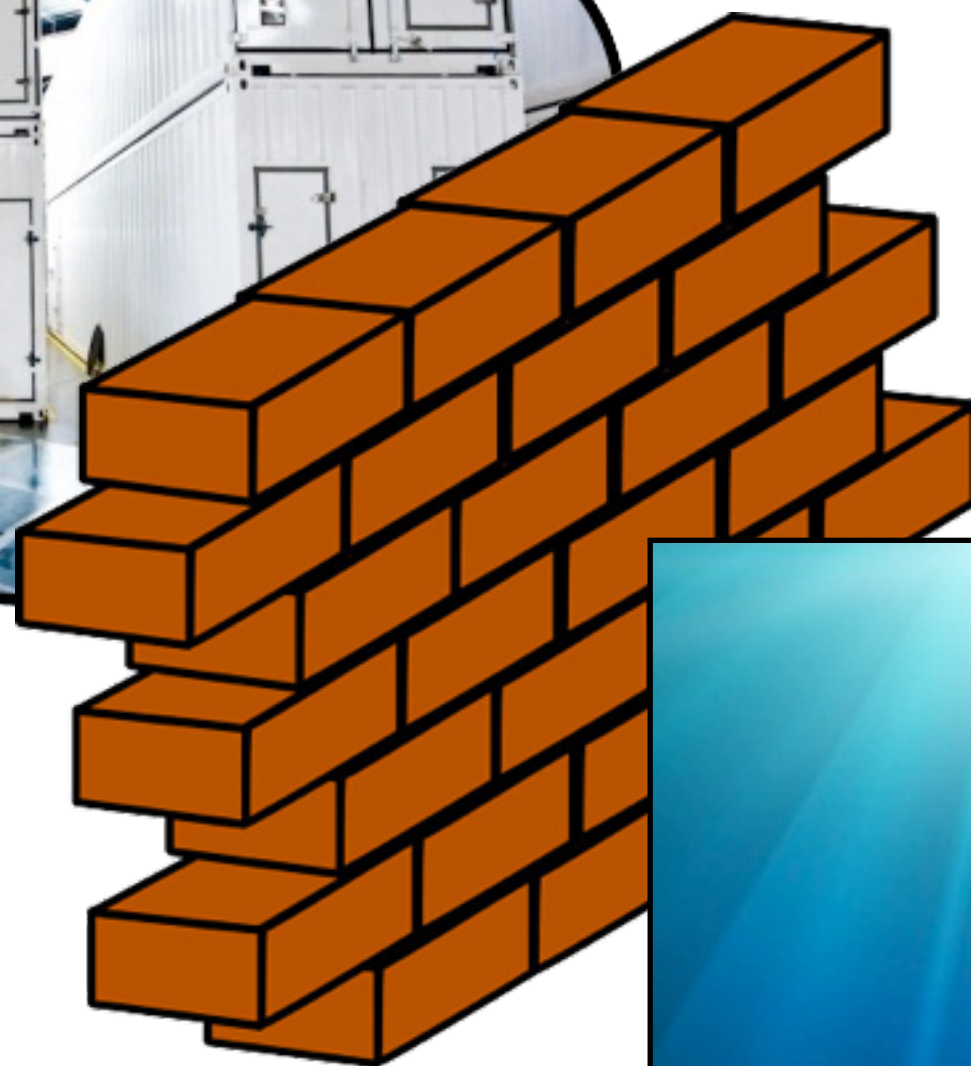
Production Platform

Boot Service





Production Platform

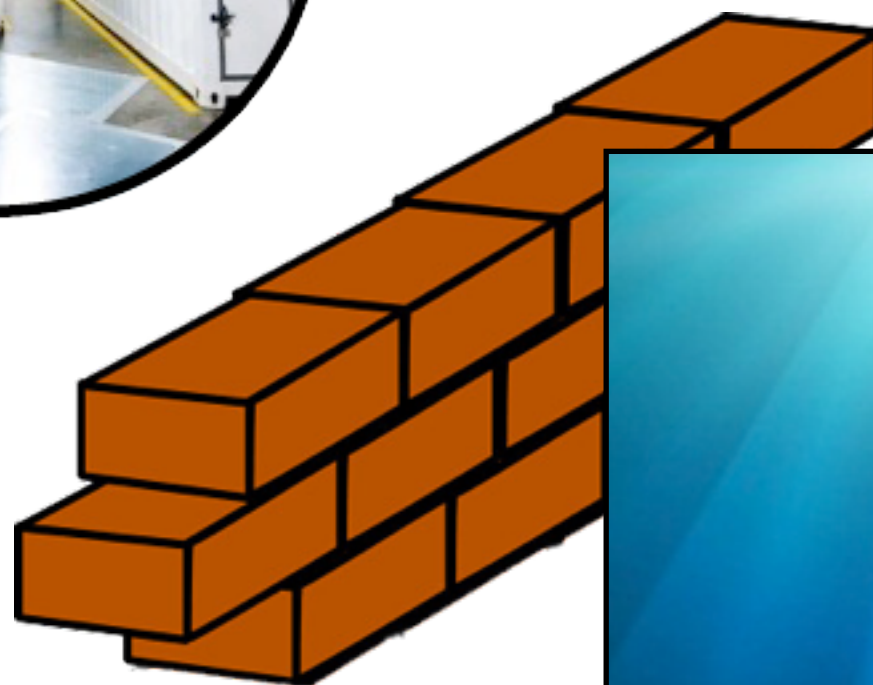


Boot Service

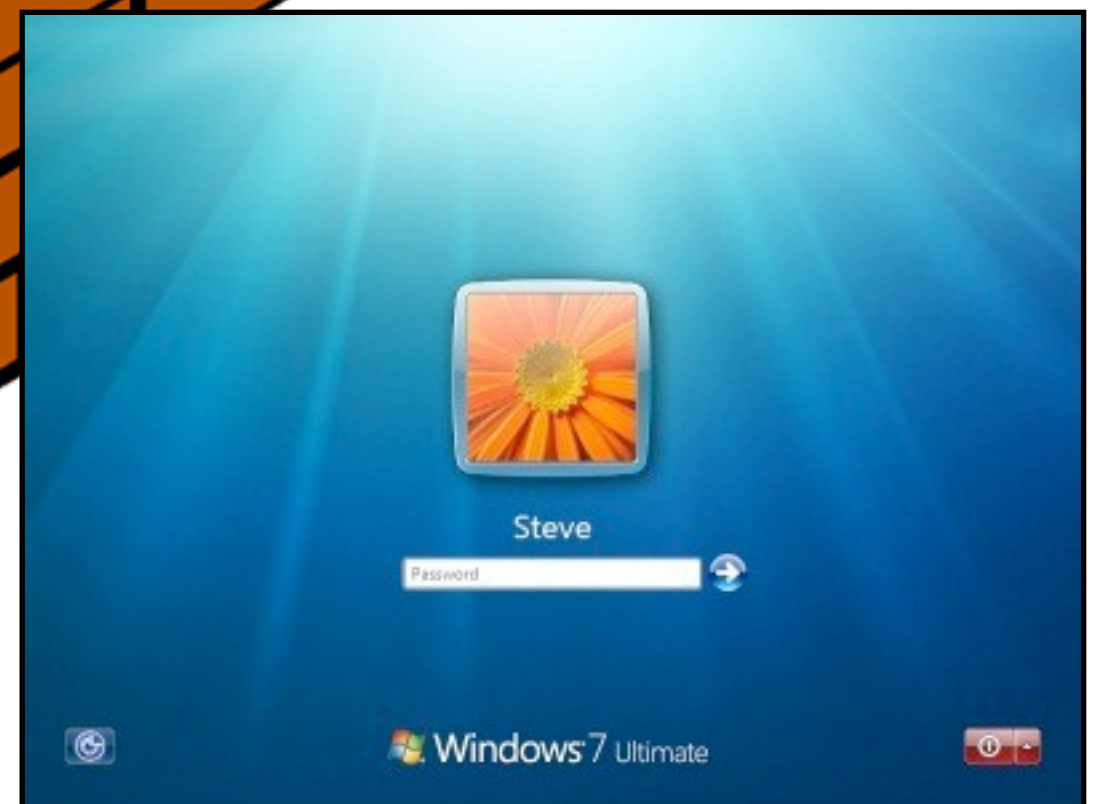




Production Platform

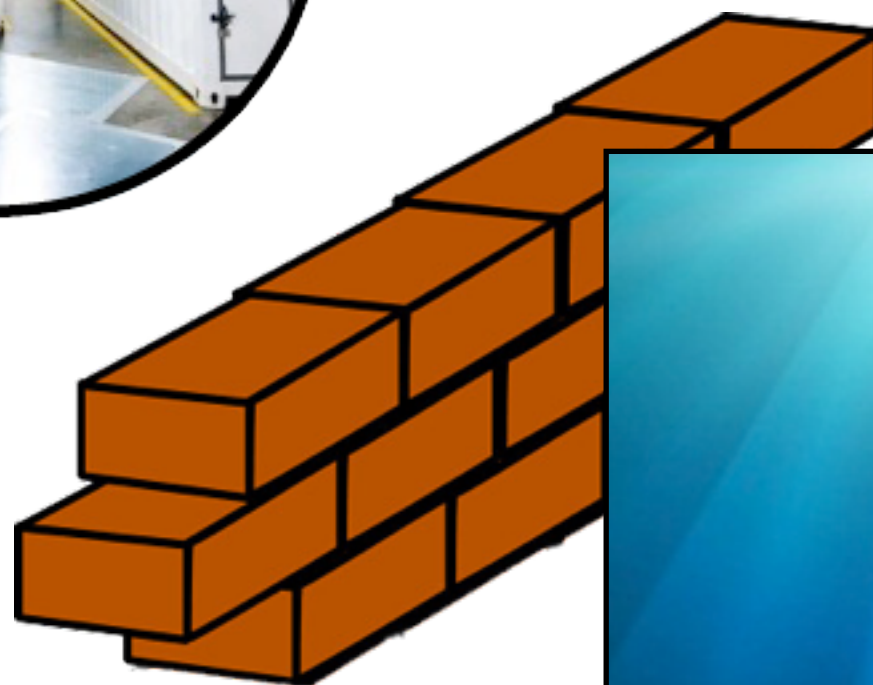


Boot Service

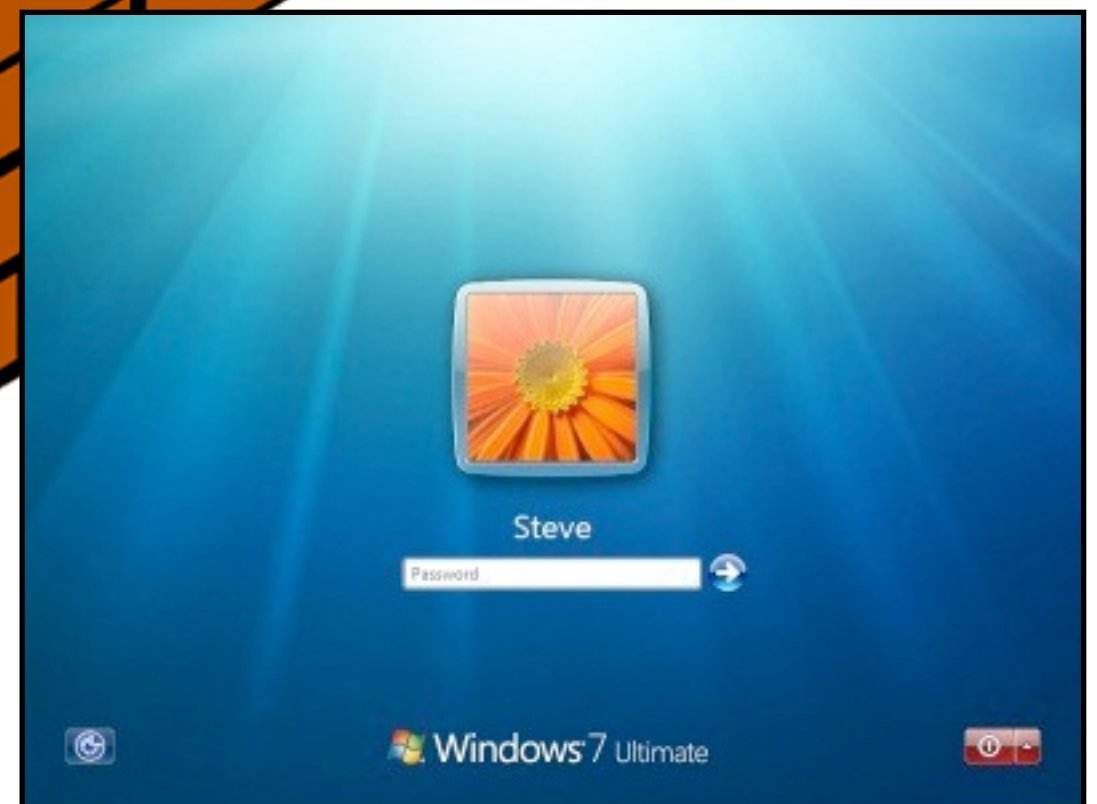




Production Platform

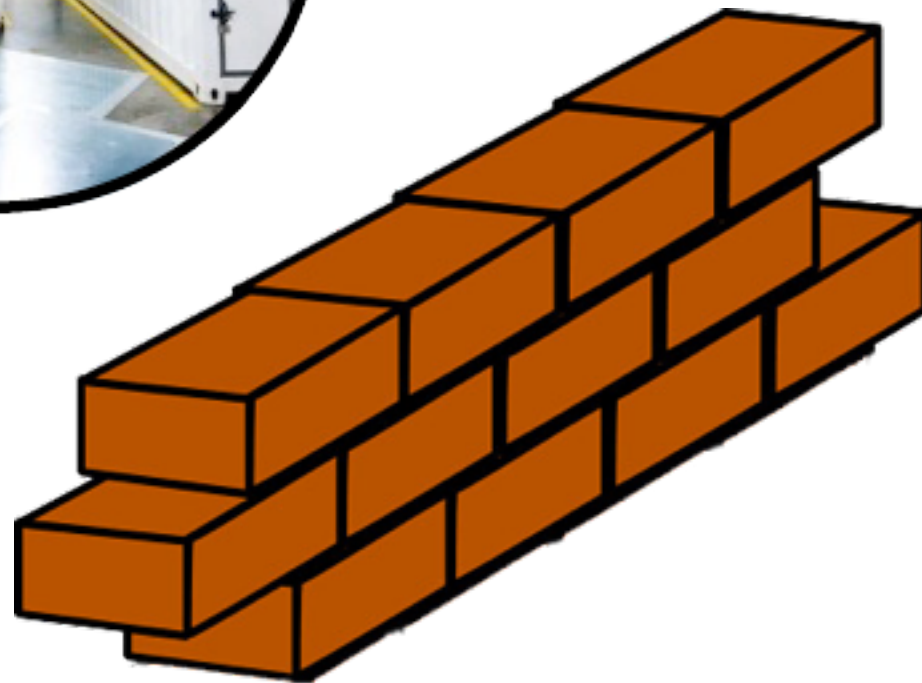


Boot Service

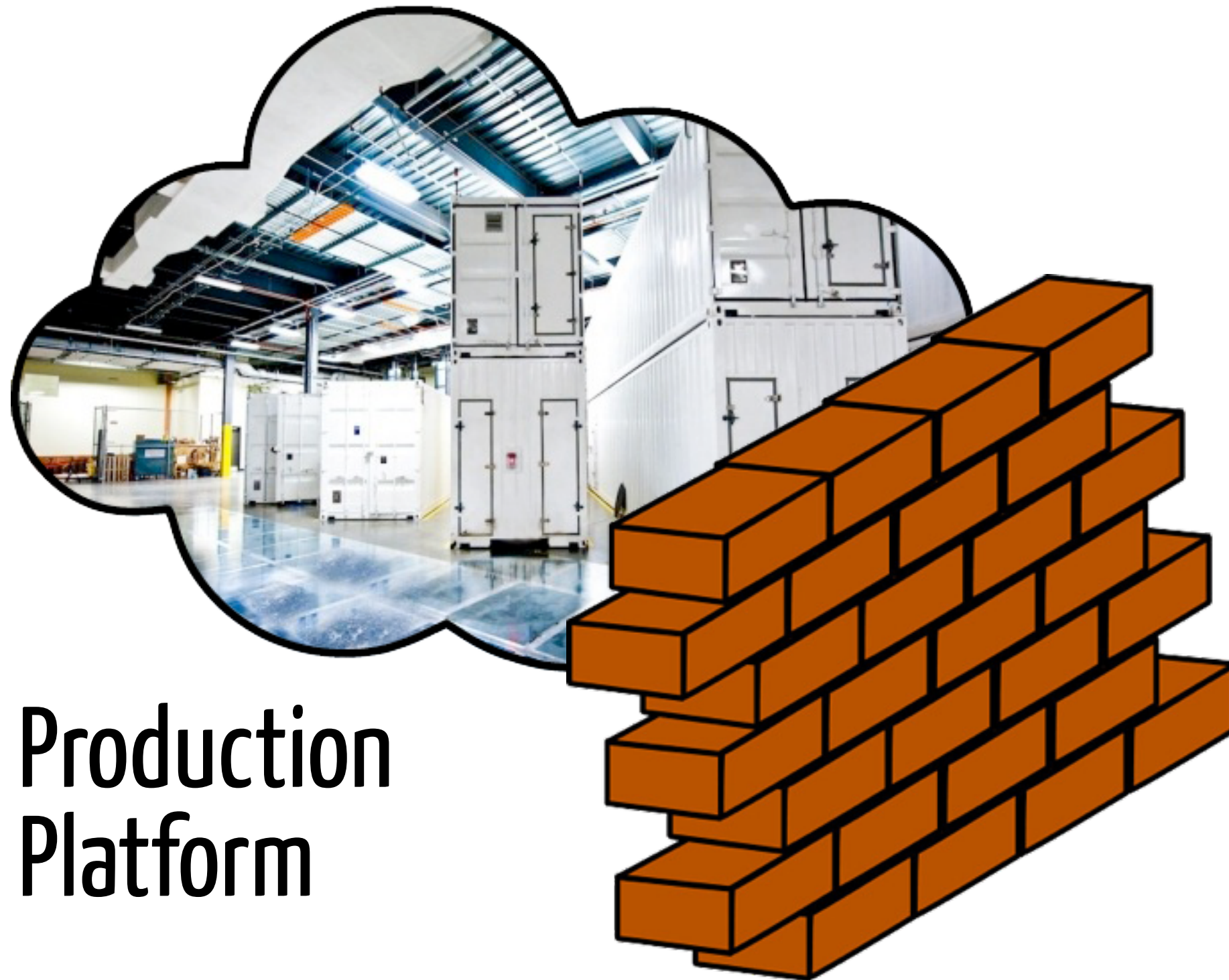




Production Platform



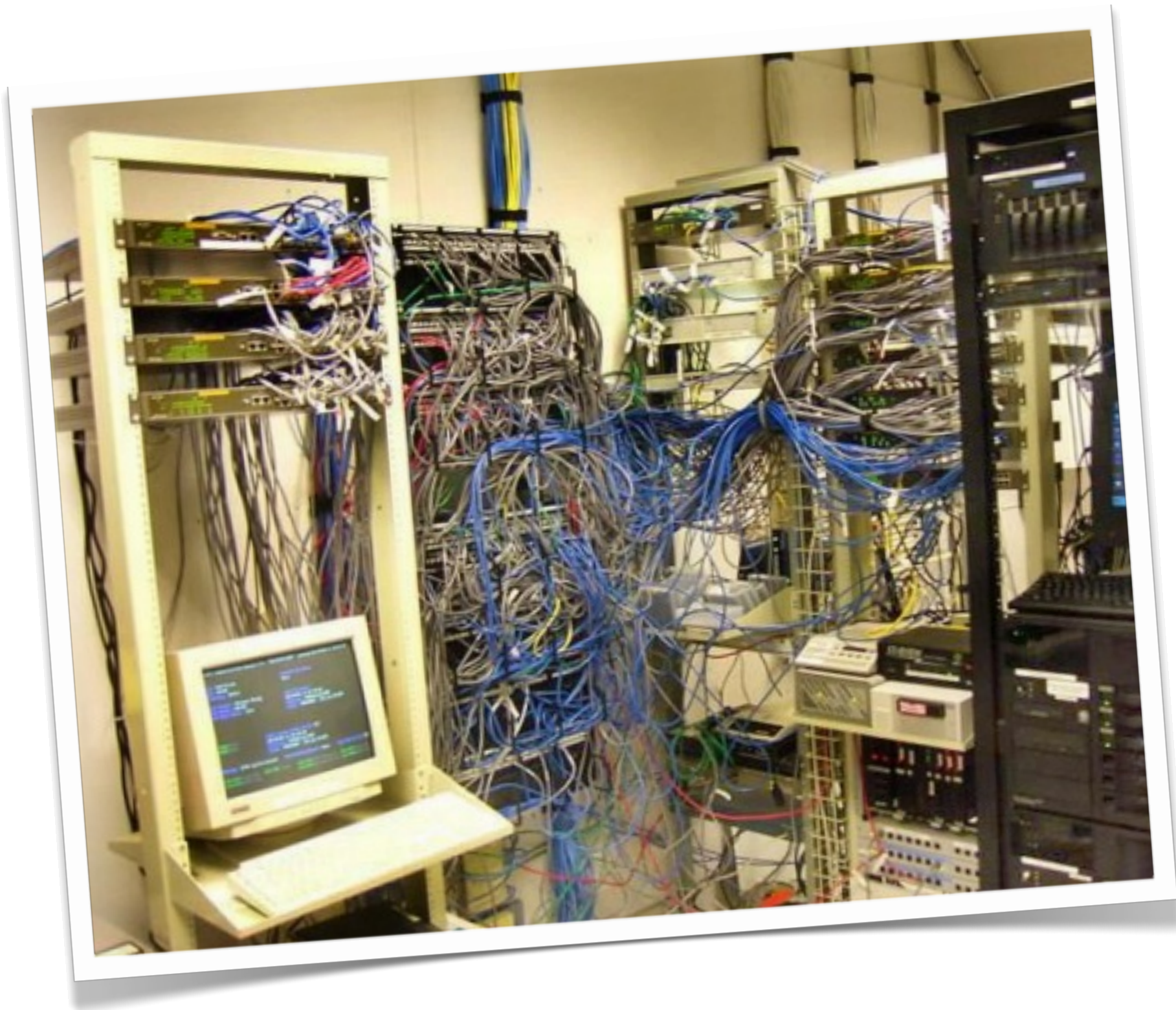
Boot Service



**Production
Platform**

**Boot
Service**

Proposal



Participatory Networking

Participatory Networking

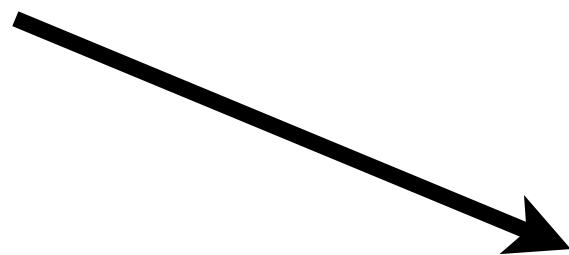
Participatory Networking



PANE

Participatory Networking

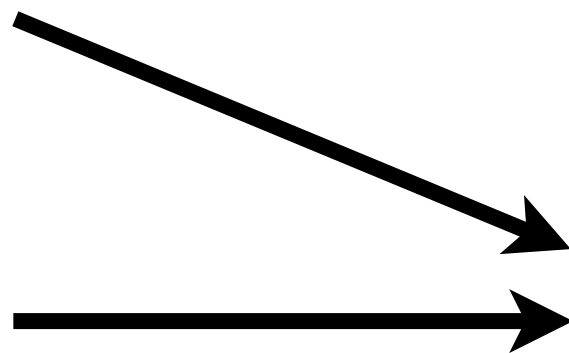
1. Requests



PANE

Participatory Networking

- 1. Requests
- 2. Hints



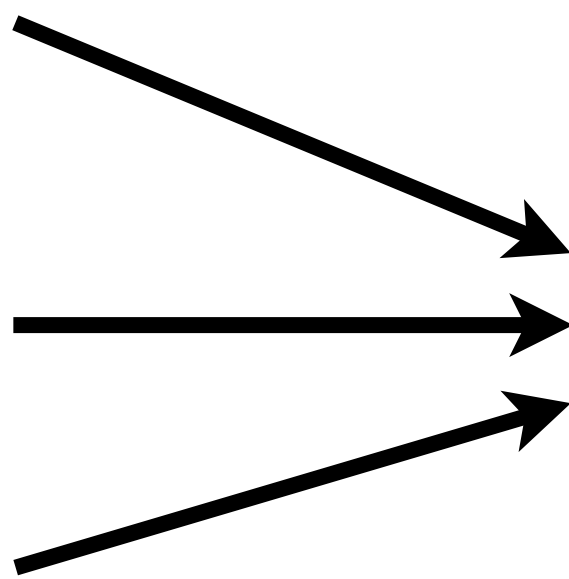
PANE

Participatory Networking

1. Requests

2. Hints

3. Queries



PANE

Participatory Networking

Participatory Networking

Safe?

Participatory Networking

Safe?

Secure?

Participatory Networking

Safe?

Secure?

Fair?

Participatory Networking

Safe?

Secure?

Fair?

Practical?

Participatory Networking

Safe?

Secure?

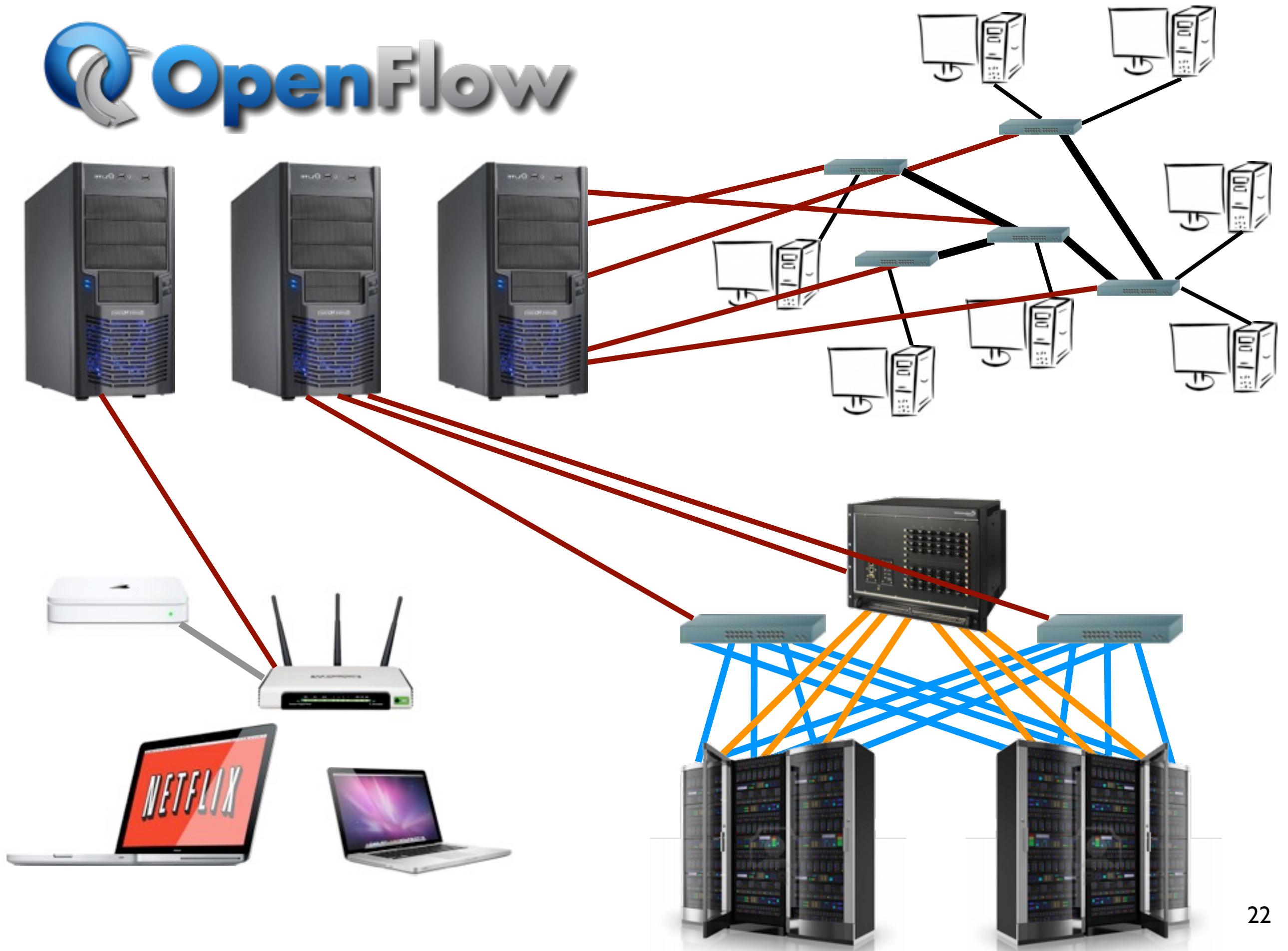
Fair?

Practical?

Efficient?



OpenFlow



Participatory Networking

Participatory Networking

- **End-user API for SDNs**

Participatory Networking

- End-user API for SDNs
- Exposes existing mechanisms

Participatory Networking

- End-user API for SDNs
- Exposes existing mechanisms
- No effect on unmodified applications

Outline of PANE

1. Privilege Delegation Semantics

Outline of PANE

1. Privilege Delegation Semantics

2. Protocol Sketch

Outline of PANE

1. Privilege Delegation Semantics
2. Protocol Sketch
3. Dynamic Flow Processing

Outline of PANE

1. Privilege Delegation Semantics
2. Protocol Sketch
3. Dynamic Flow Processing
4. Current Status

Outline of PANE

Privilege Delegation Semantics

Shares

Shares

Flowgroup

--	--

Shares

Flowgroup

src=128.12/16

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

	Privileges
--	------------

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

Privileges
deny, allow

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

Privileges
deny, allow

bandwidth: 5Mb/s

limit: 10Mb/s

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

Privileges

deny, allow

bandwidth: 5Mb/s

limit: 10Mb/s

hint

query

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

Speakers

Privileges

deny, allow

bandwidth: 5Mb/s

limit: 10Mb/s

hint

query

Shares

Flowgroup

$\text{src}=128.12/16 \wedge \text{dst.port} \leq 1024$

Speakers

Alice

Bob

Privileges

deny, allow

bandwidth: 5Mb/s

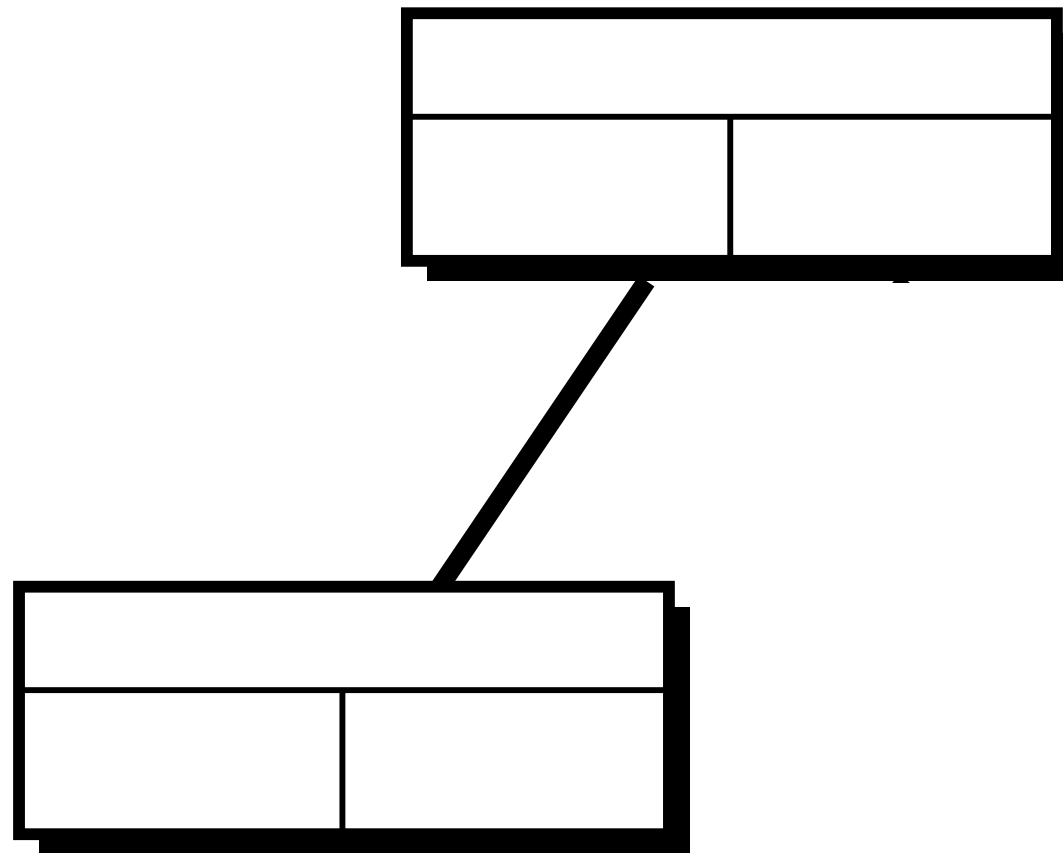
limit: 10Mb/s

hint

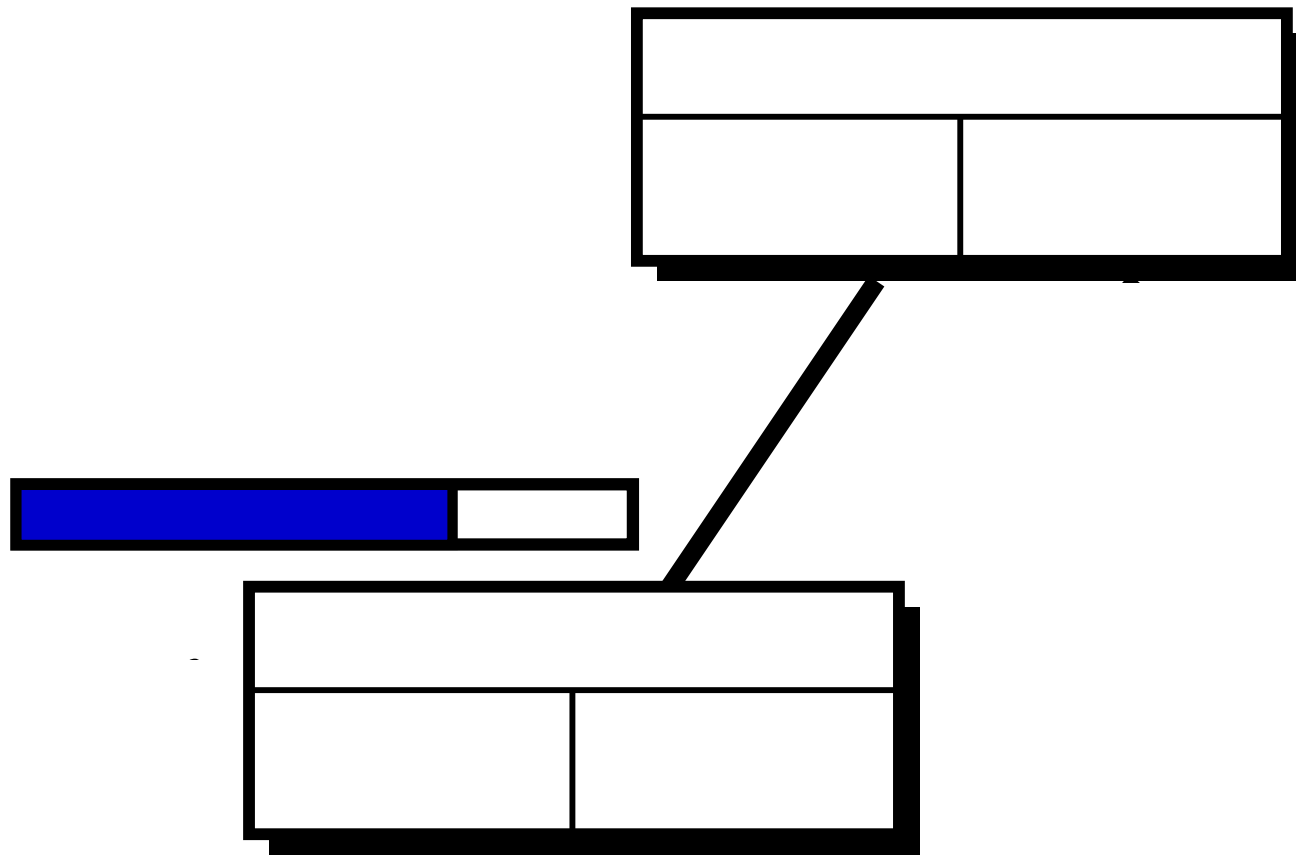
query

Shares

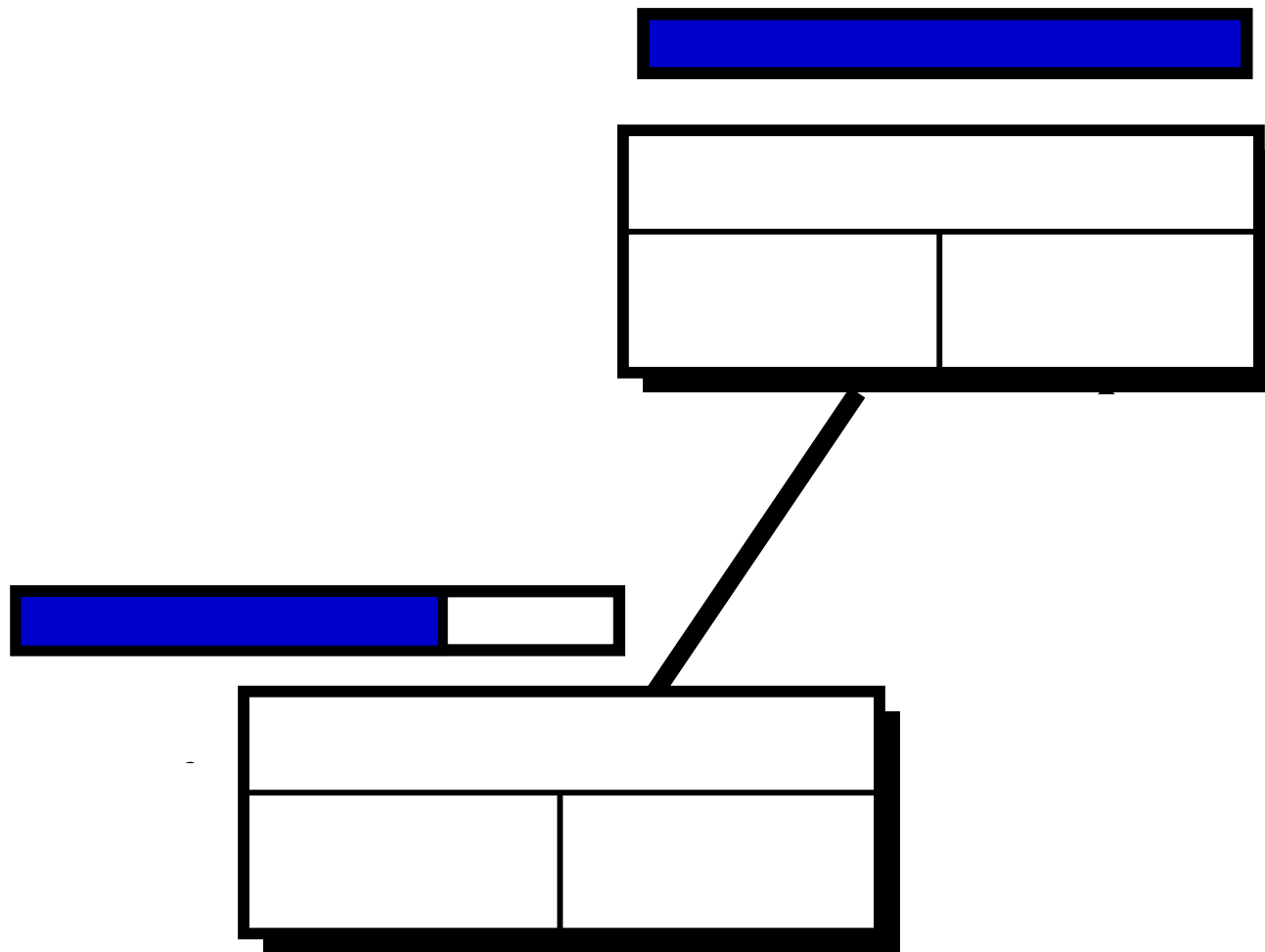
Delegation



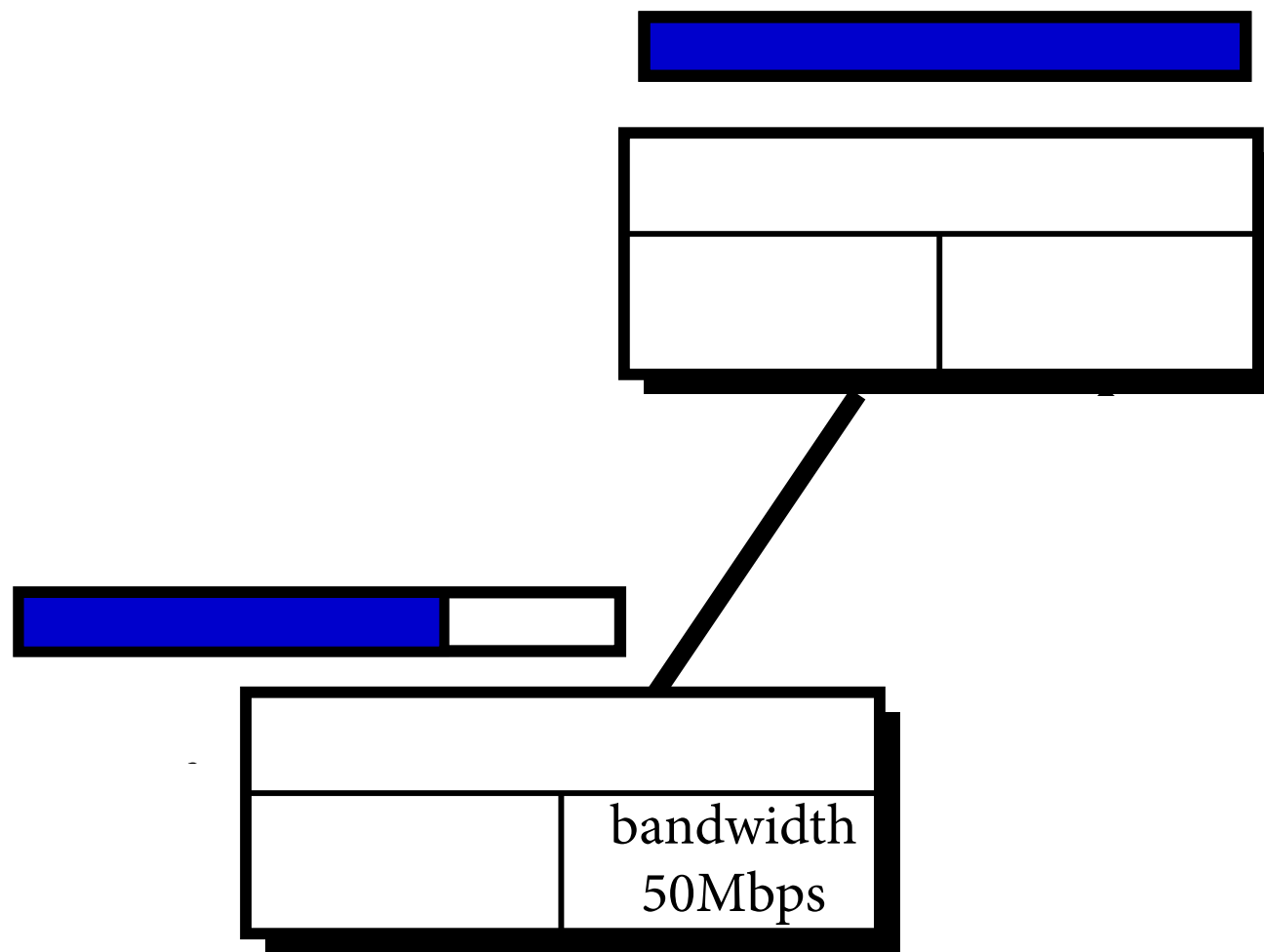
Delegation



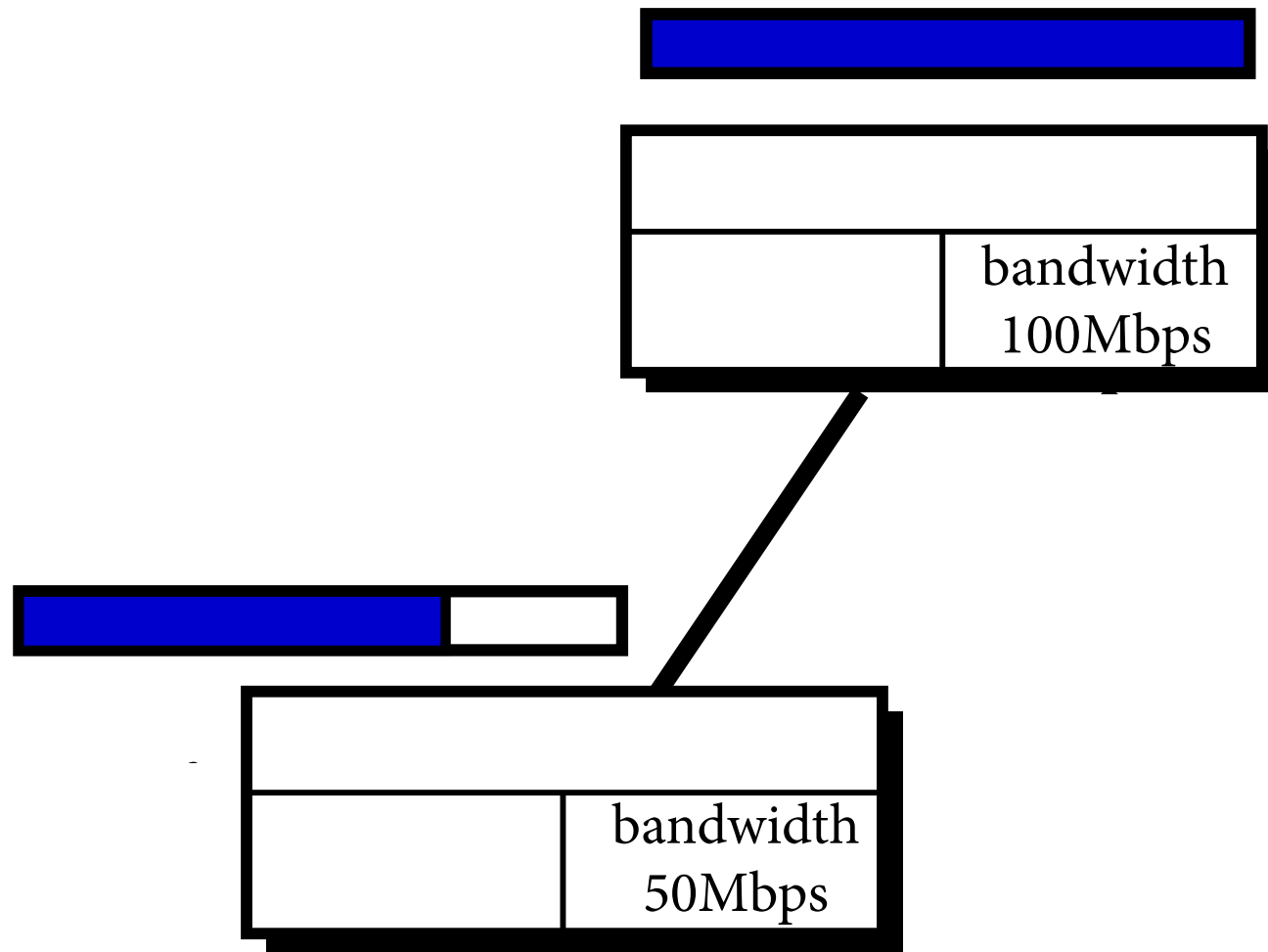
Delegation



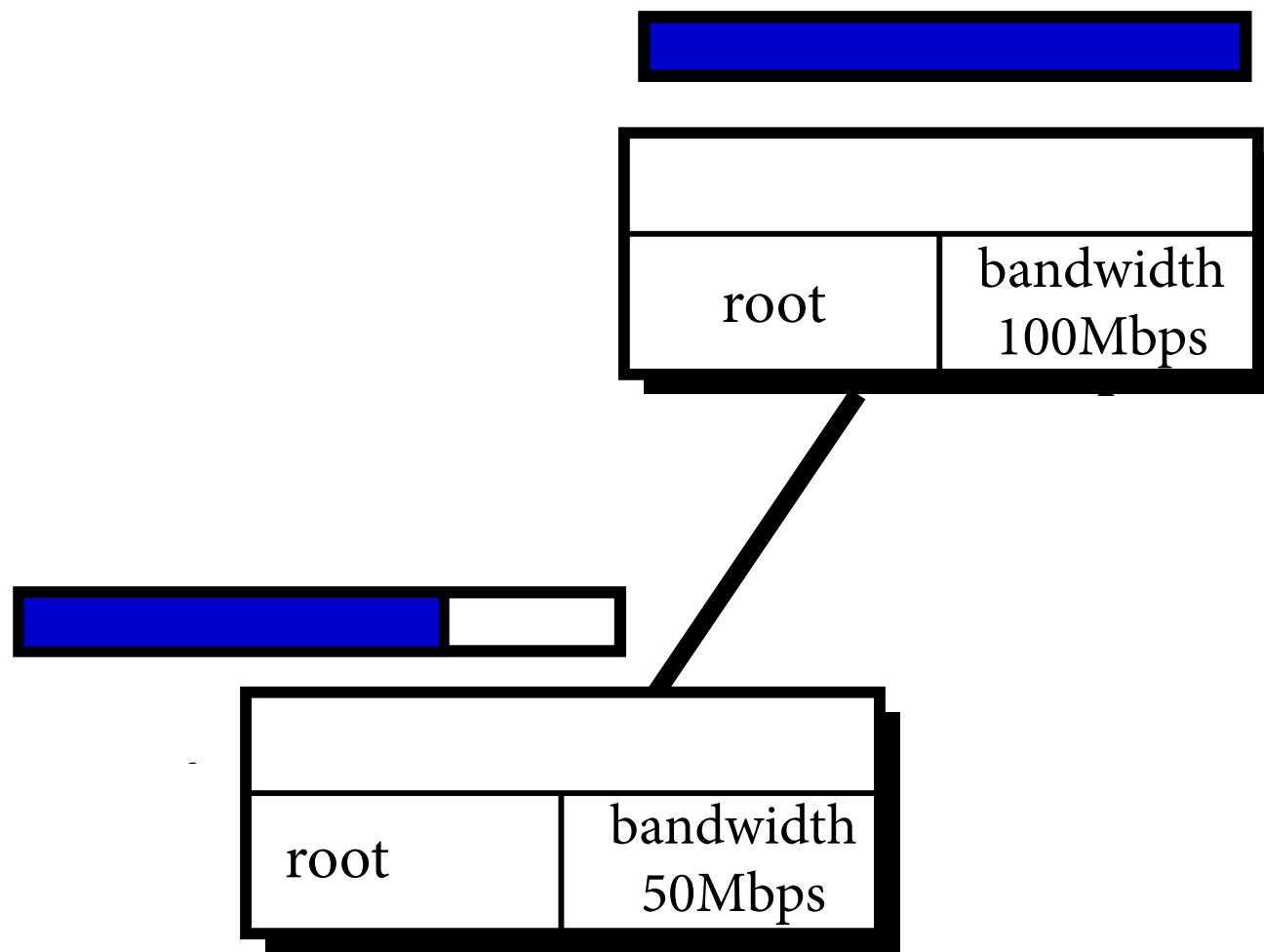
Delegation



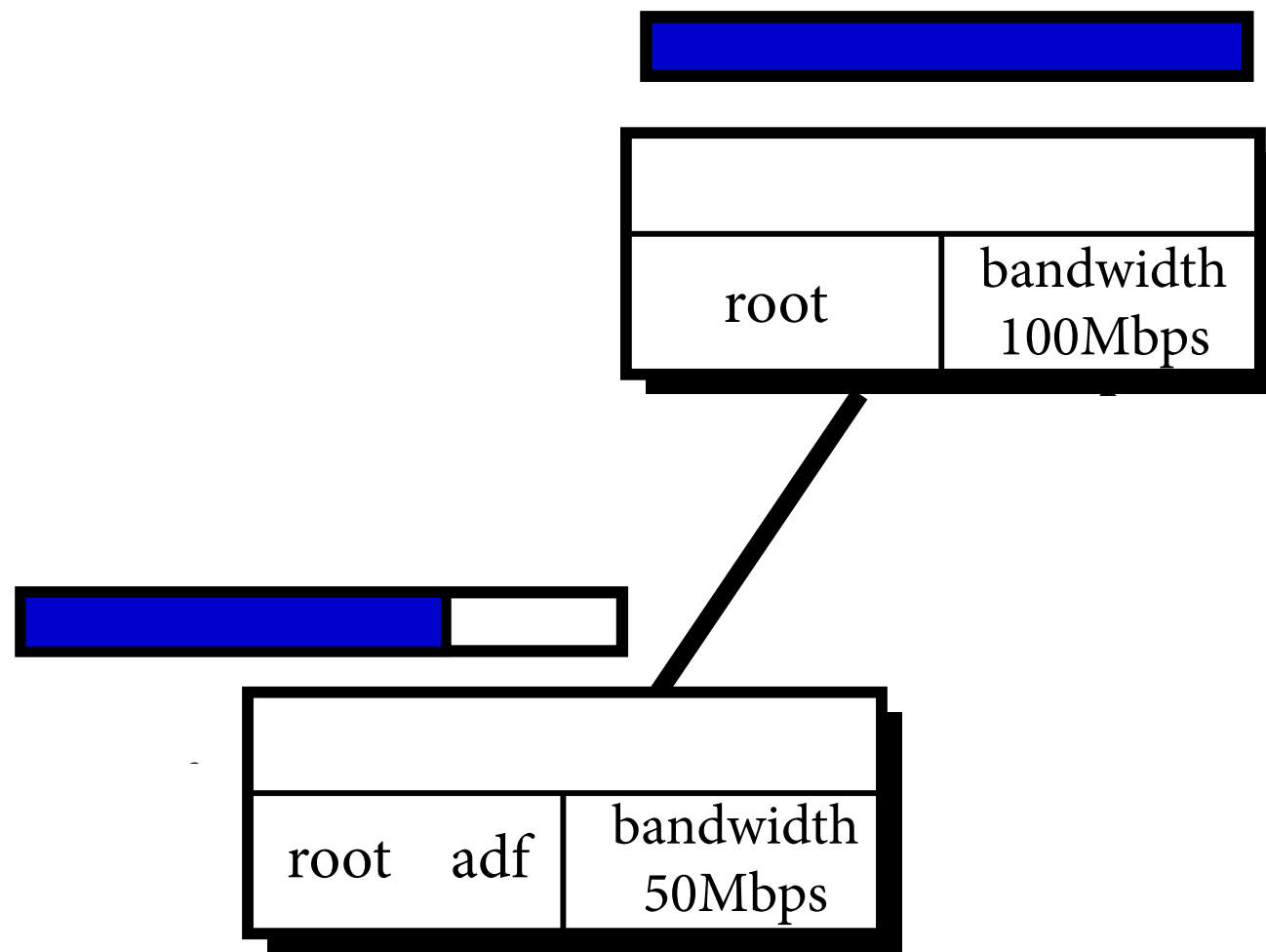
Delegation



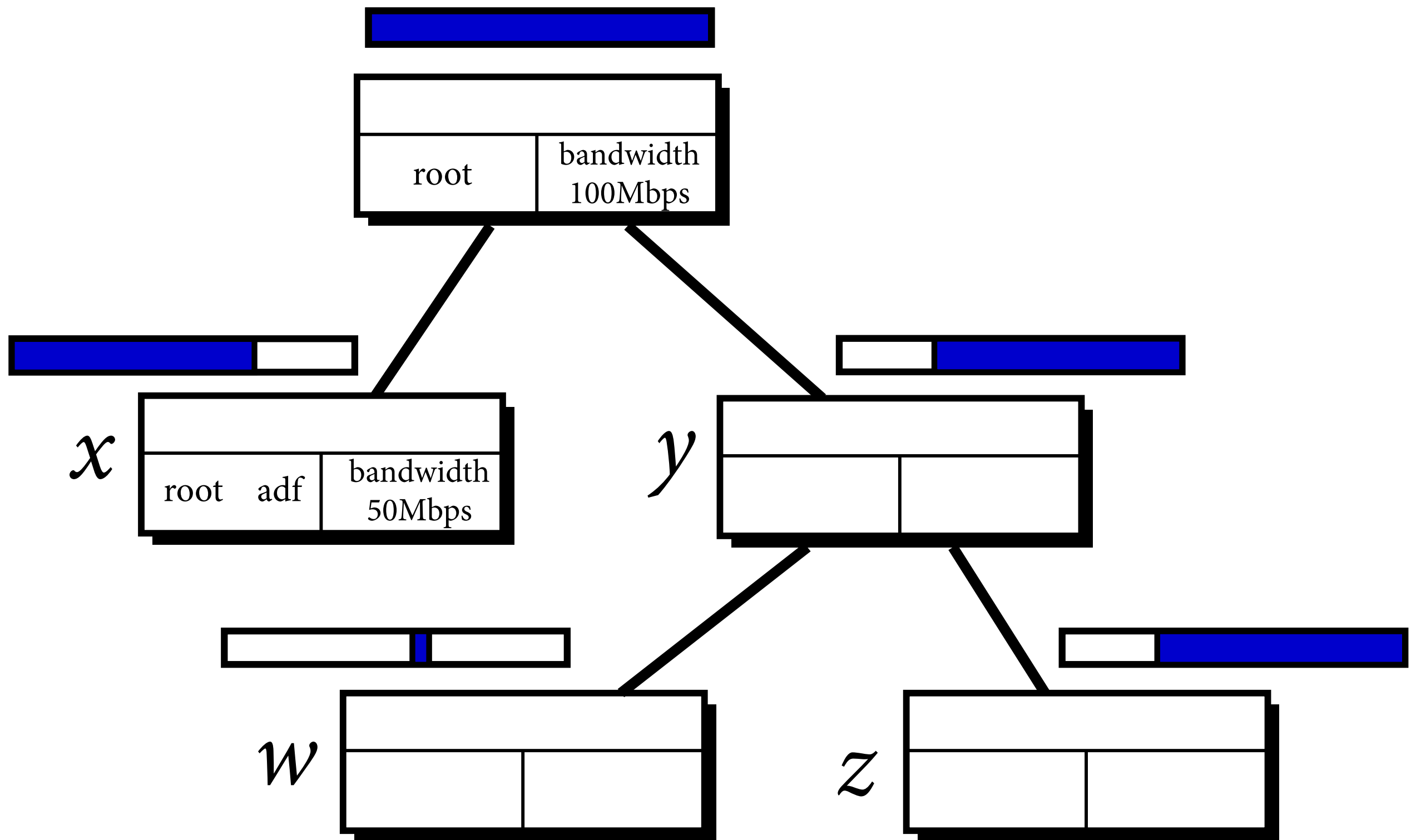
Delegation



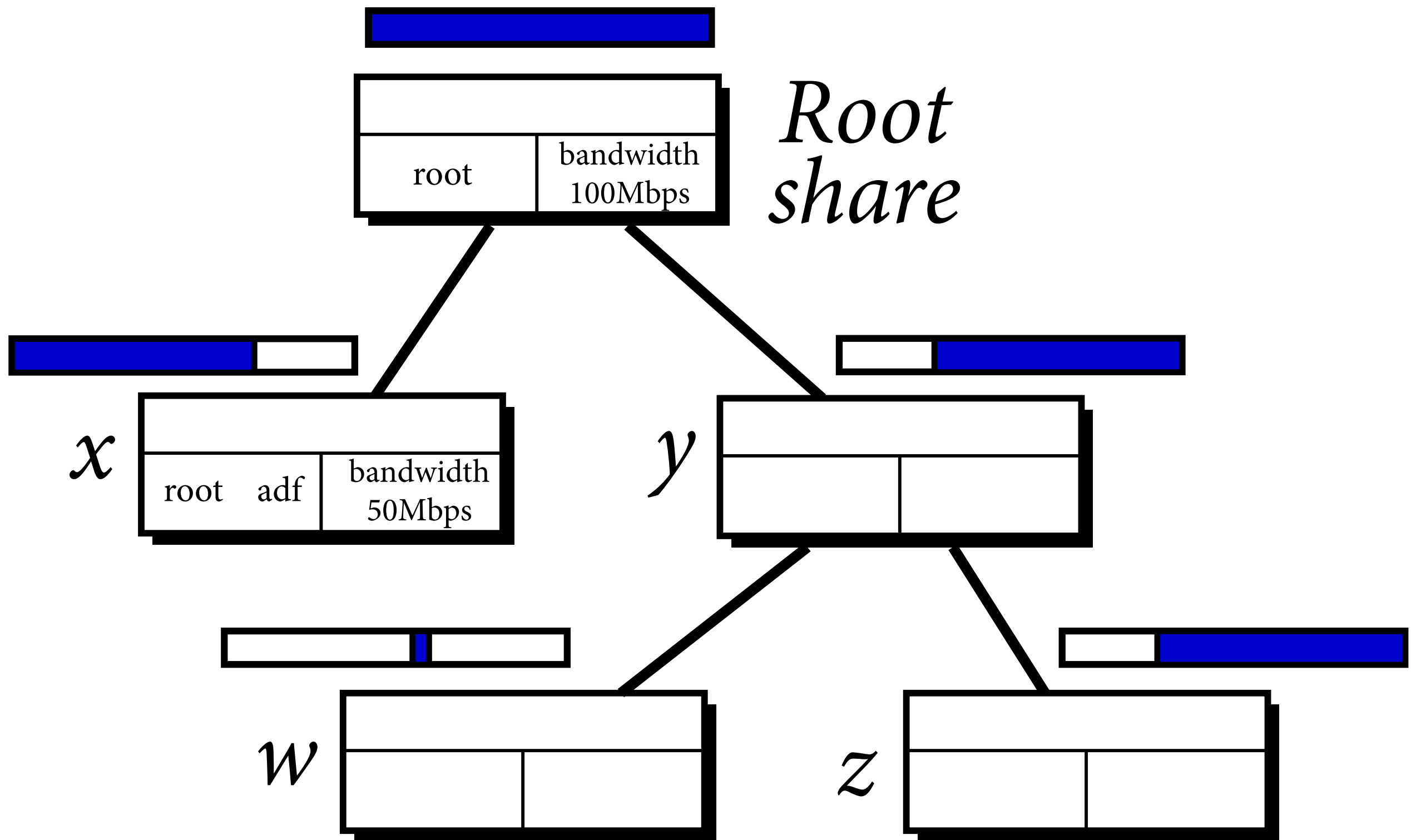
Delegation



Delegation



Delegation



Delegation

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice	deny, allow
Bob	bandwidth: 5Mb/s
	limit: 10Mb/s
	<i>hint</i>
	<i>query</i>

Dynamic Context

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice	deny, allow
Bob	bandwidth: 5Mb/s
	limit: 10Mb/s
	<i>hint</i>
	<i>query</i>

Dynamic Context



PANE

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice	deny, allow
Bob	bandwidth: 5Mb/s
	limit: 10Mb/s
	<i>hint</i>
	<i>query</i>

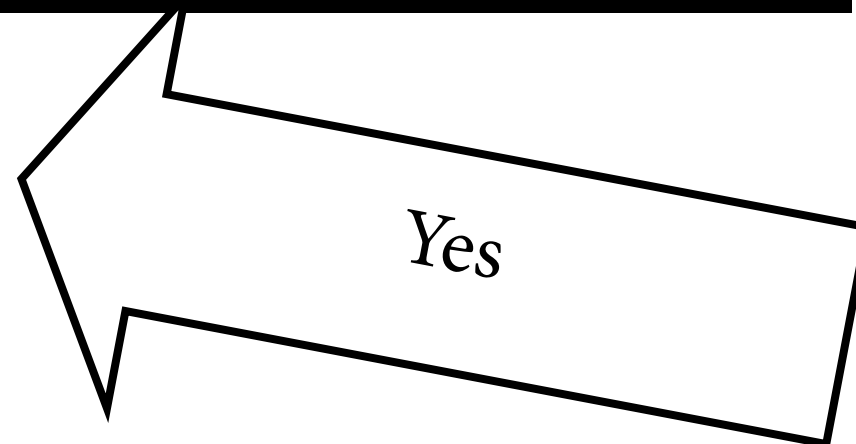
Reserve 2 Mbps
from now to +5min?

Dynamic Context



PANE

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice Bob	deny, allow bandwidth: 5Mb/s limit: 10Mb/s <i>hint</i> <i>query</i>



Yes

Dynamic Context



PANE

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice	deny, allow
Bob	bandwidth: 5Mb/s
	limit: 10Mb/s
	<i>hint</i>
	<i>query</i>

This traffic will be short and bursty

Dynamic Context



PANE

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice	deny, allow
Bob	bandwidth: 5Mb/s
	limit: 10Mb/s
	<i>hint</i>
	<i>query</i>

Dynamic Context



OK

PANE

Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice	deny, allow
Bob	bandwidth: 5Mb/s
	limit: 10Mb/s
	<i>hint</i>
	<i>query</i>

How much web traffic
in the last hour?

Dynamic Context



PANE

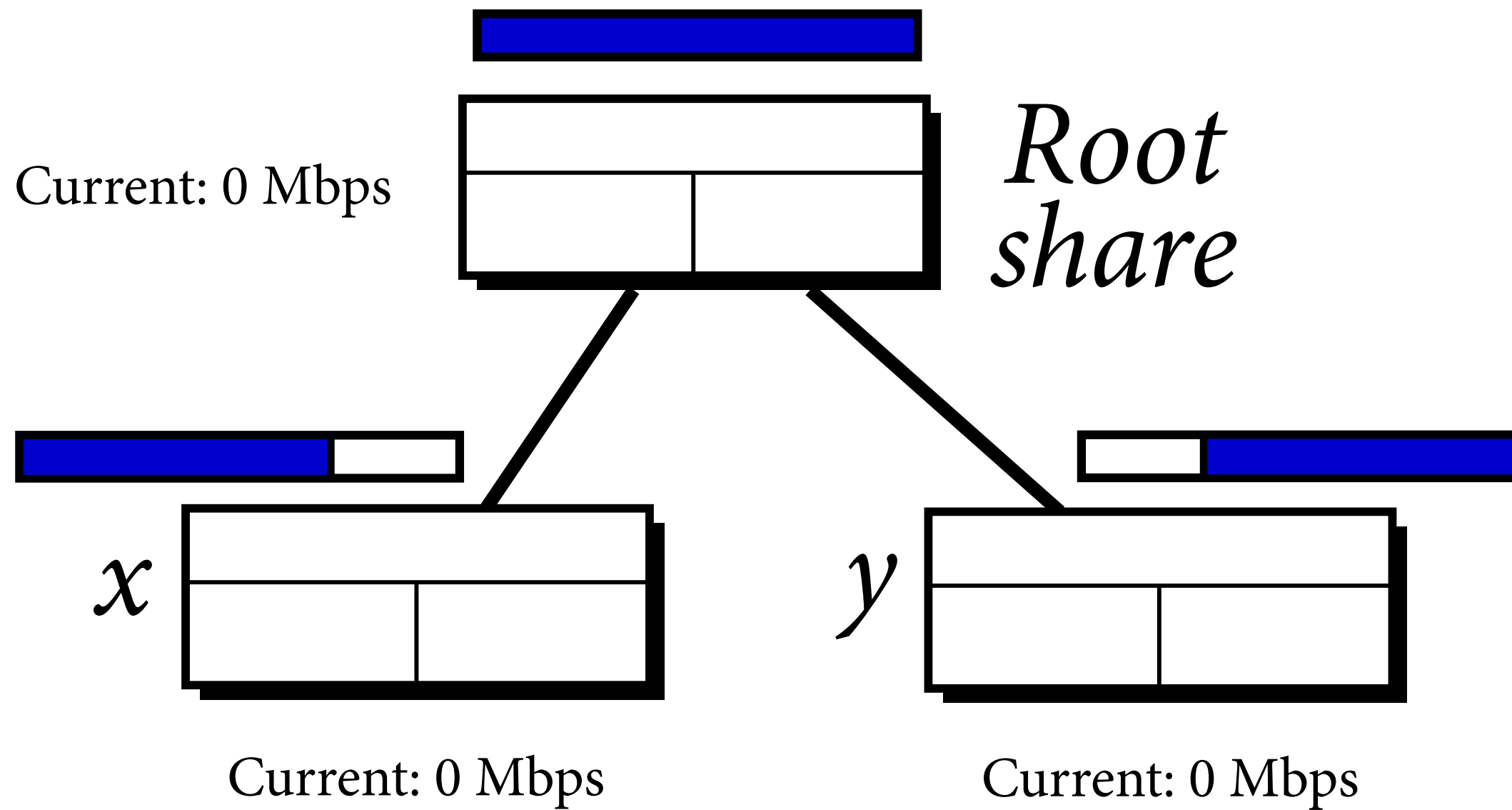
Flowgroup	
src=128.12/16 $\wedge$ dst.port $\leq$ 1024	
Speakers	Privileges
Alice Bob	deny, allow bandwidth: 5Mb/s limit: 10Mb/s <i>hint</i> <i>query</i>

67,560 bytes

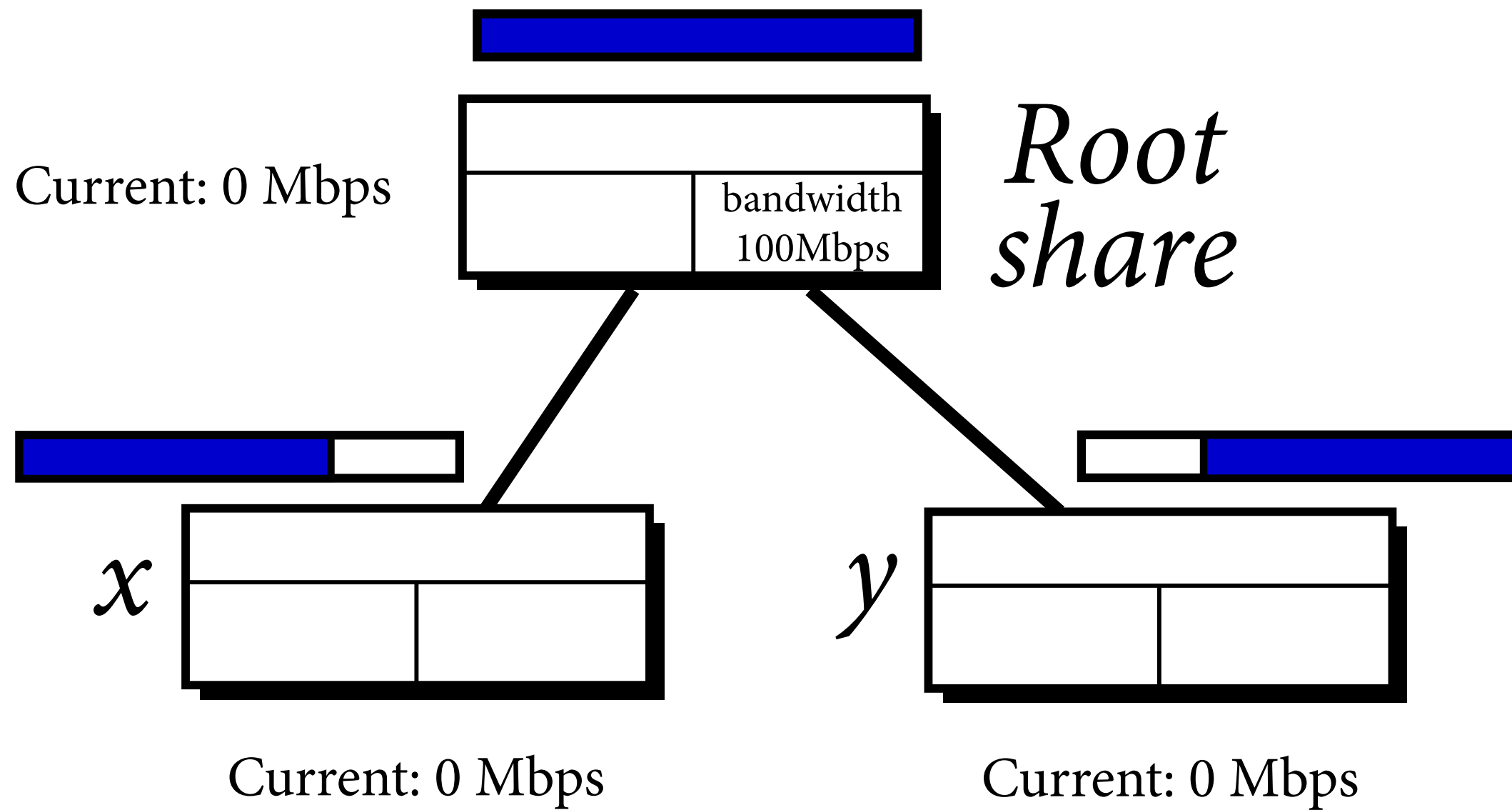
Dynamic Context



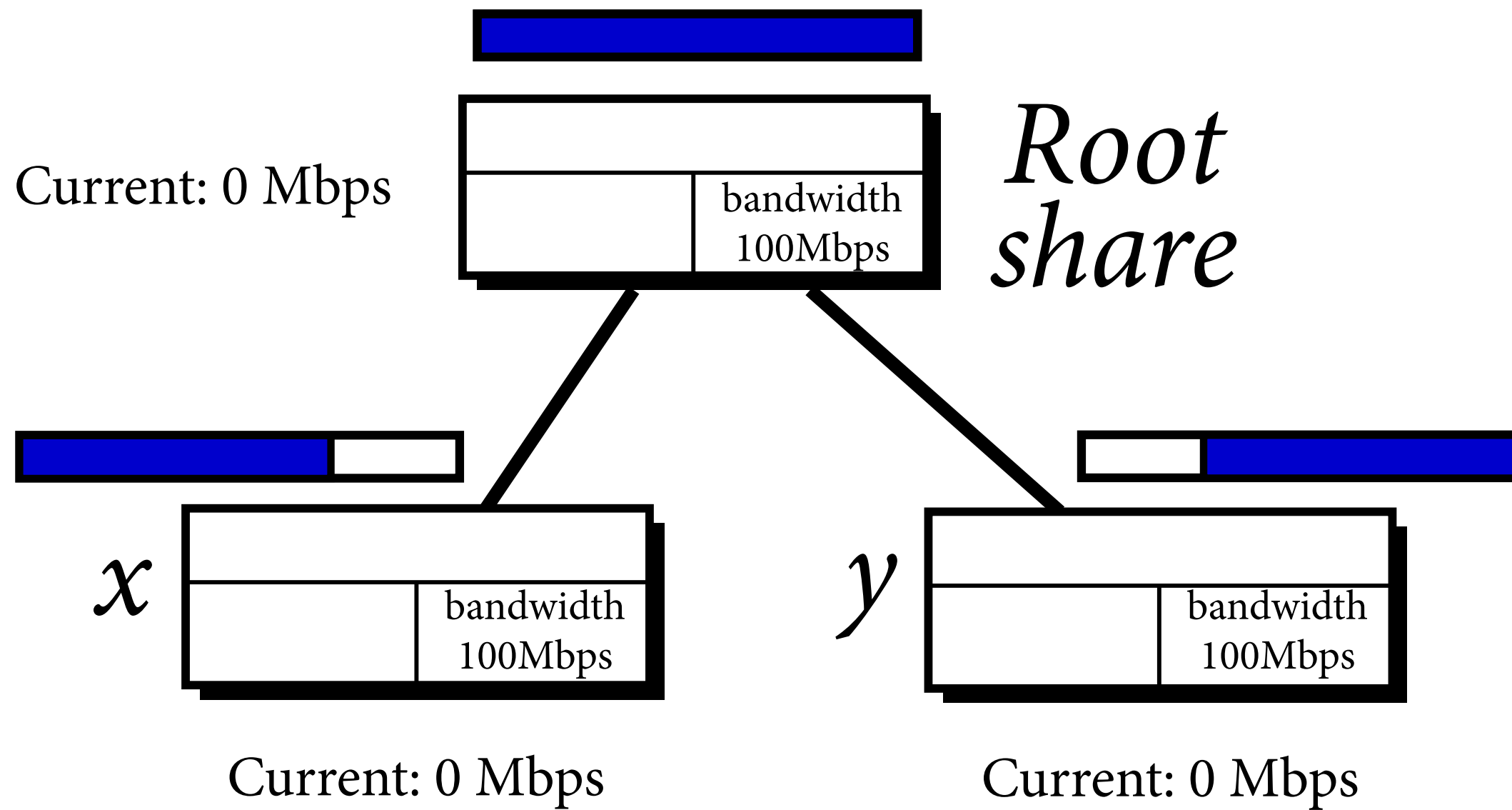
PANE



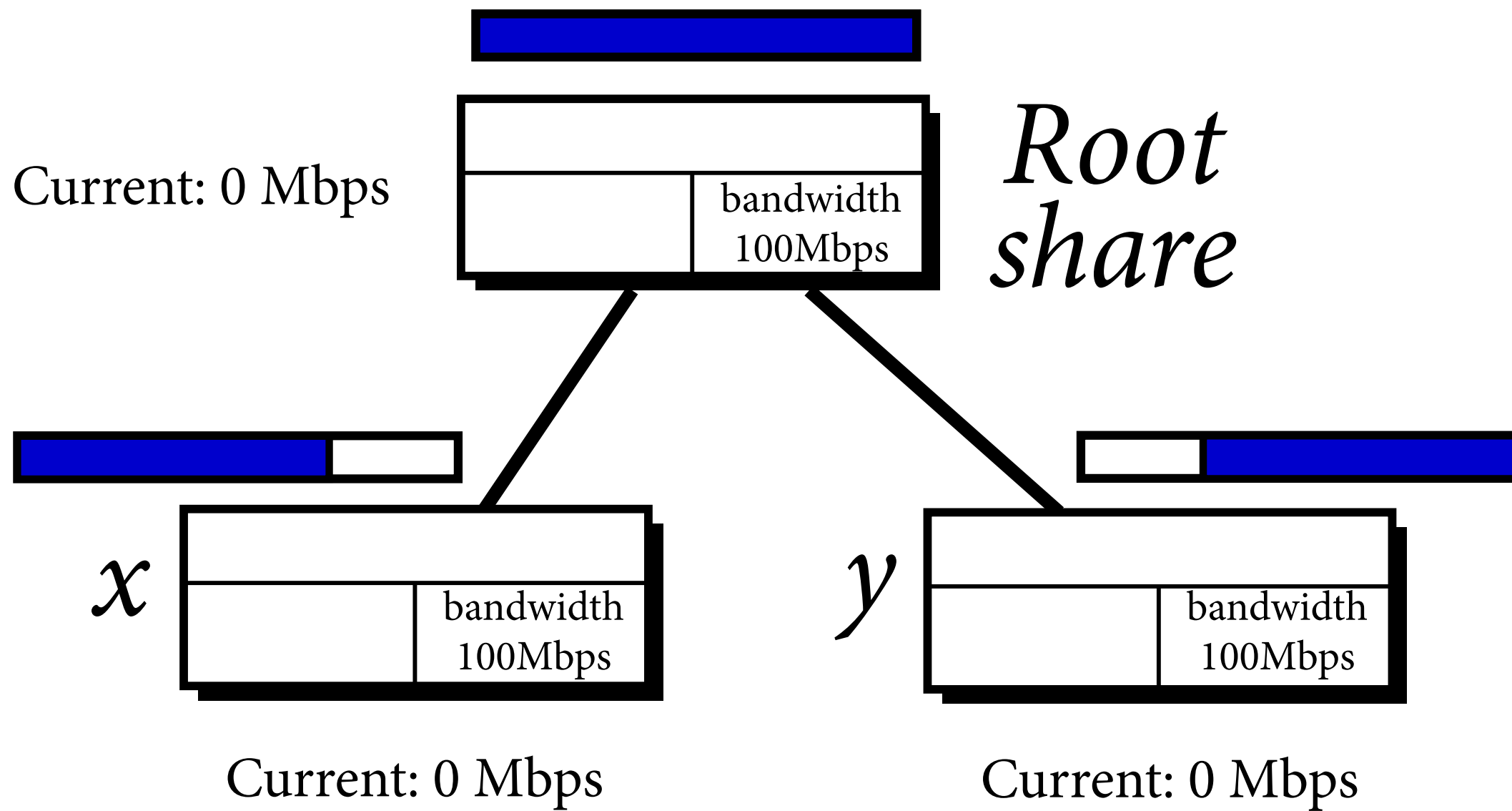
PANE



PANE



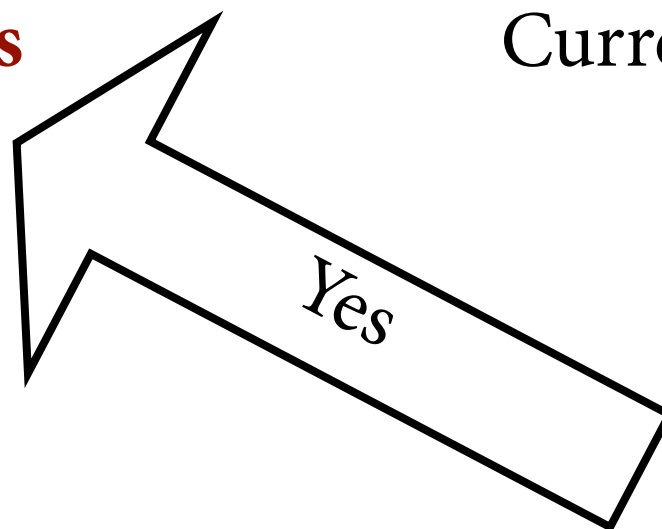
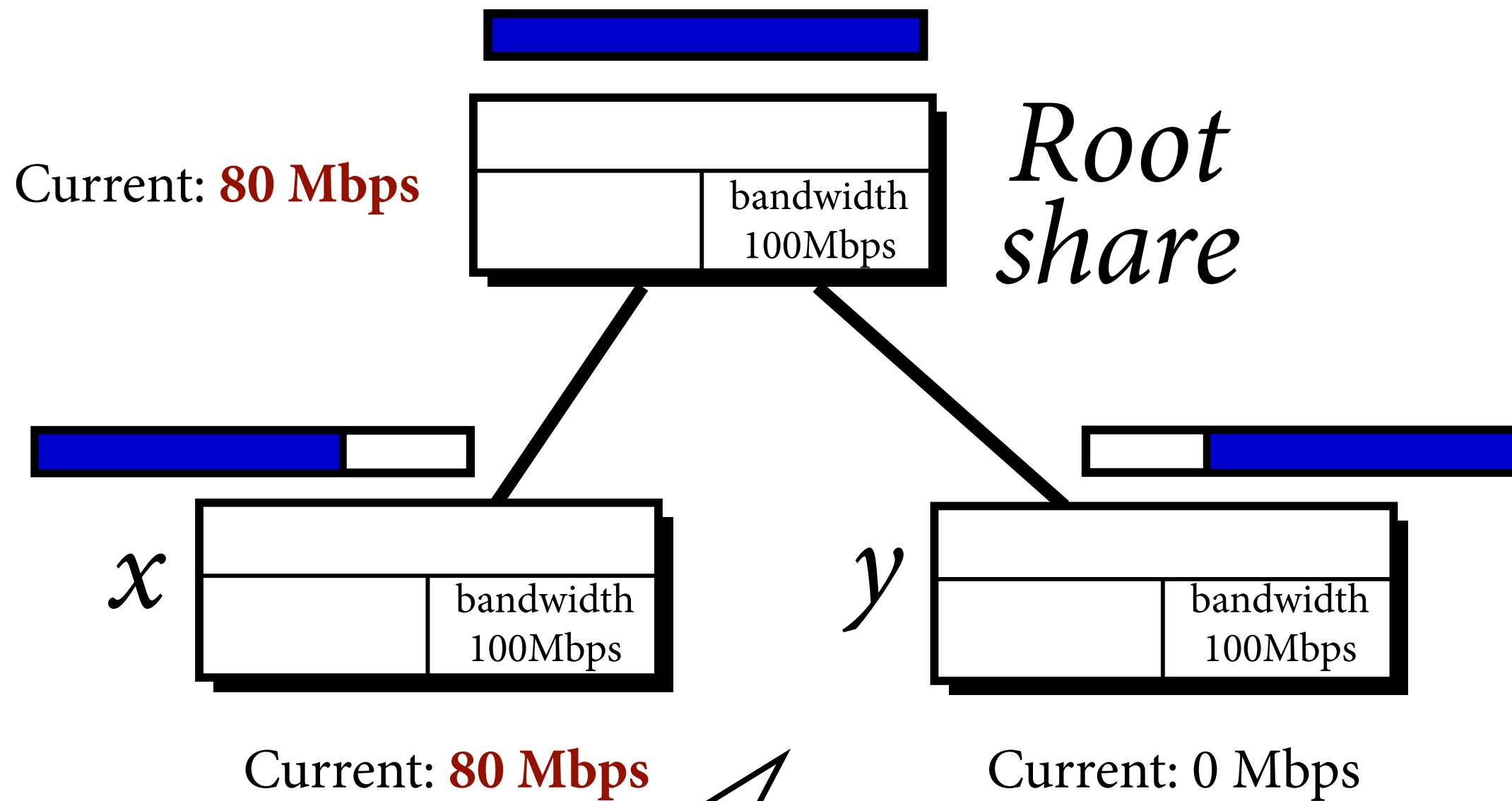
PANE



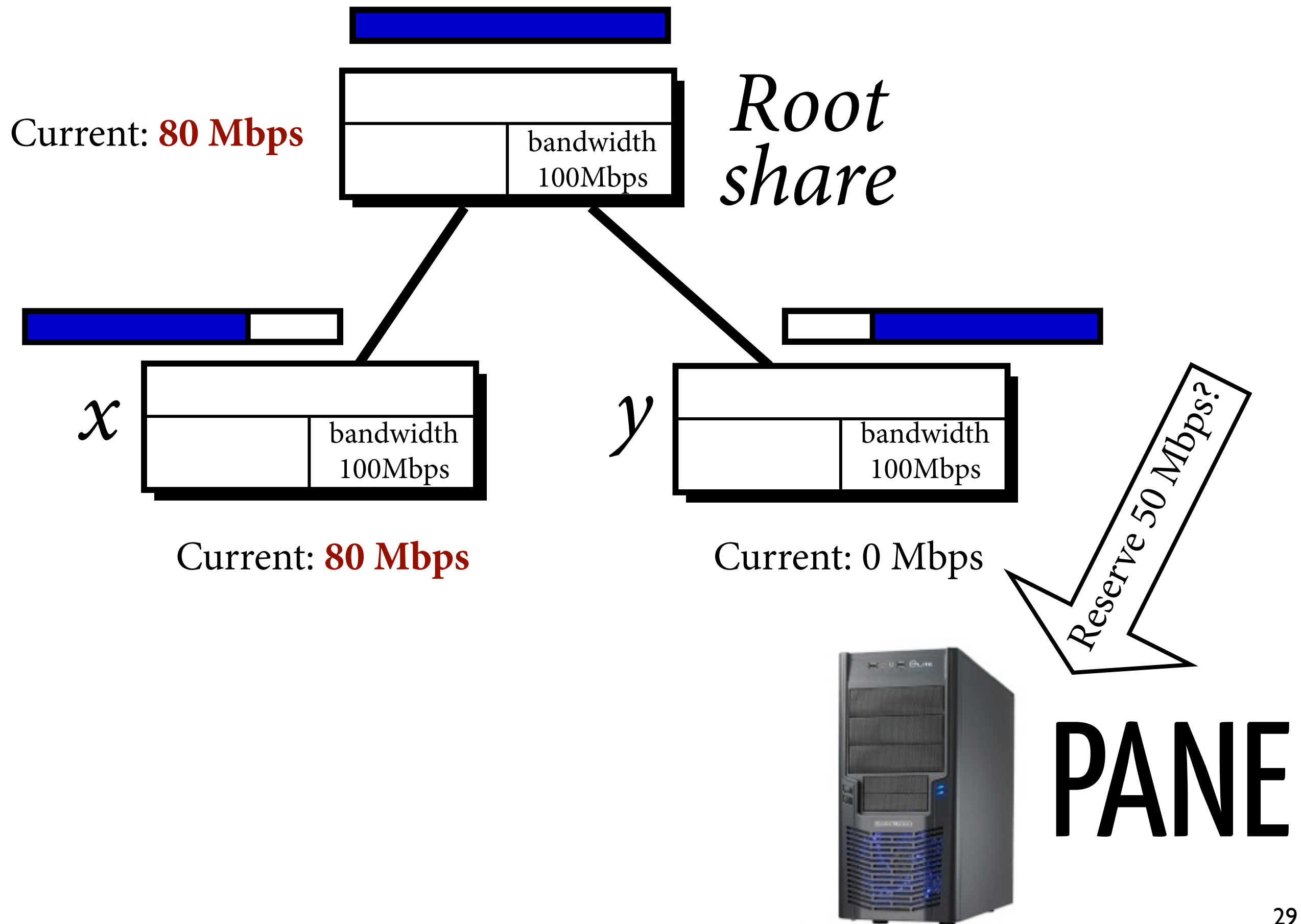
Reserve 80 Mbps?

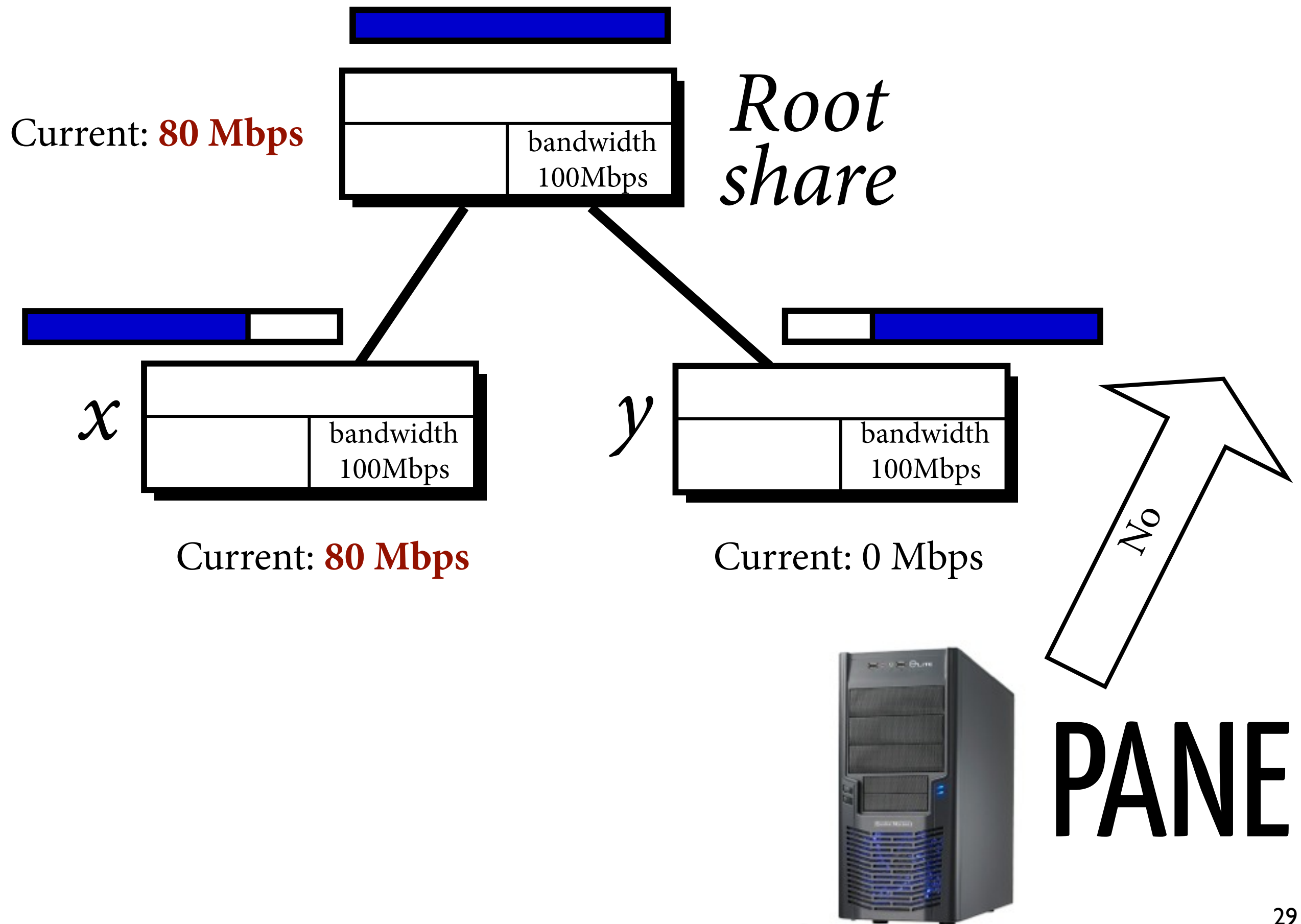


PANE

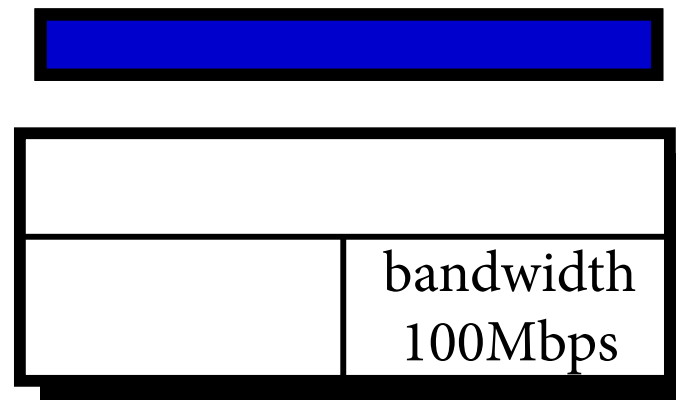


PANE



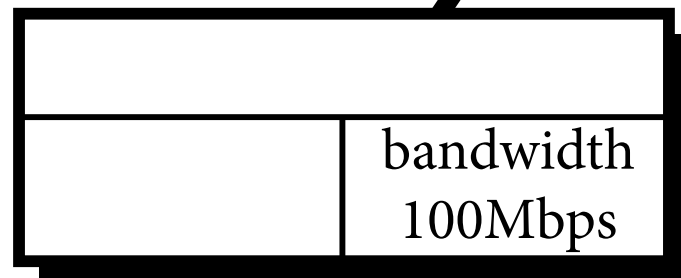


Current: **80 Mbps**



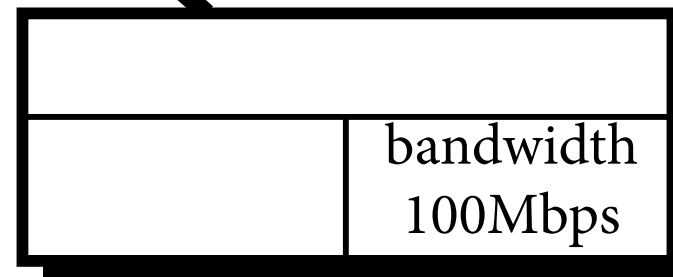
*Root
share*

x

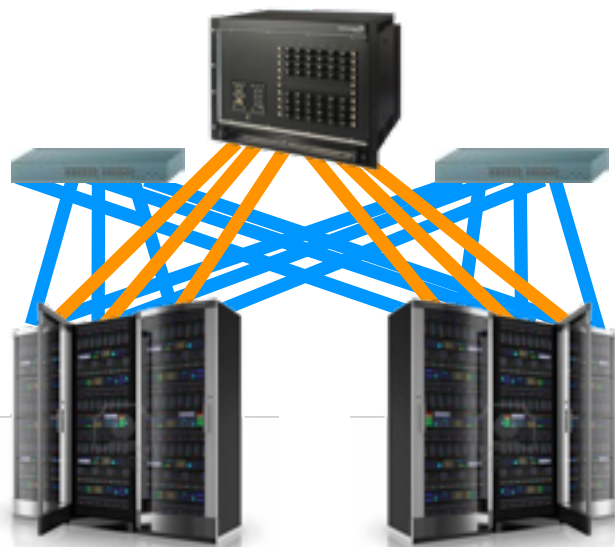


Current: **80 Mbps**

y



Current: 0 Mbps

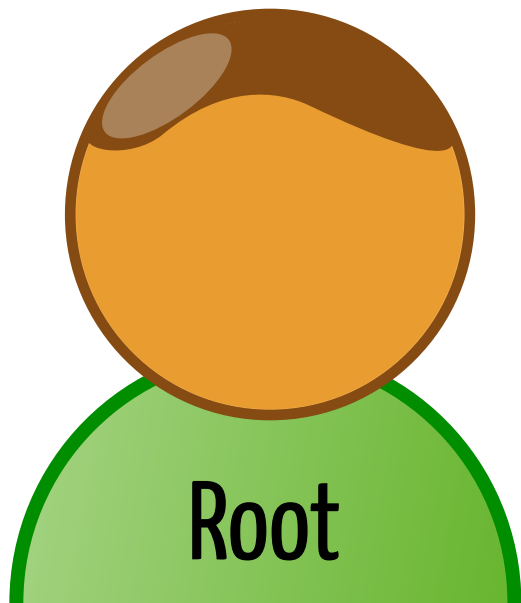


PANE

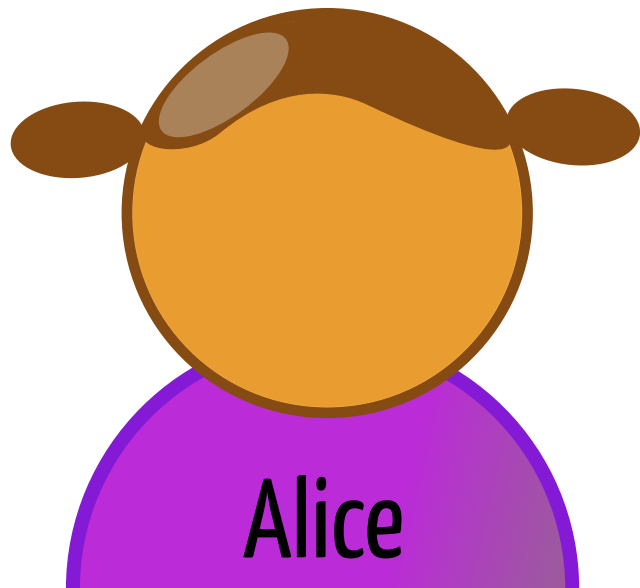
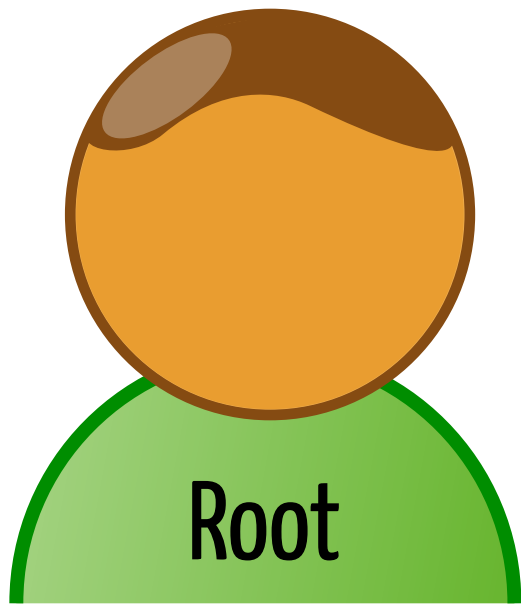
Protocol Sketch



PANE

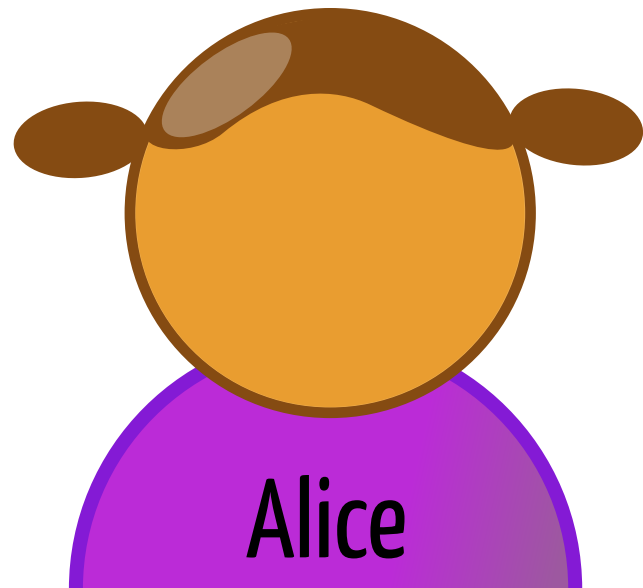
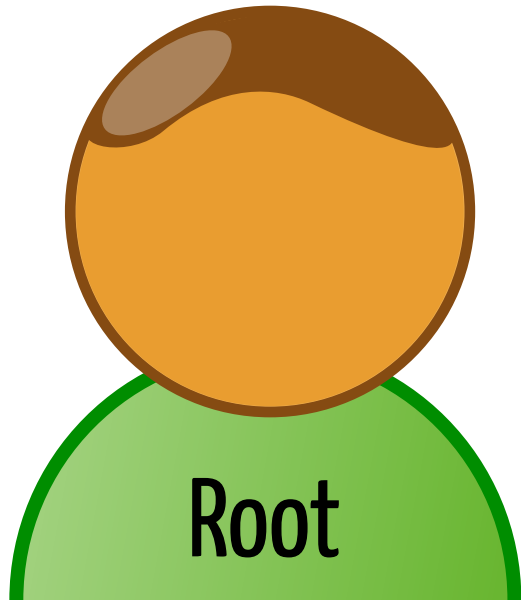


PANE

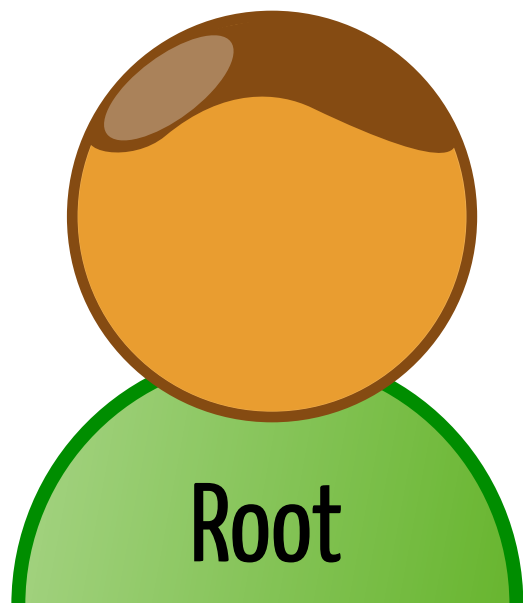


PANE

NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

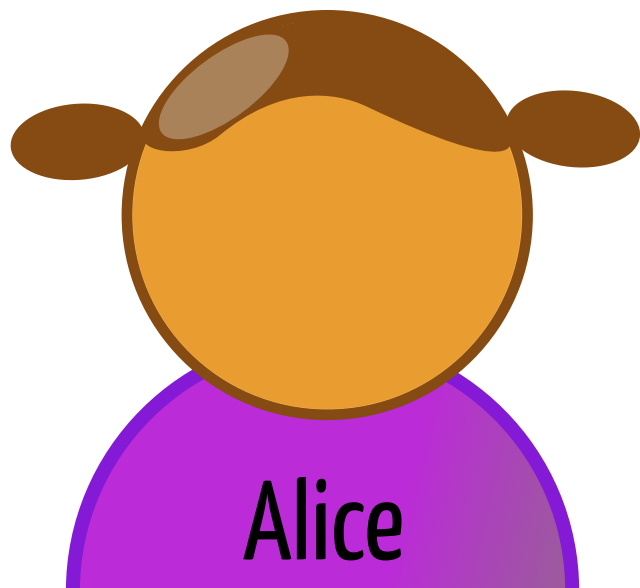


PANE

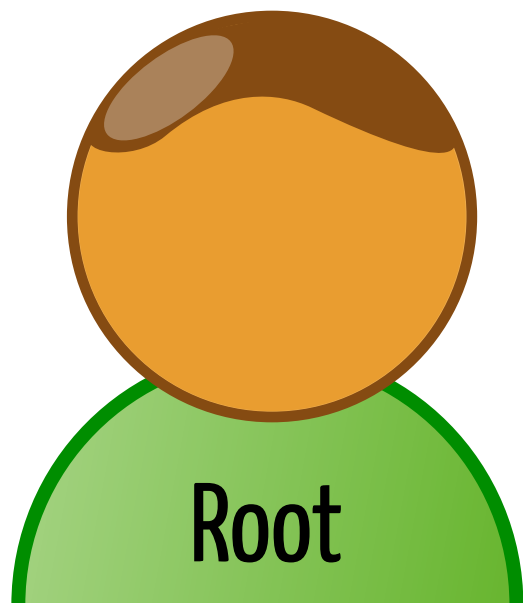


NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

OK



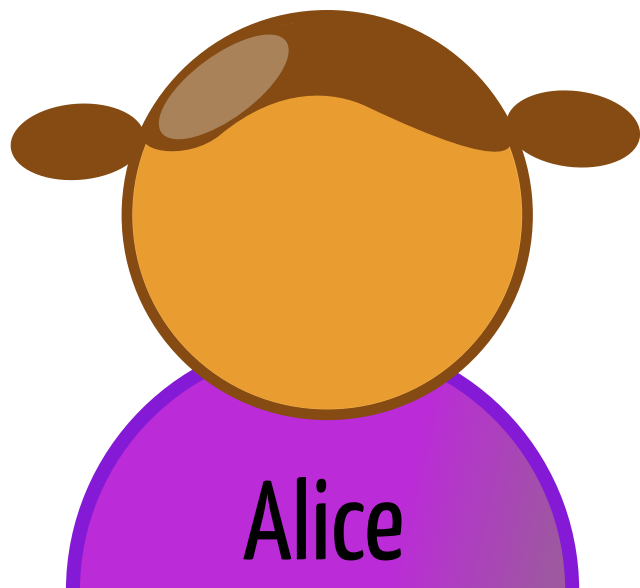
PANE



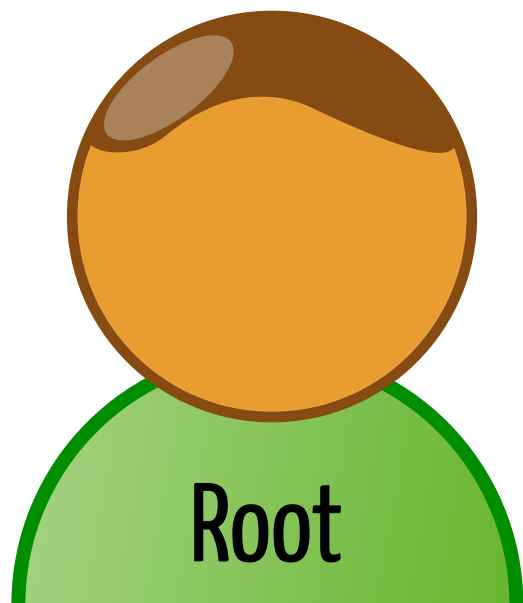
NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

OK

Grant A to Alice.



PANE

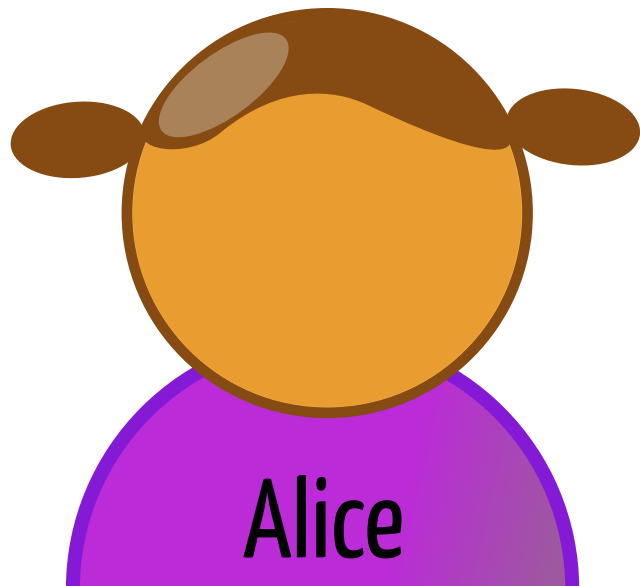


NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

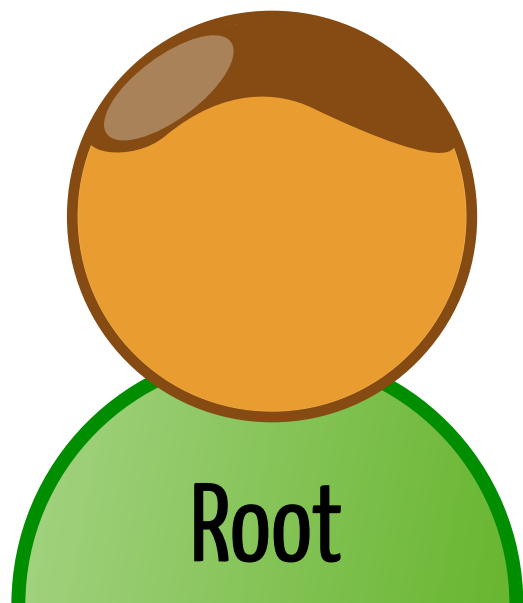
OK

Grant A to Alice.

OK



PANE

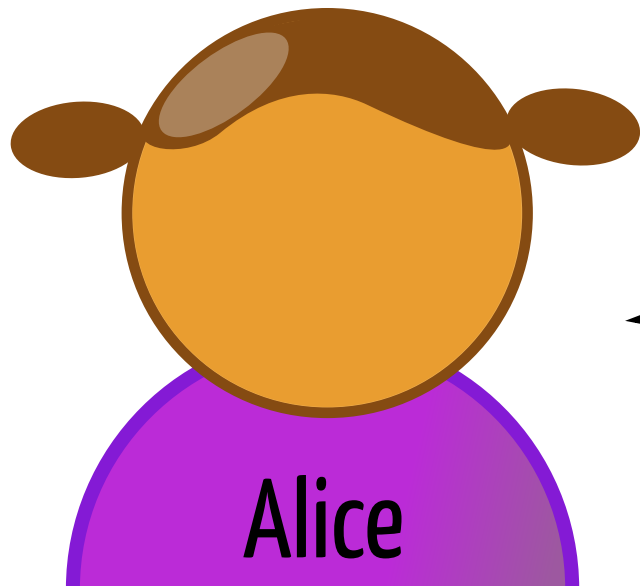


NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

OK

Grant A to Alice.

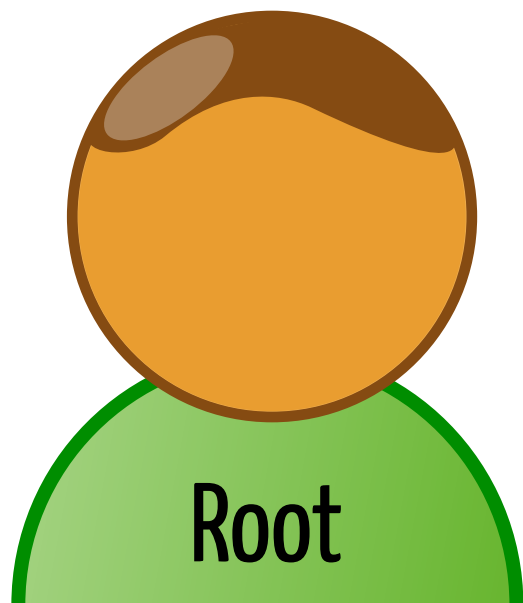
OK



reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.



PANE

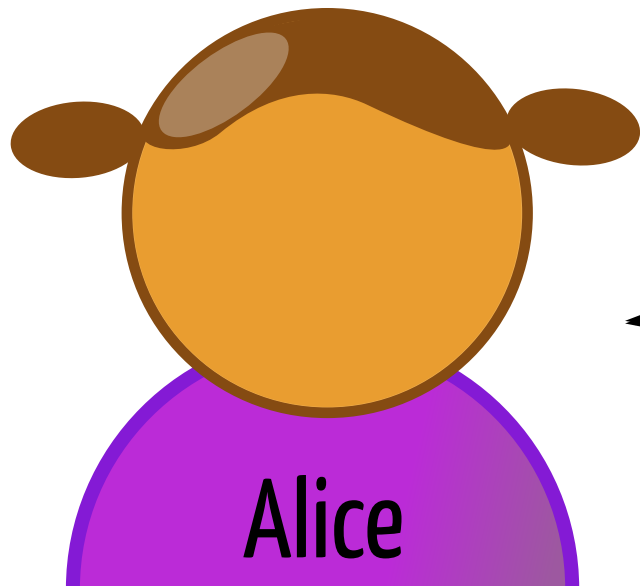


NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

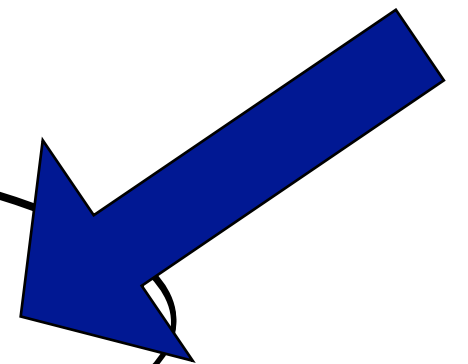
OK

Grant A to Alice.

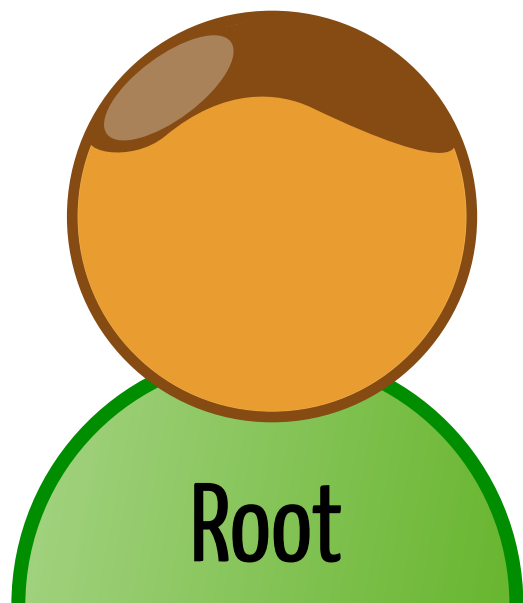
OK



reserve(user=Alice,
dstPort=80) = 5Mb on **A**
from now to +10min.



PANE



NewShare A for
(user=Alice) [reserve <= 10Mb]
on rootShare.

OK

Grant A to Alice.

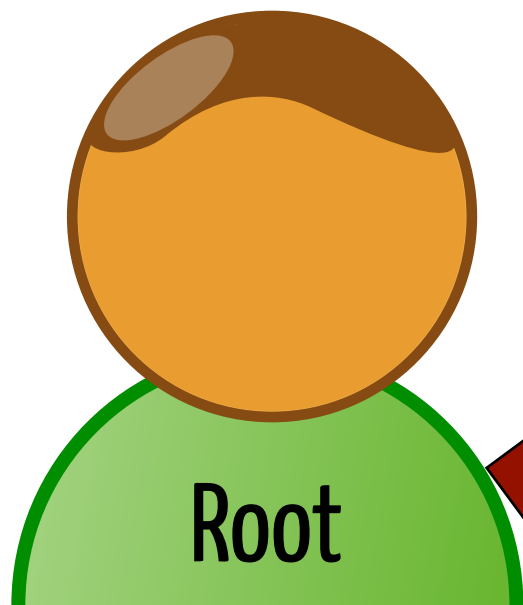
OK



reserve(**user=Alice**,
dstPort=80) = 5Mb on A
from now to +10min.



PANE



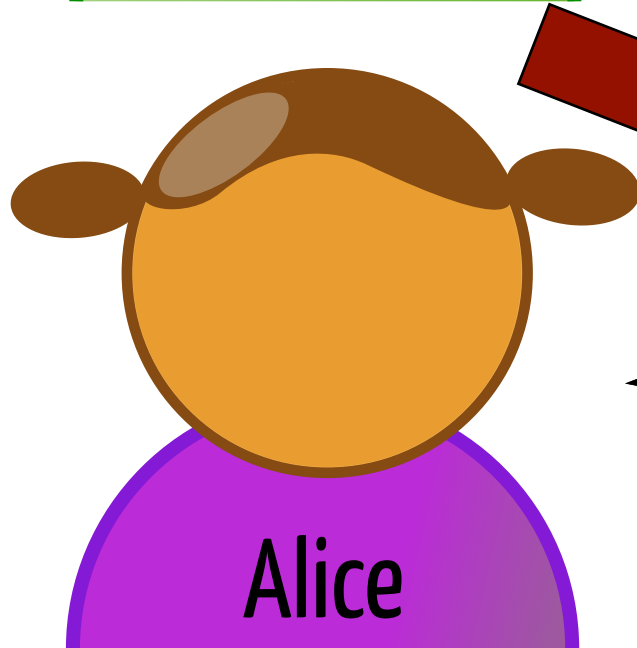
Root

NewShare A for
(**user=Alice**) [reserve <= 10Mb]
on rootShare.

OK

Grant A to Alice.

OK



Alice

reserve(**user=Alice**,
dstPort=80) = 5Mb on A
from now to +10min.

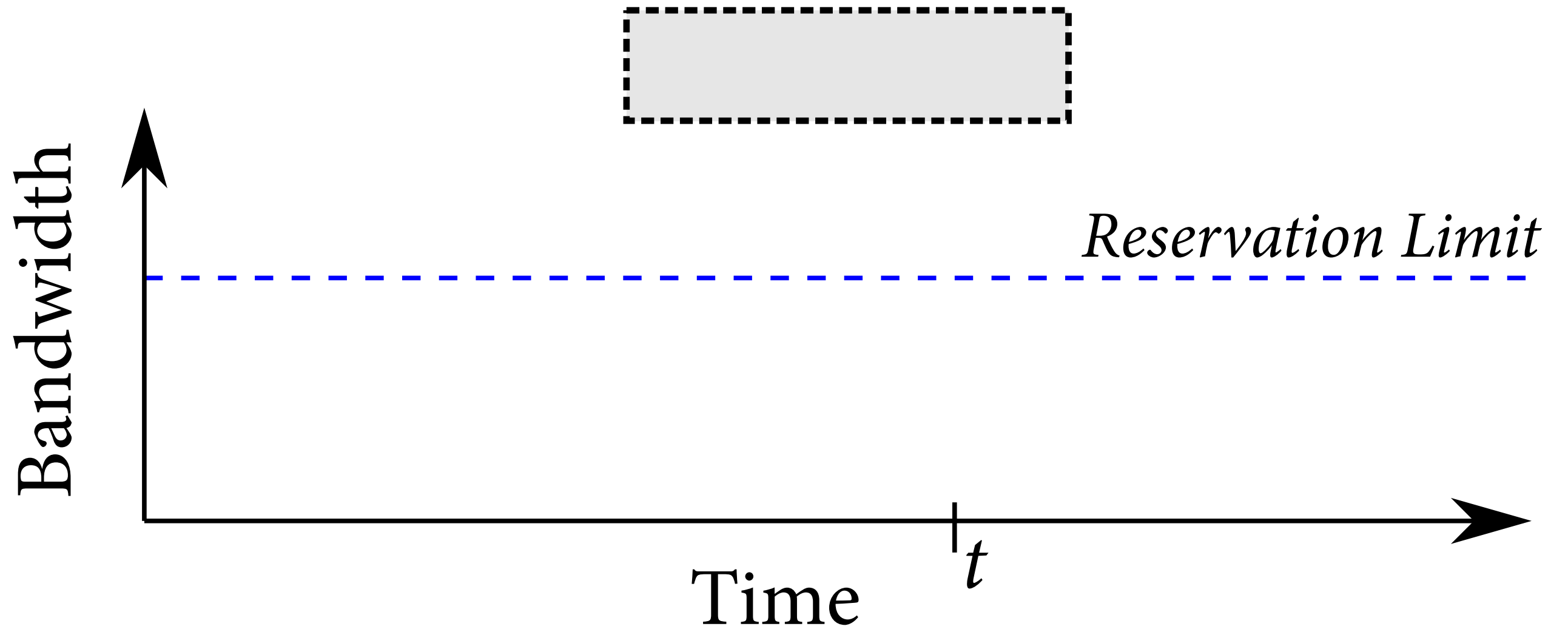


PANE

`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`



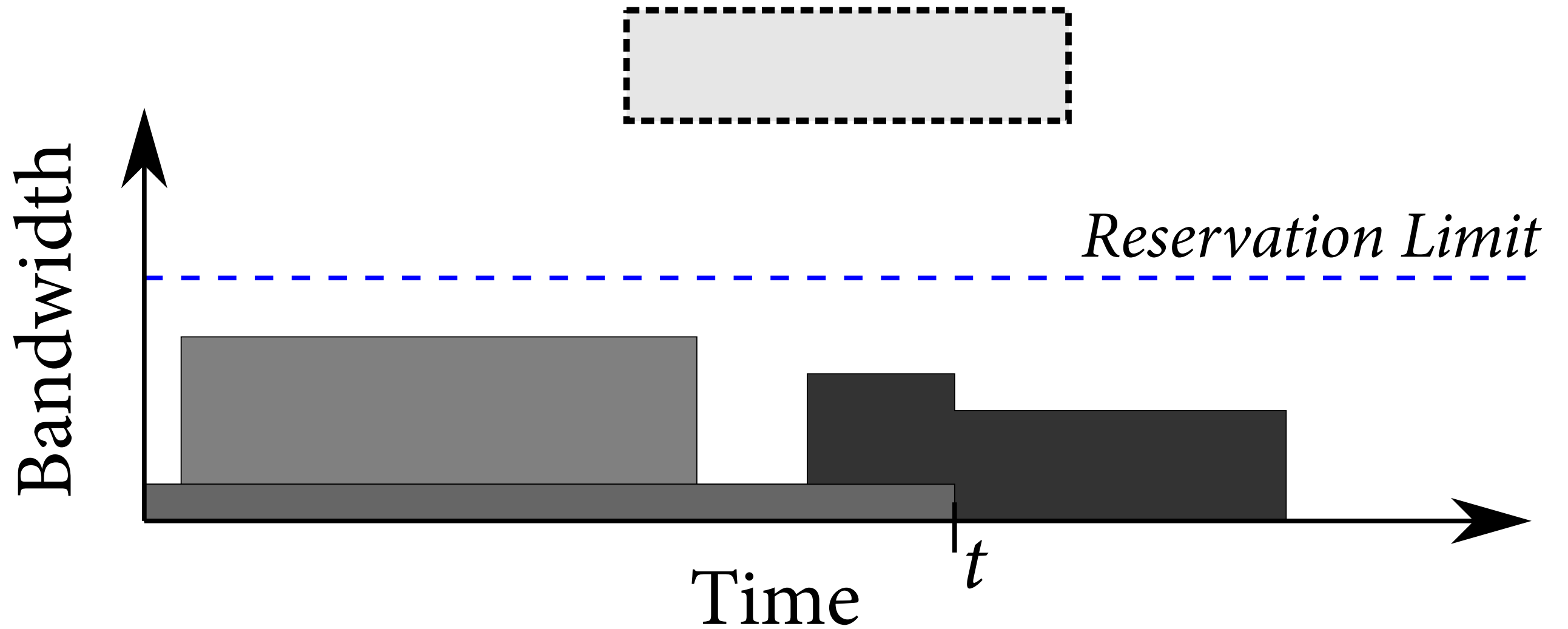
PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`



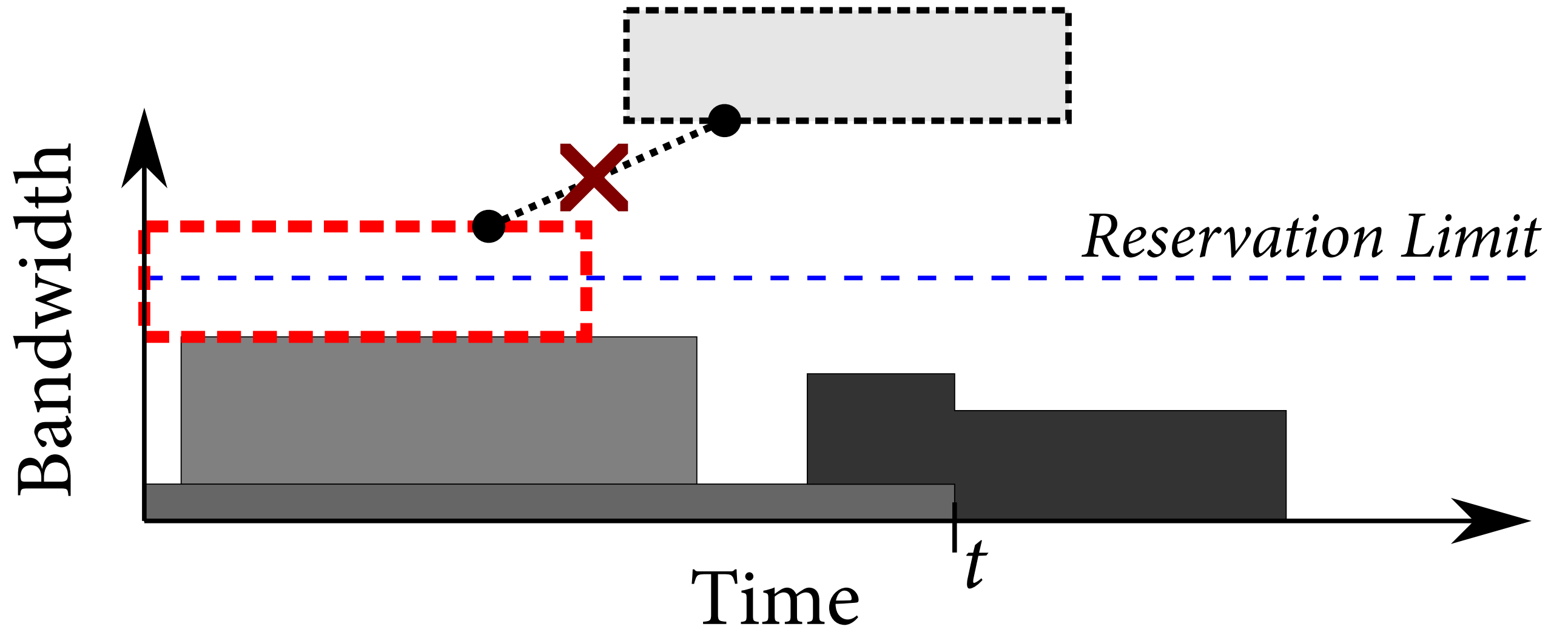
PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`



PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`



PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`

NO



PANE

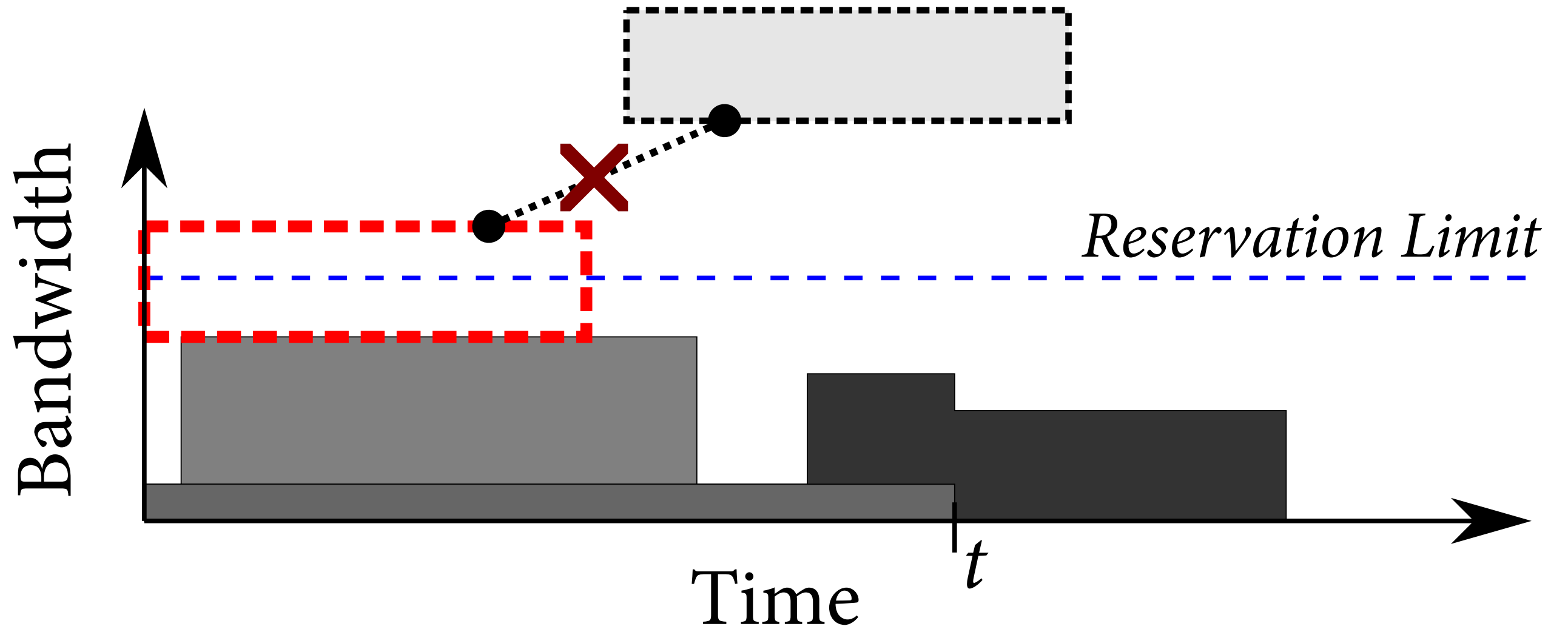


`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`

NO



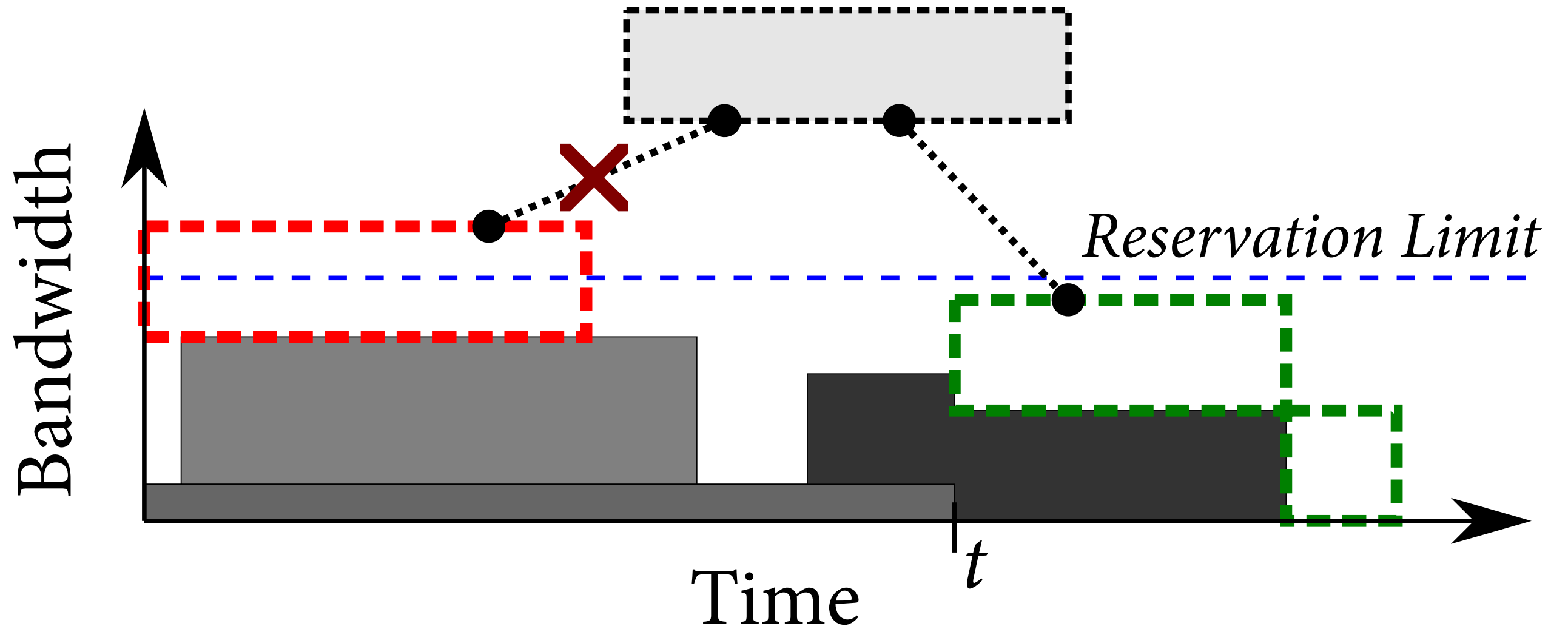
PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from +20min to +30min.`



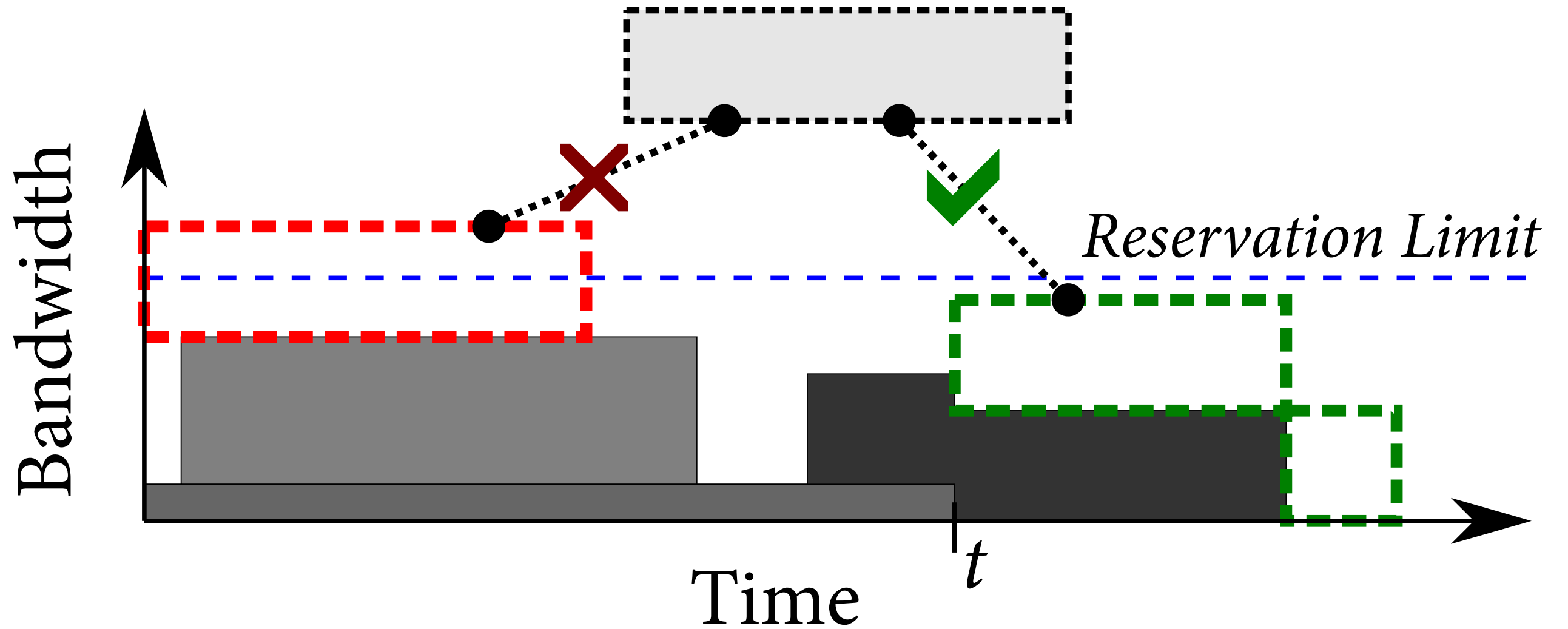
PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from +20min to +30min.`



PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from +20min to +30min.`



PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`

NO

`reserve(user=Alice,
dstPort=80) = 5Mb on A
from +20min to +30min.`

OK



PANE



`reserve(user=Alice,
dstPort=80) = 5Mb on A
from now to +10min.`

NO

`reserve(user=Alice,
dstPort=80) = 5Mb on A
from +20min to +30min.`

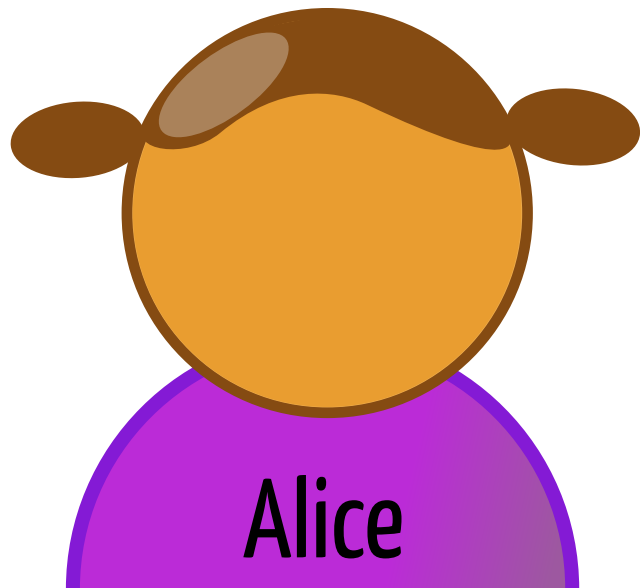
OK



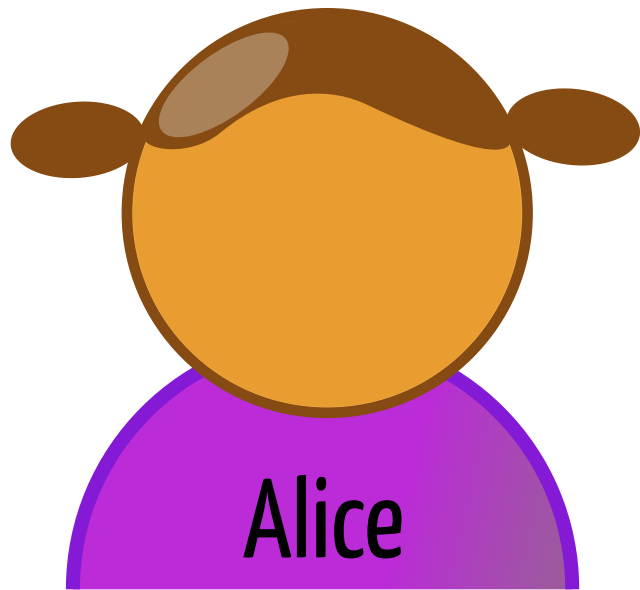
PANE



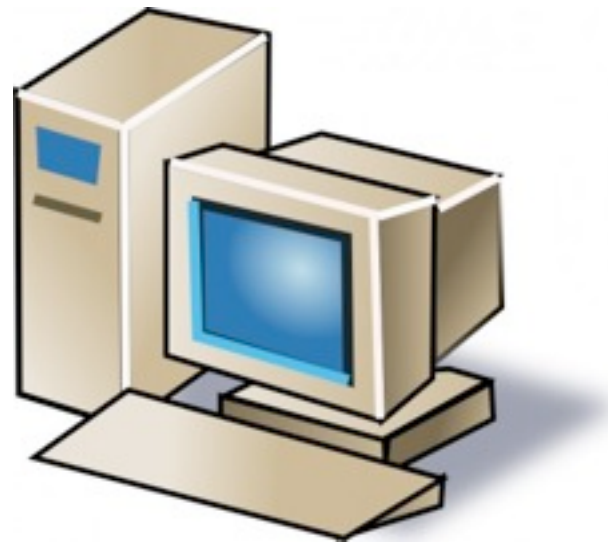
PANE



PANE



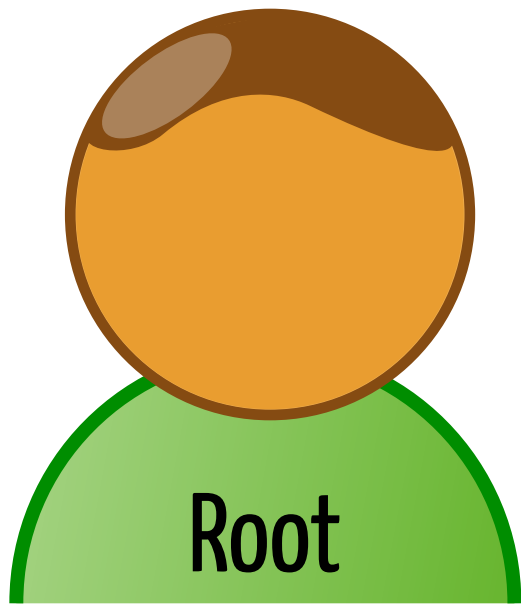
Alice



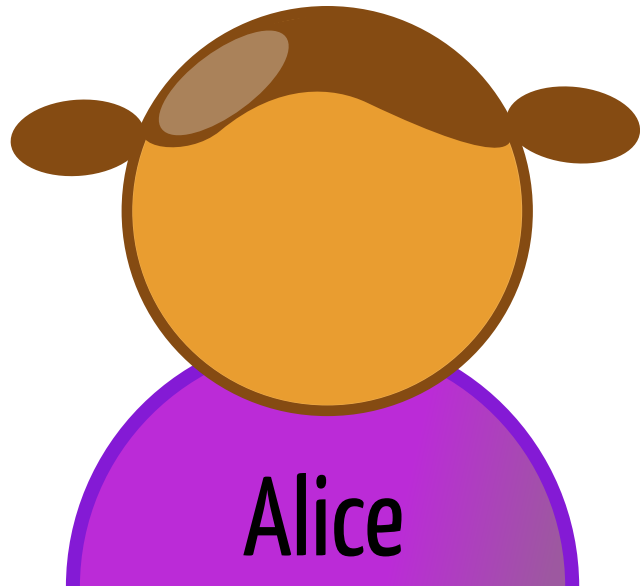
10.0.0.2



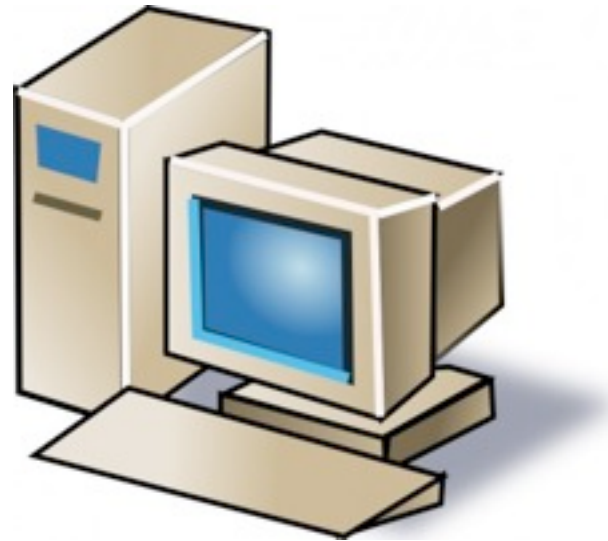
PANE



Root



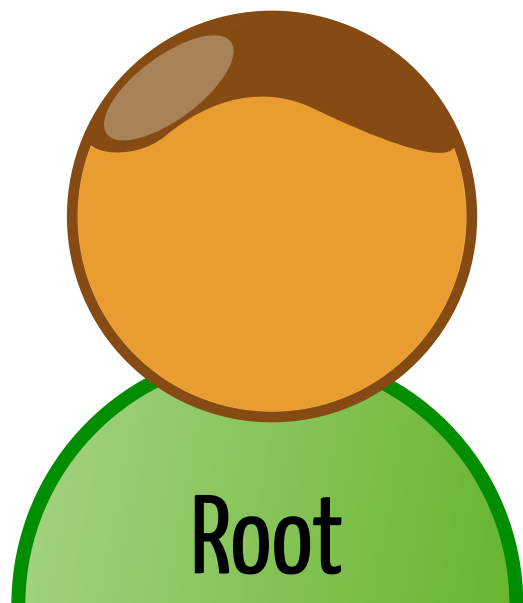
Alice



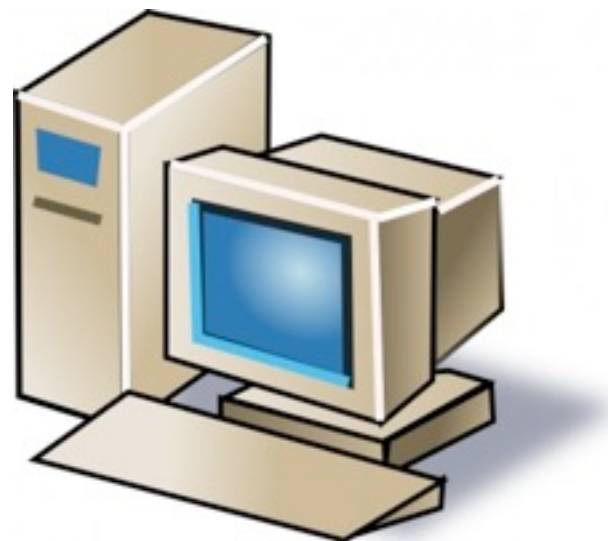
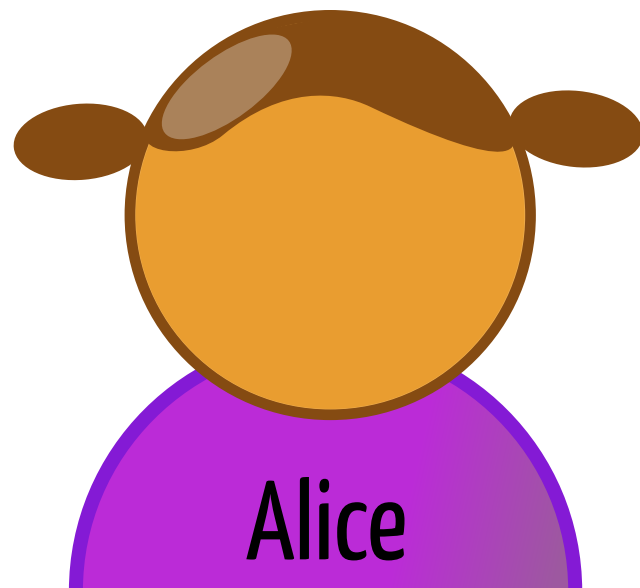
10.0.0.2



PANE



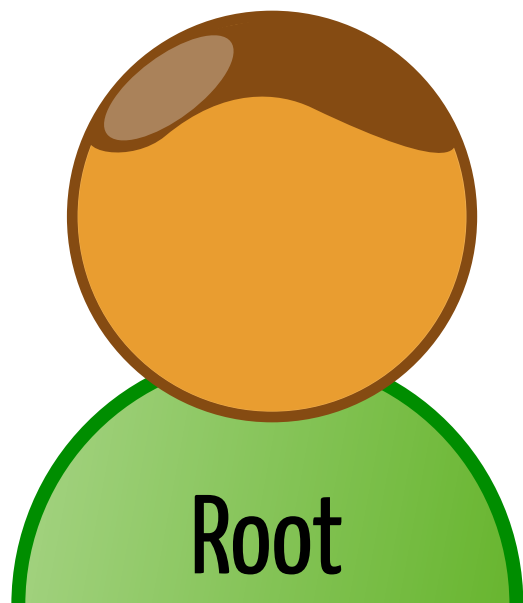
NewShare aAC for
(dstHost=10.0.0.2) [deny = True]
on rootShare.



10.0.0.2

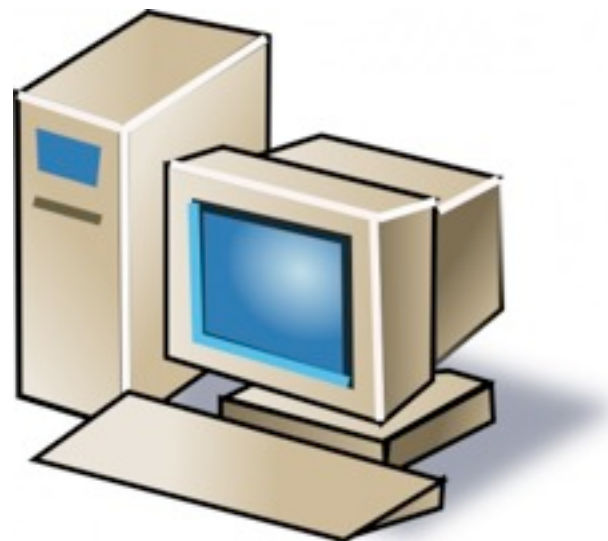
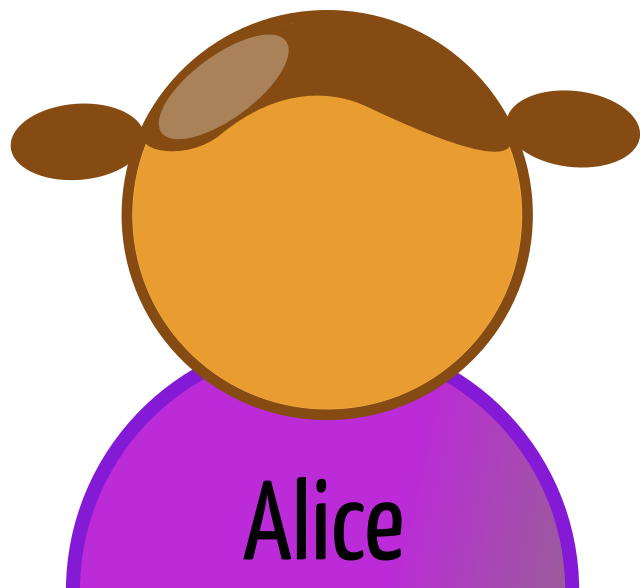


PANE



NewShare aAC for
(dstHost=10.0.0.2) [deny = True]
on rootShare.

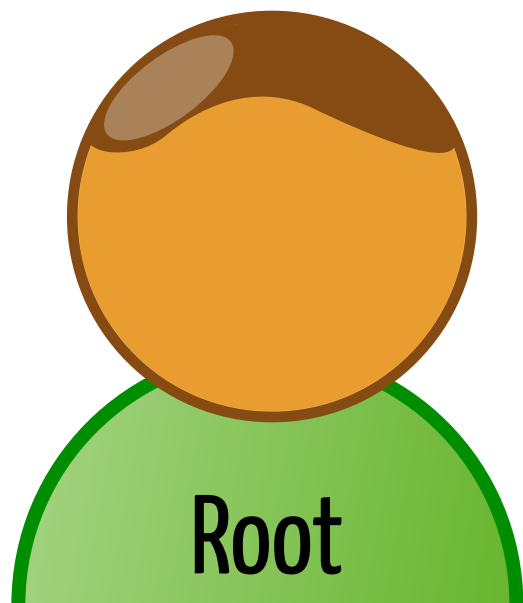
OK



10.0.0.2



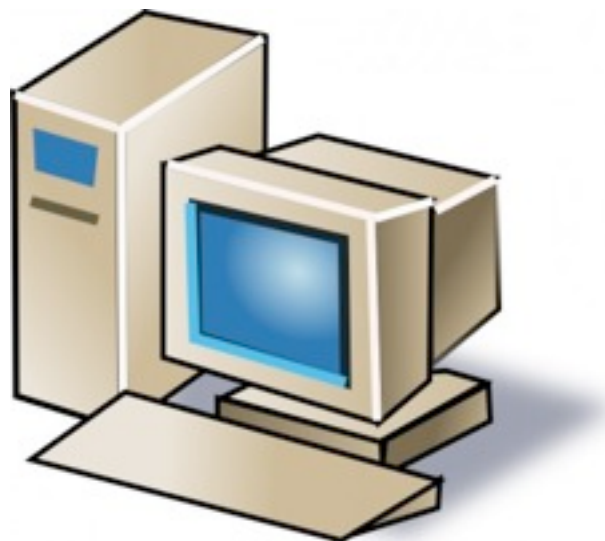
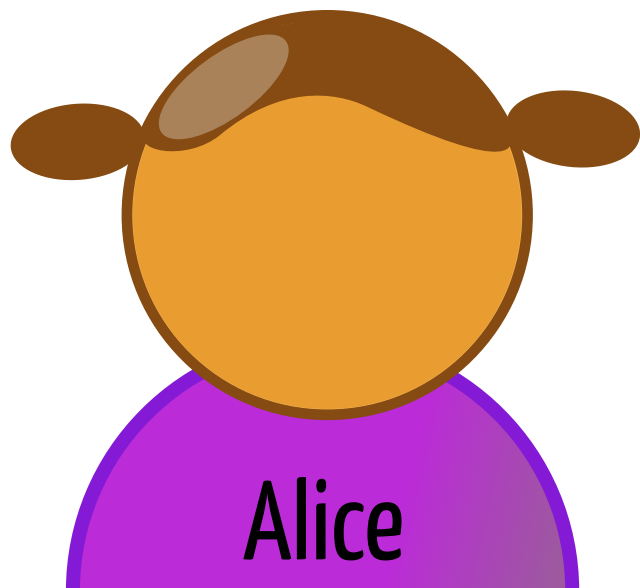
PANE



NewShare aAC for
(dstHost=10.0.0.2) [deny = True]
on rootShare.

OK

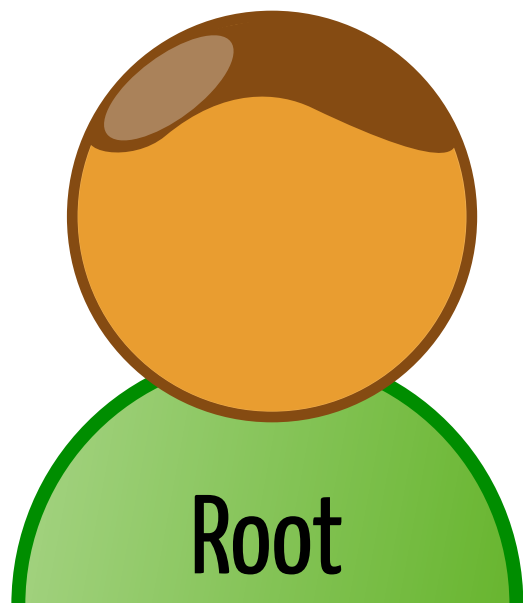
Grant aAC to Alice.



10.0.0.2



PANE

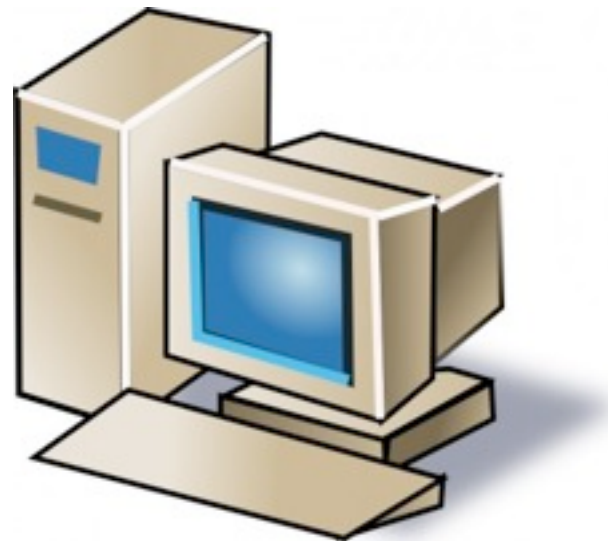
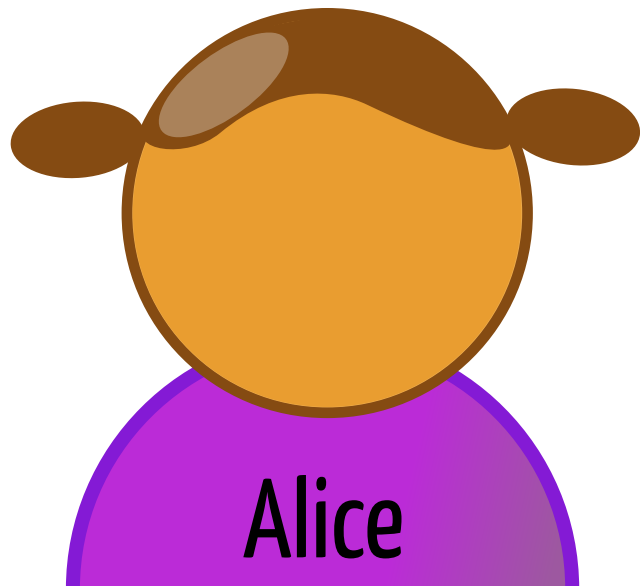


NewShare aAC for
(dstHost=10.0.0.2) [deny = True]
on rootShare.

OK

Grant aAC to Alice.

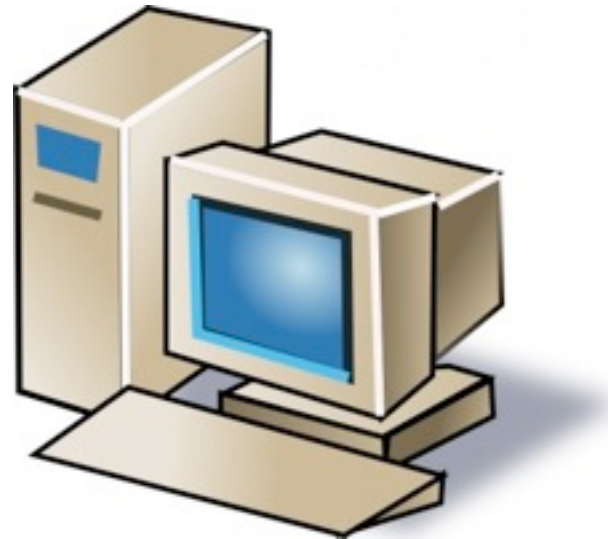
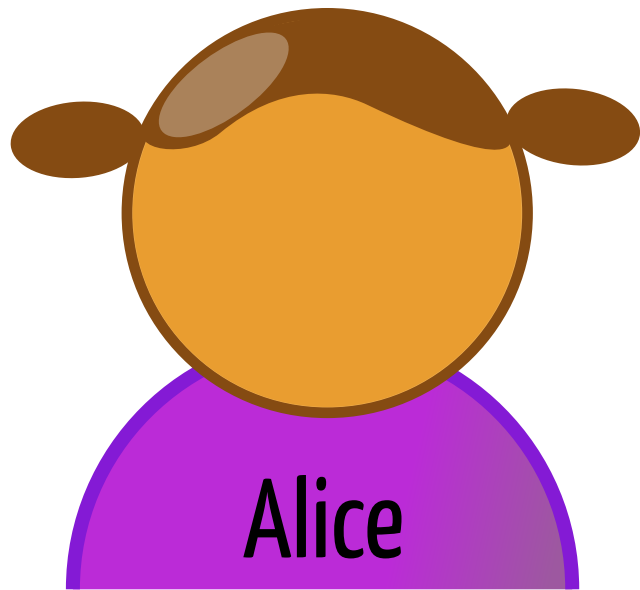
OK



10.0.0.2



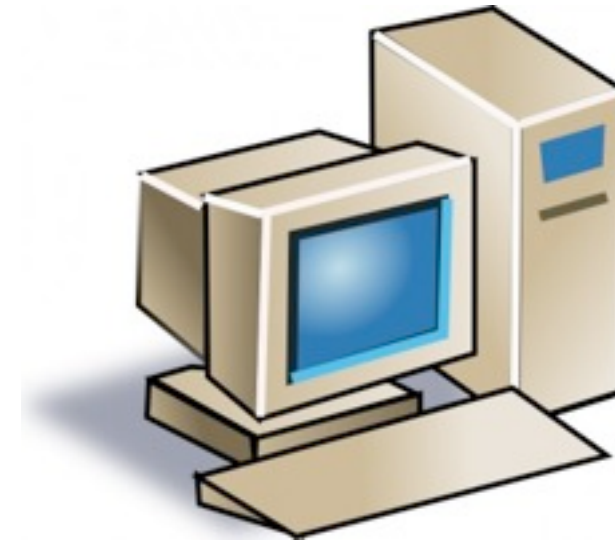
PANE



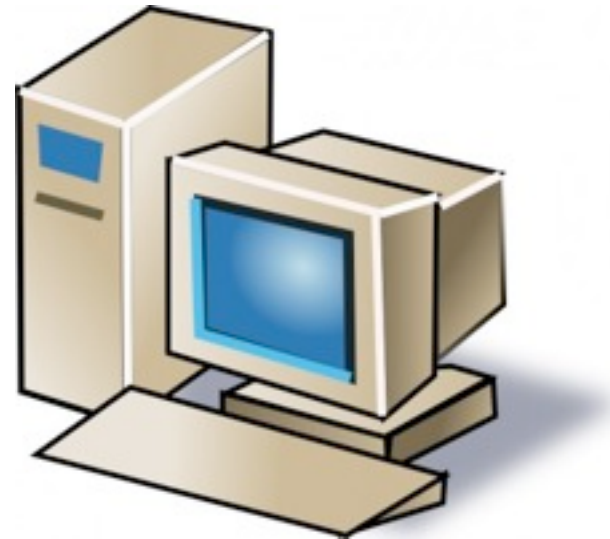
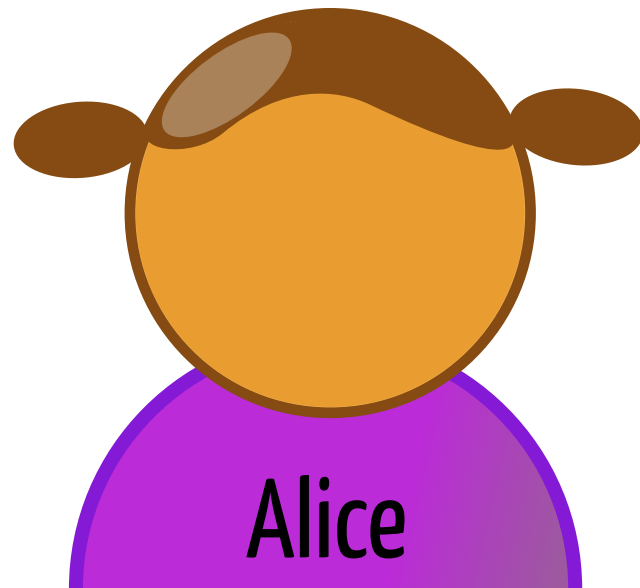
10.0.0.2



PANE



10.0.0.3

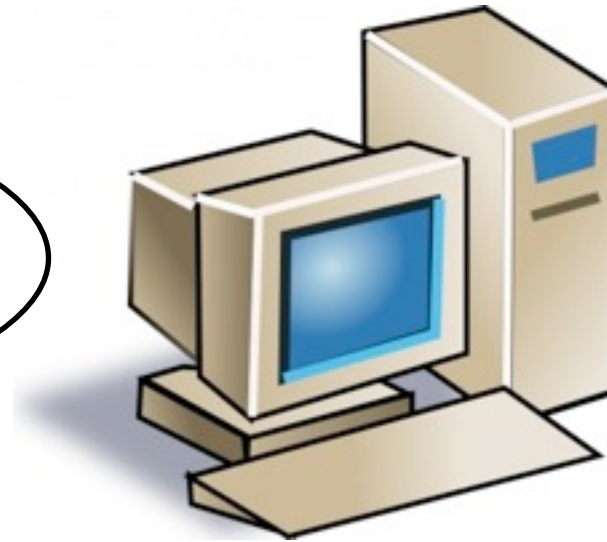


10.0.0.2

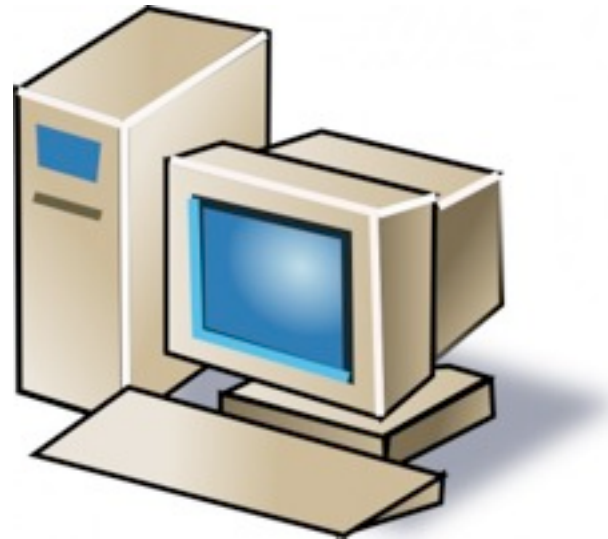
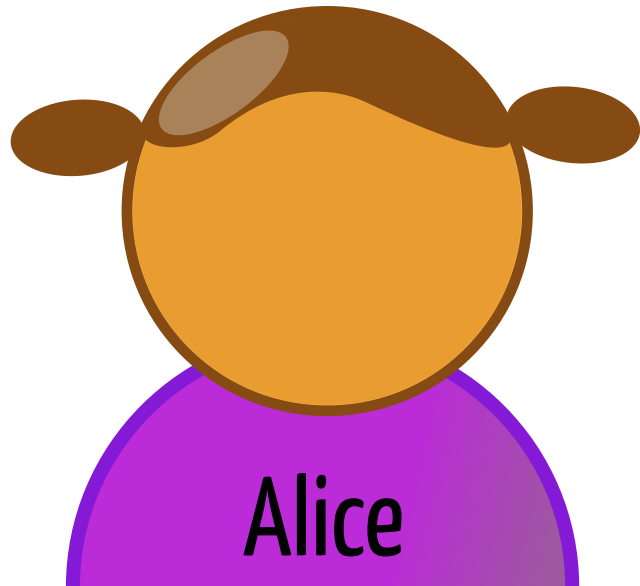


PANE

`deny(dstHost=10.0.0.2,
srcHost=10.0.0.3) on aAC
from now to +5min.`



10.0.0.3



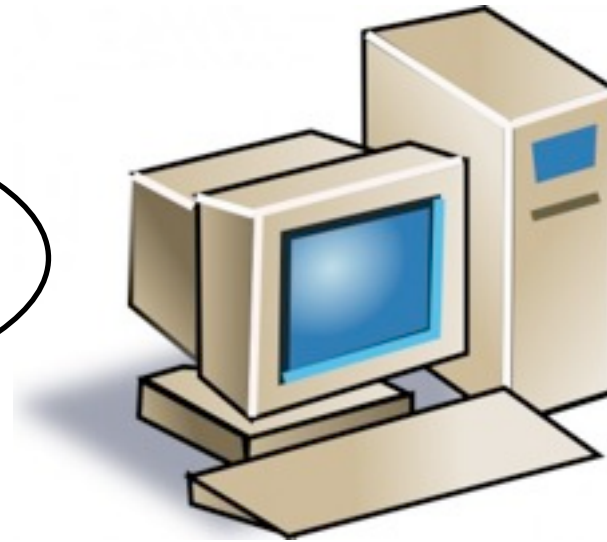
10.0.0.2



PANE

`deny(dstHost=10.0.0.2,
srcHost=10.0.0.3) on aAC
from now to +5min.`

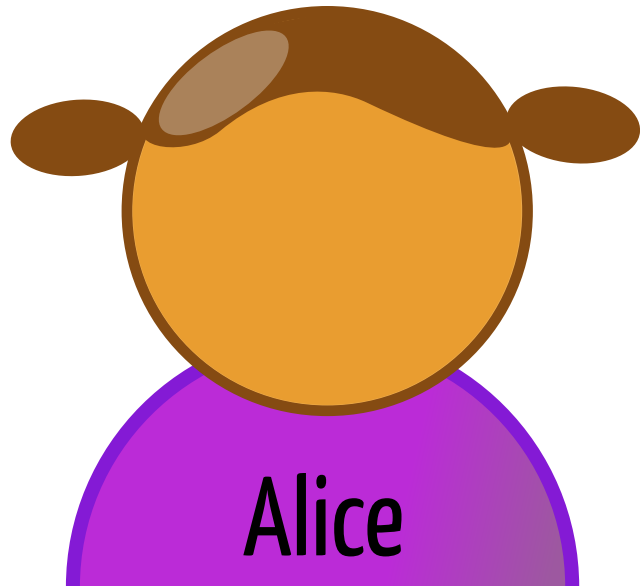
OK



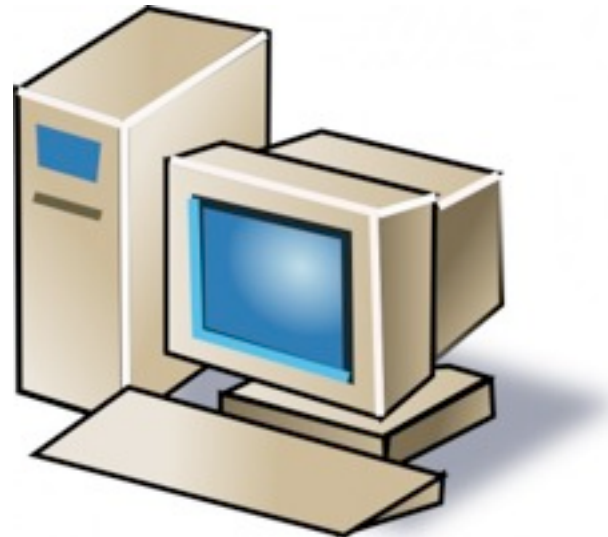
10.0.0.3



Eve



Alice



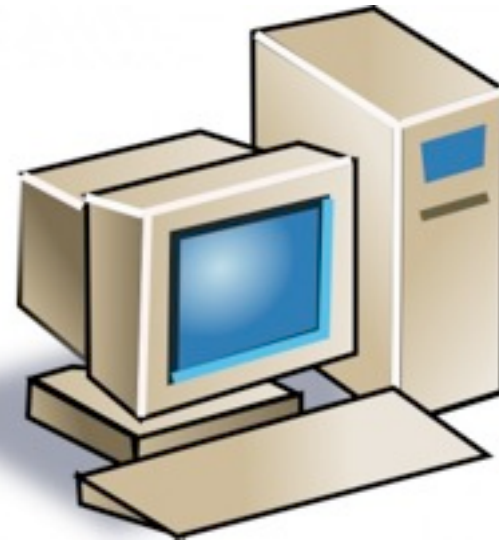
10.0.0.2



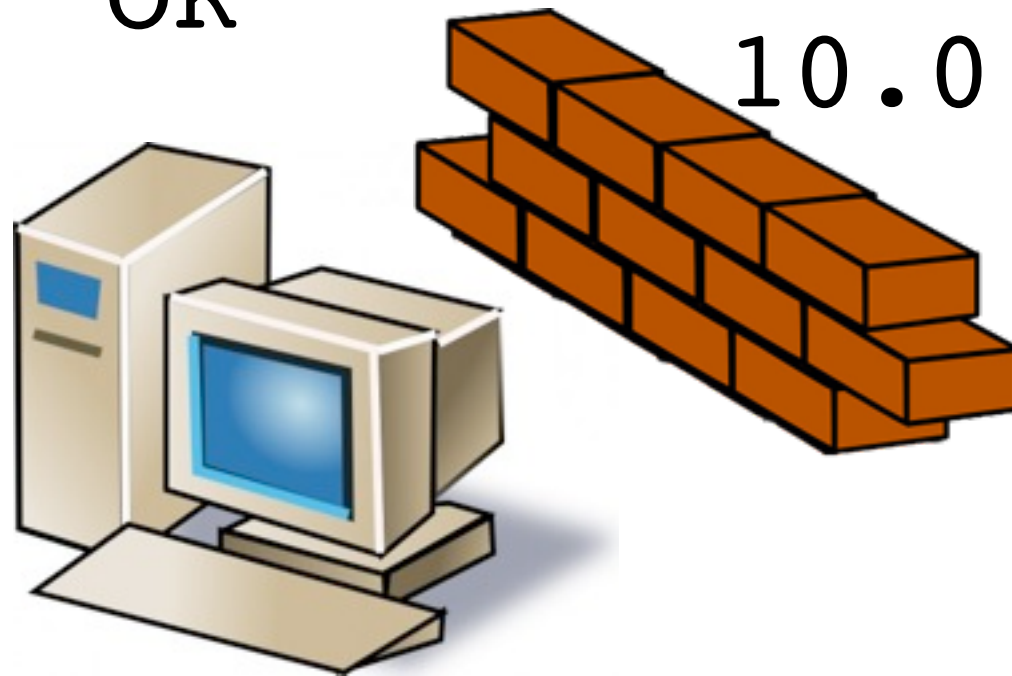
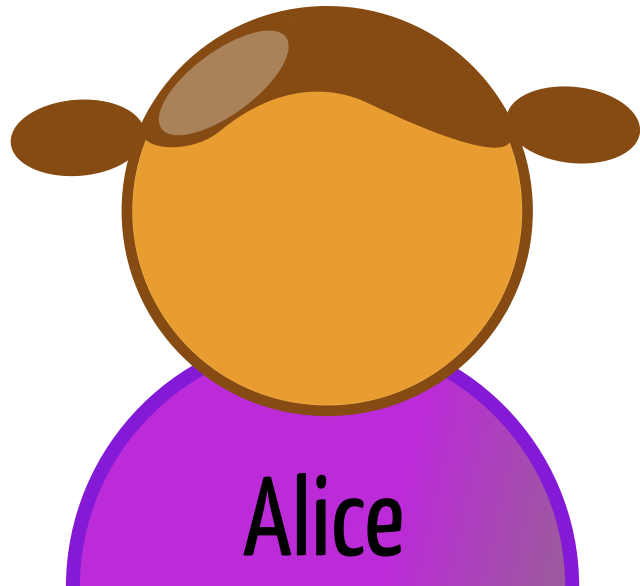
PANE

`deny(dstHost=10.0.0.2,
srcHost=10.0.0.3) on aAC
from now to +5min.`

OK



10.0.0.3

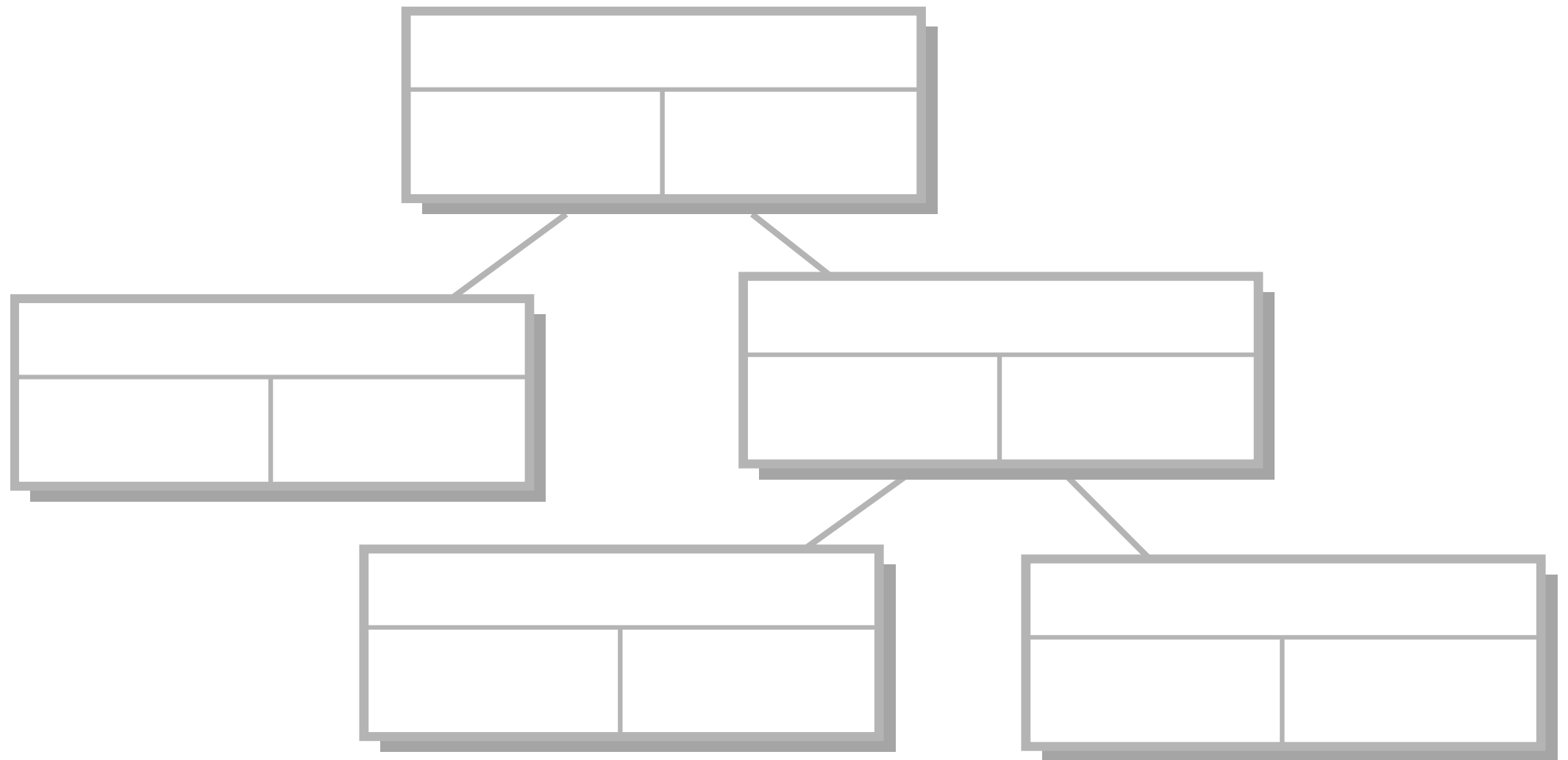


10.0.0.2

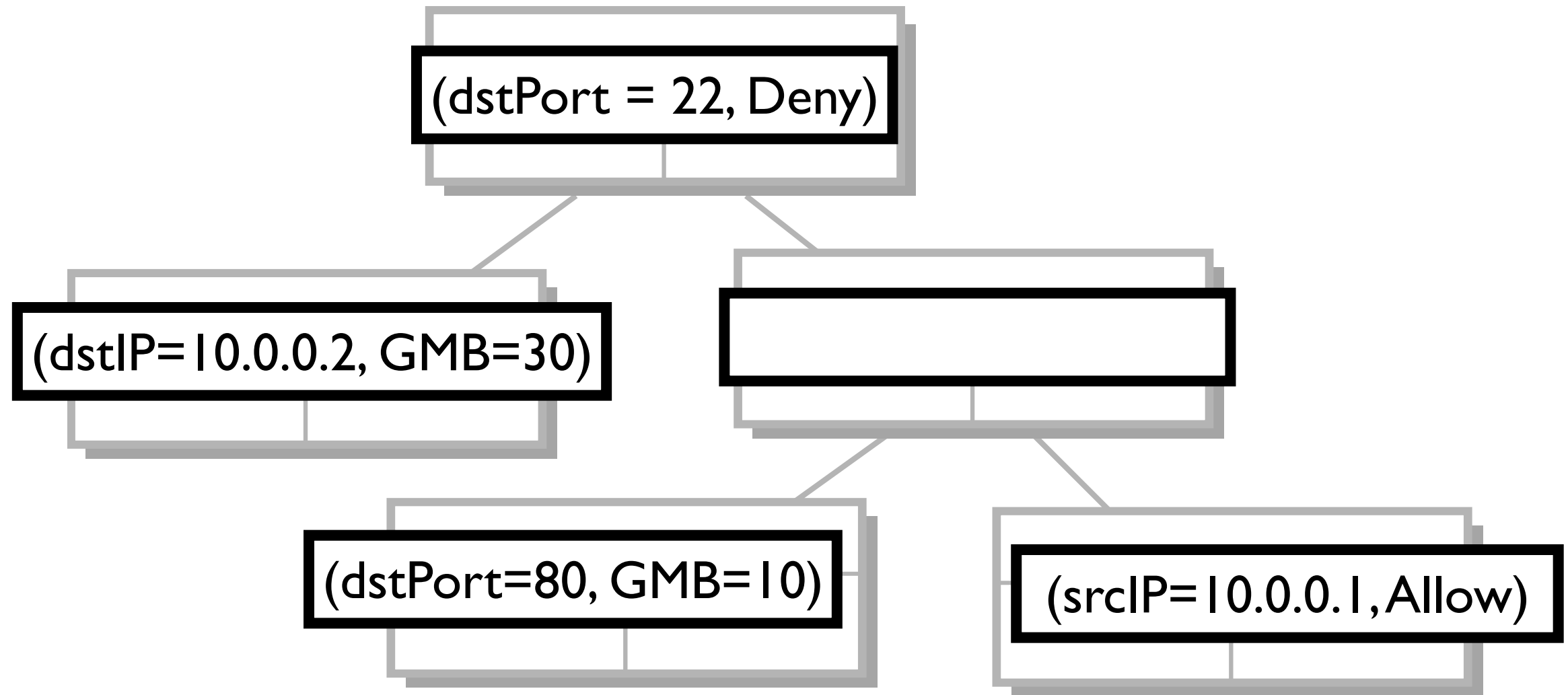


PANE

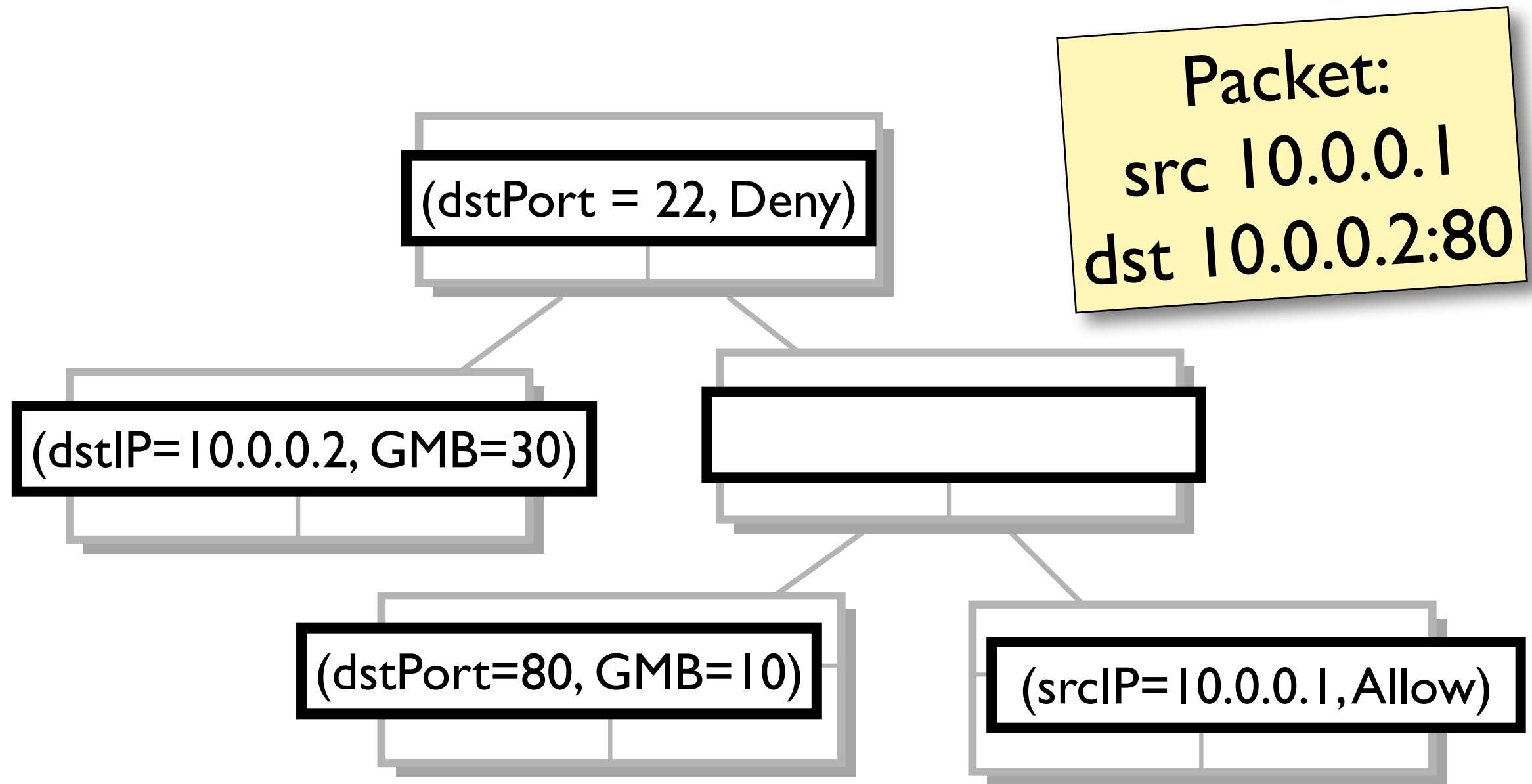
Dynamic Flow Processing



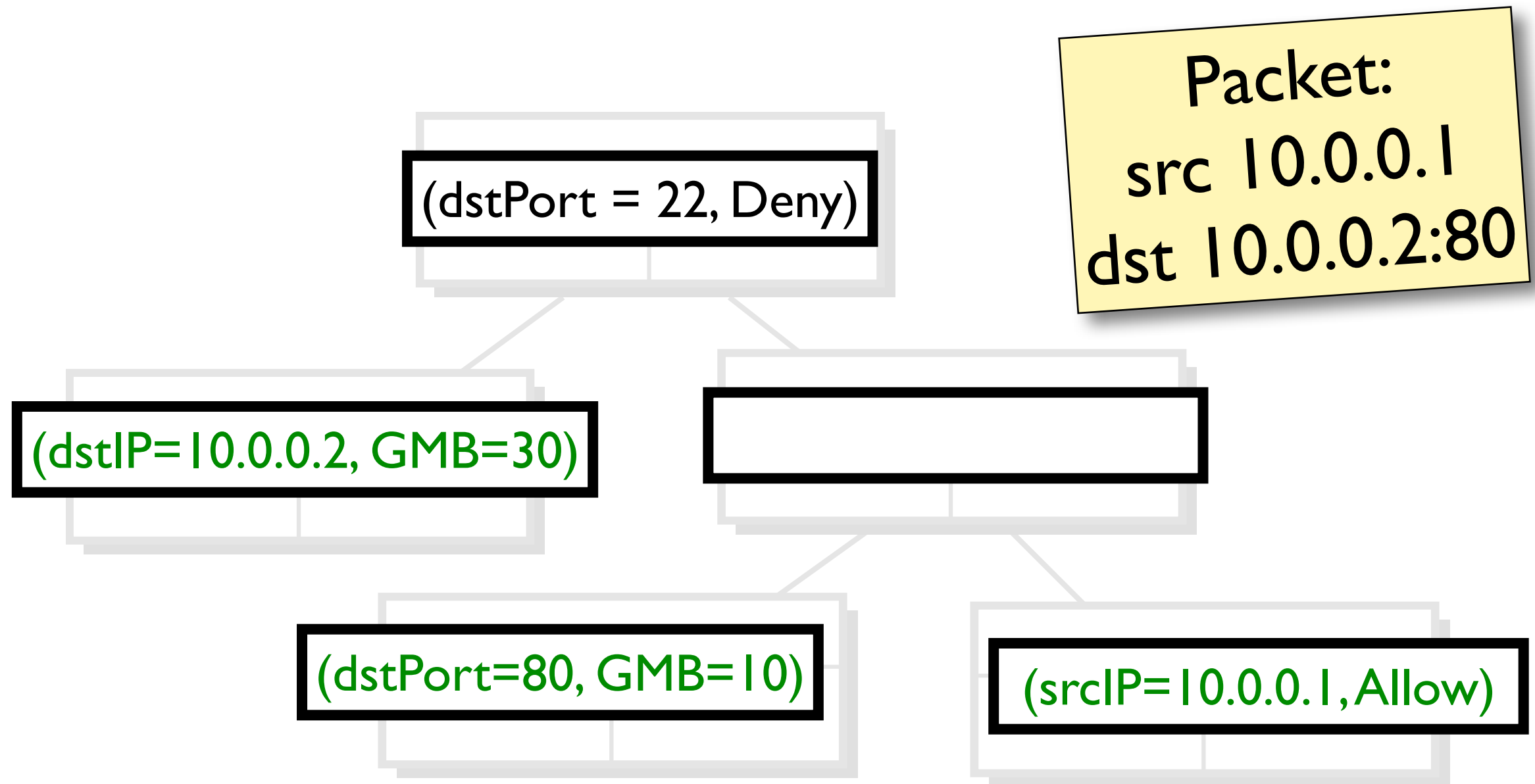
Hierarchy of Policies



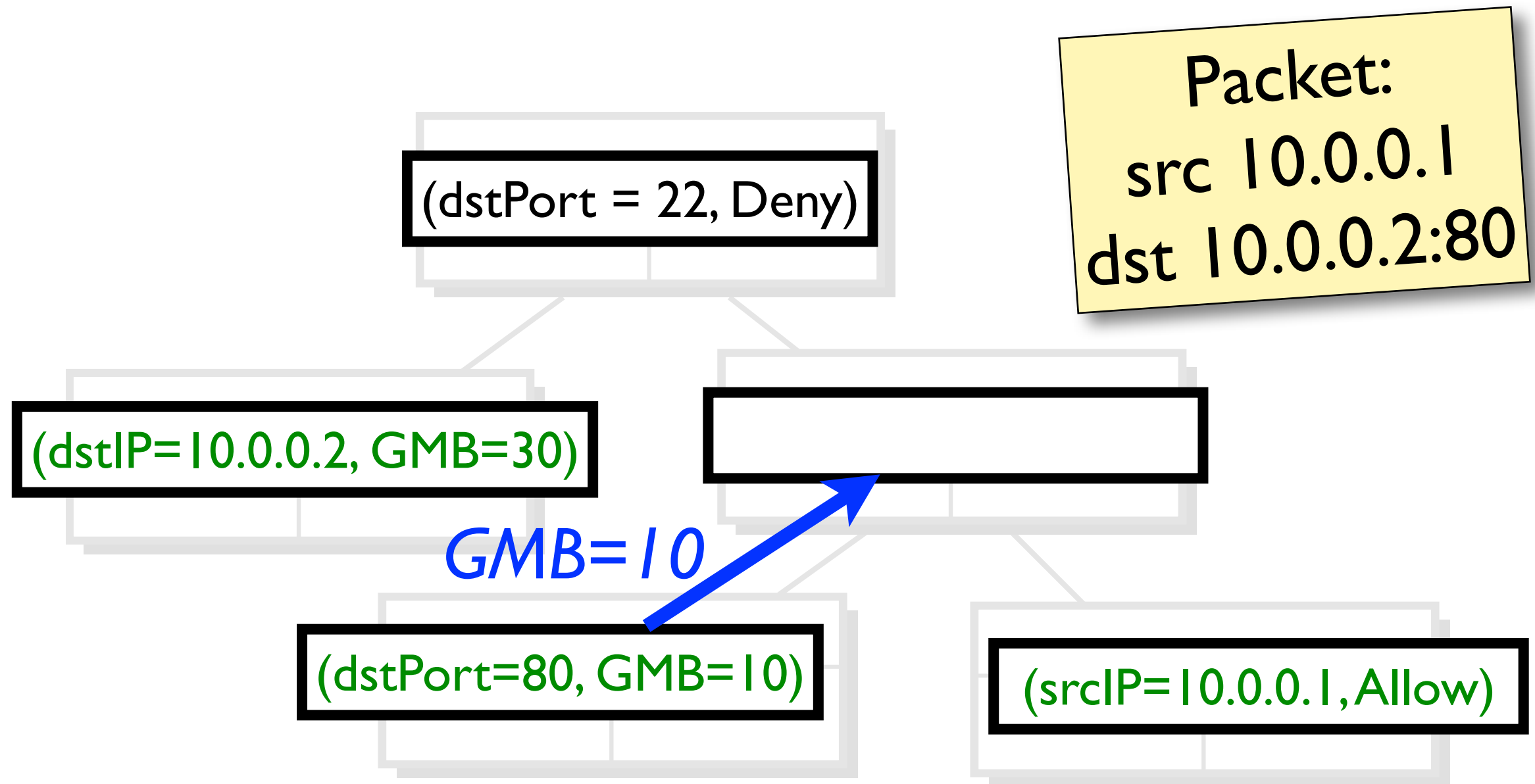
Hierarchy of Policies



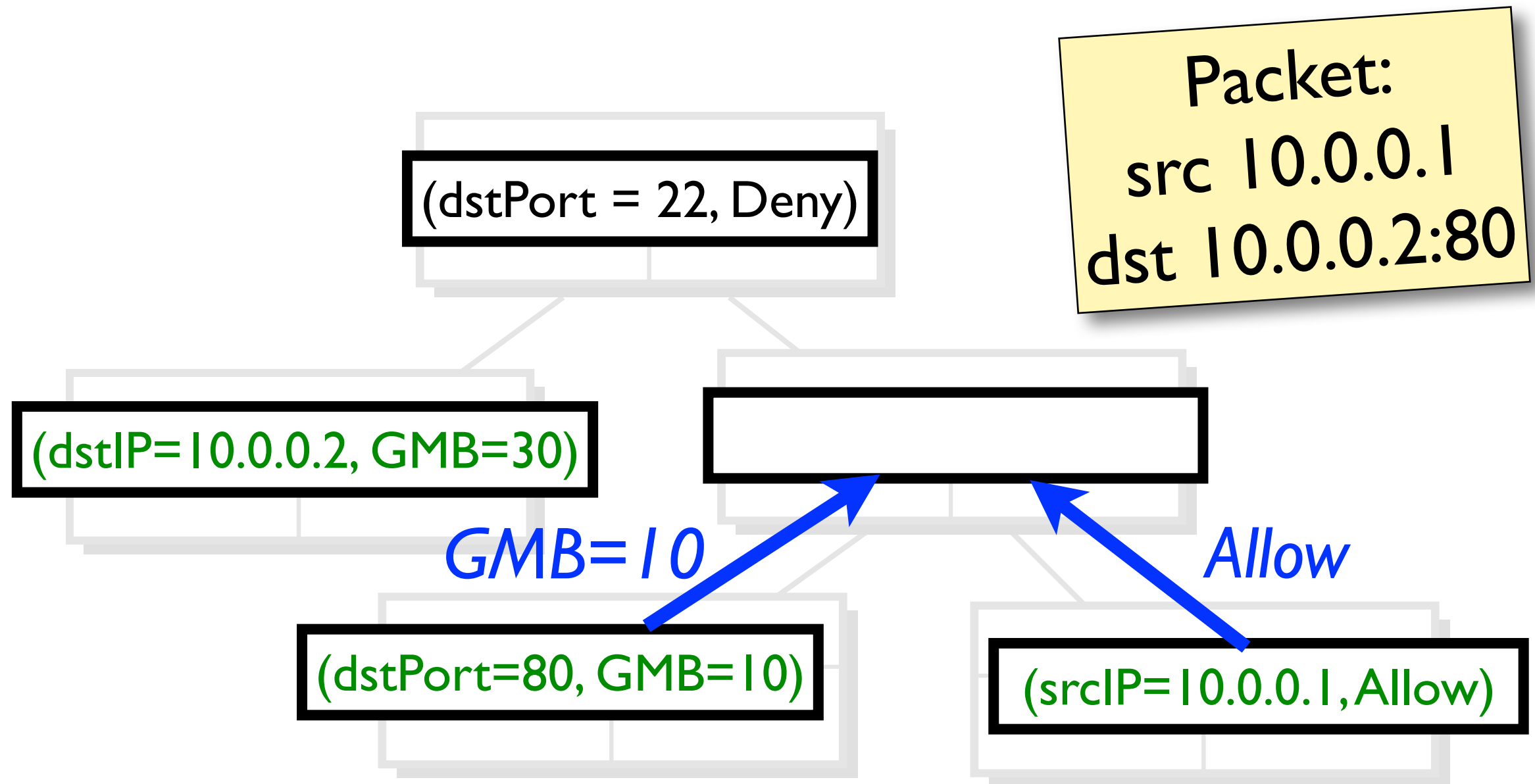
Hierarchy of Policies



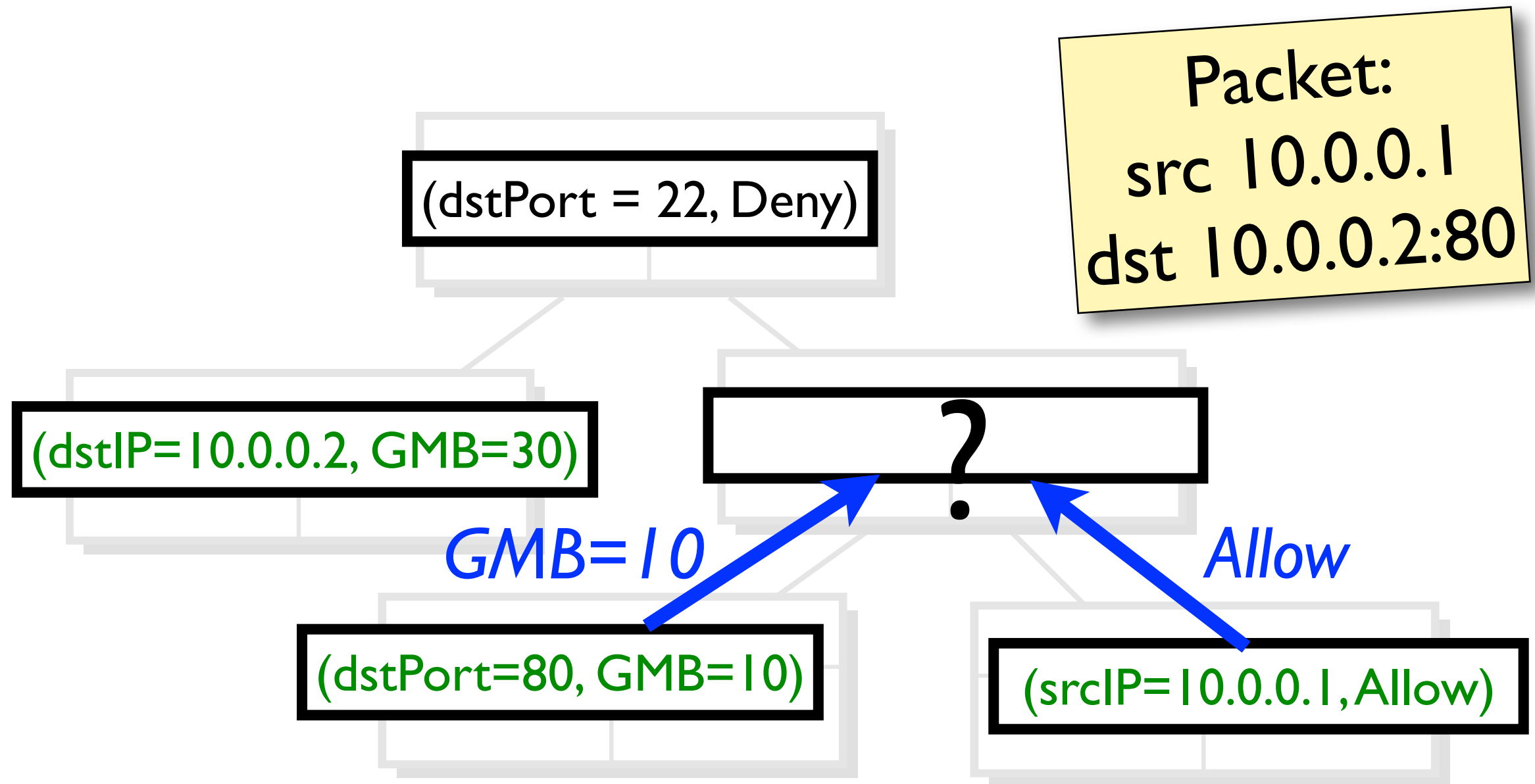
Hierarchical Flow Table



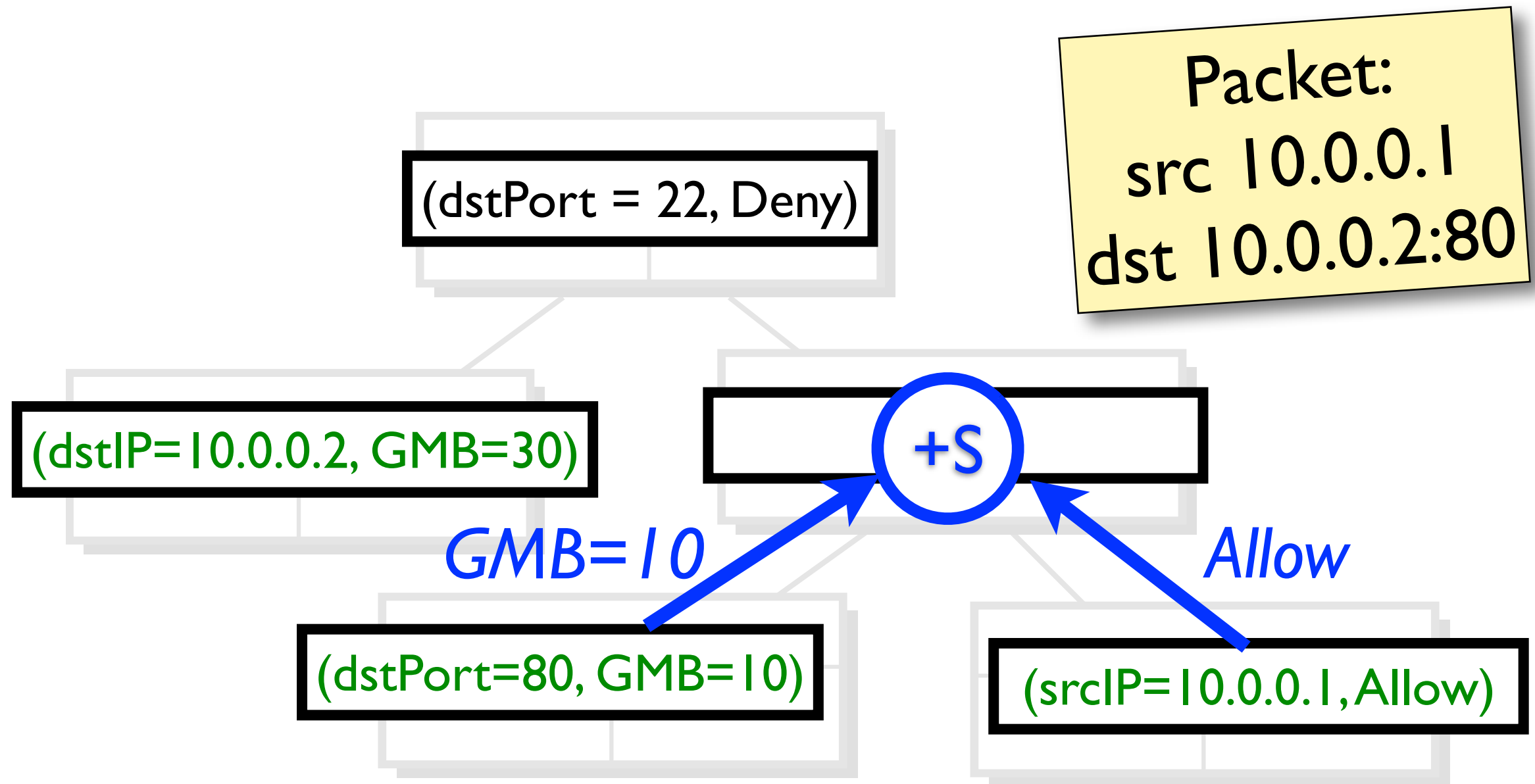
Hierarchical Flow Table



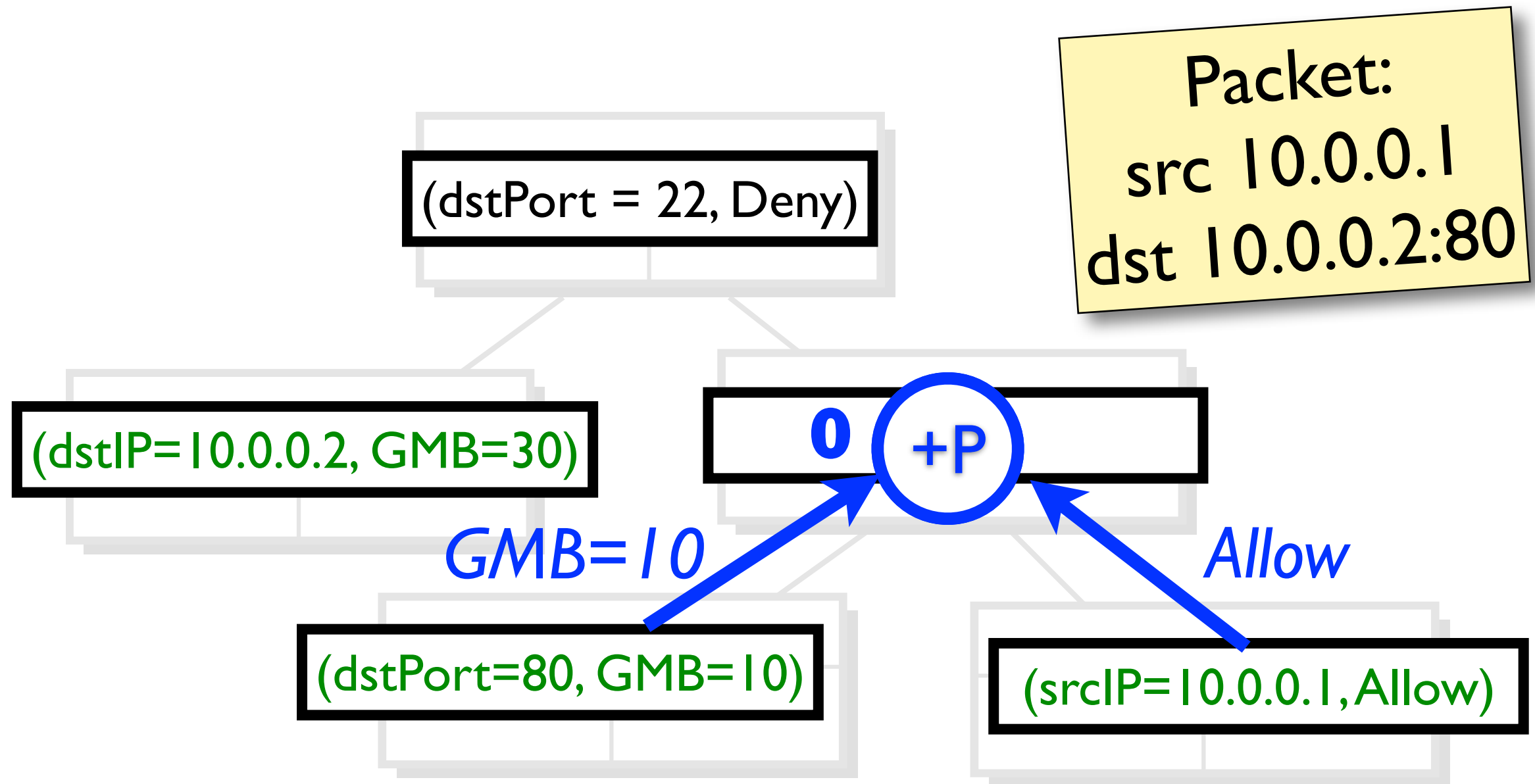
Hierarchical Flow Table



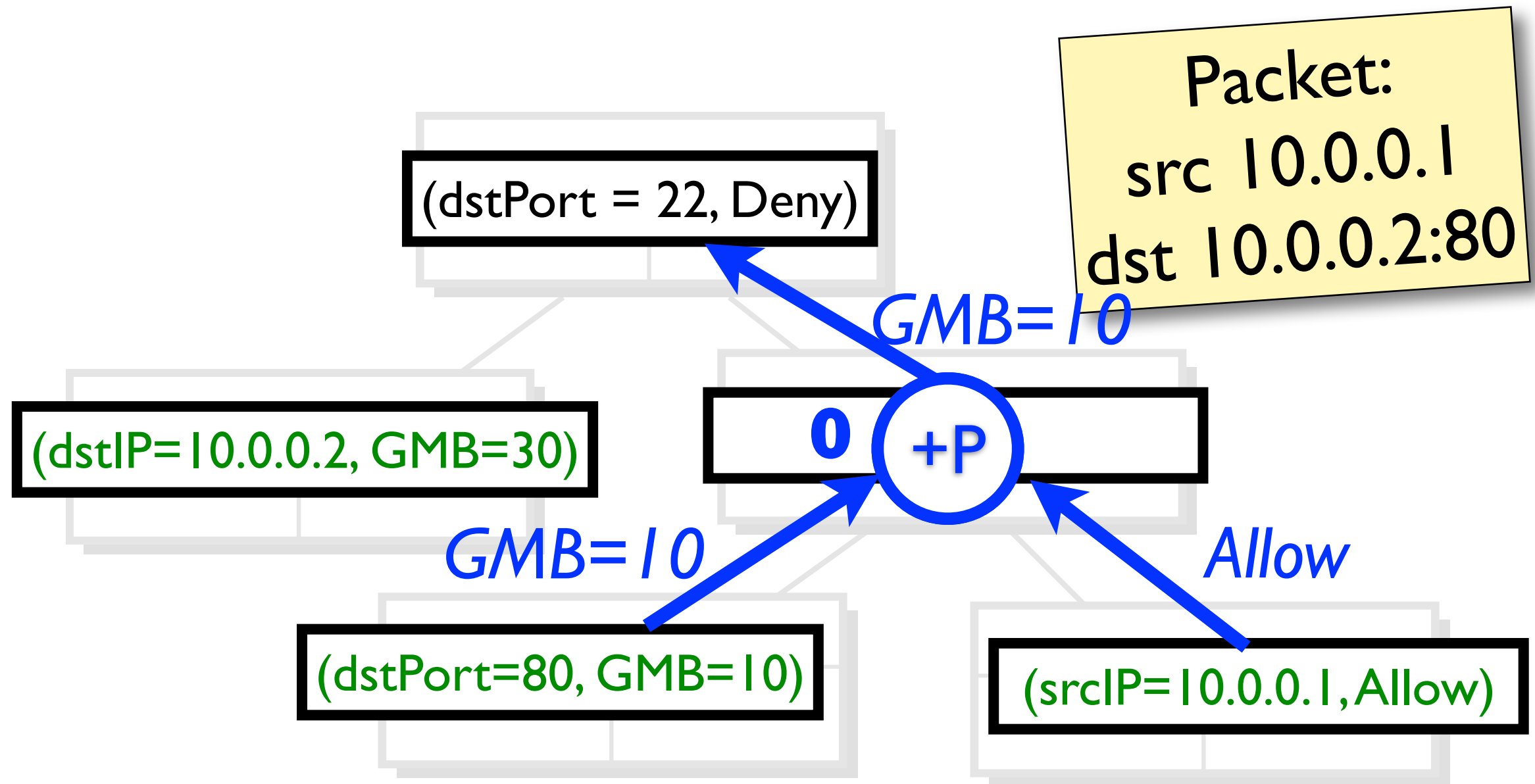
Hierarchical Flow Table



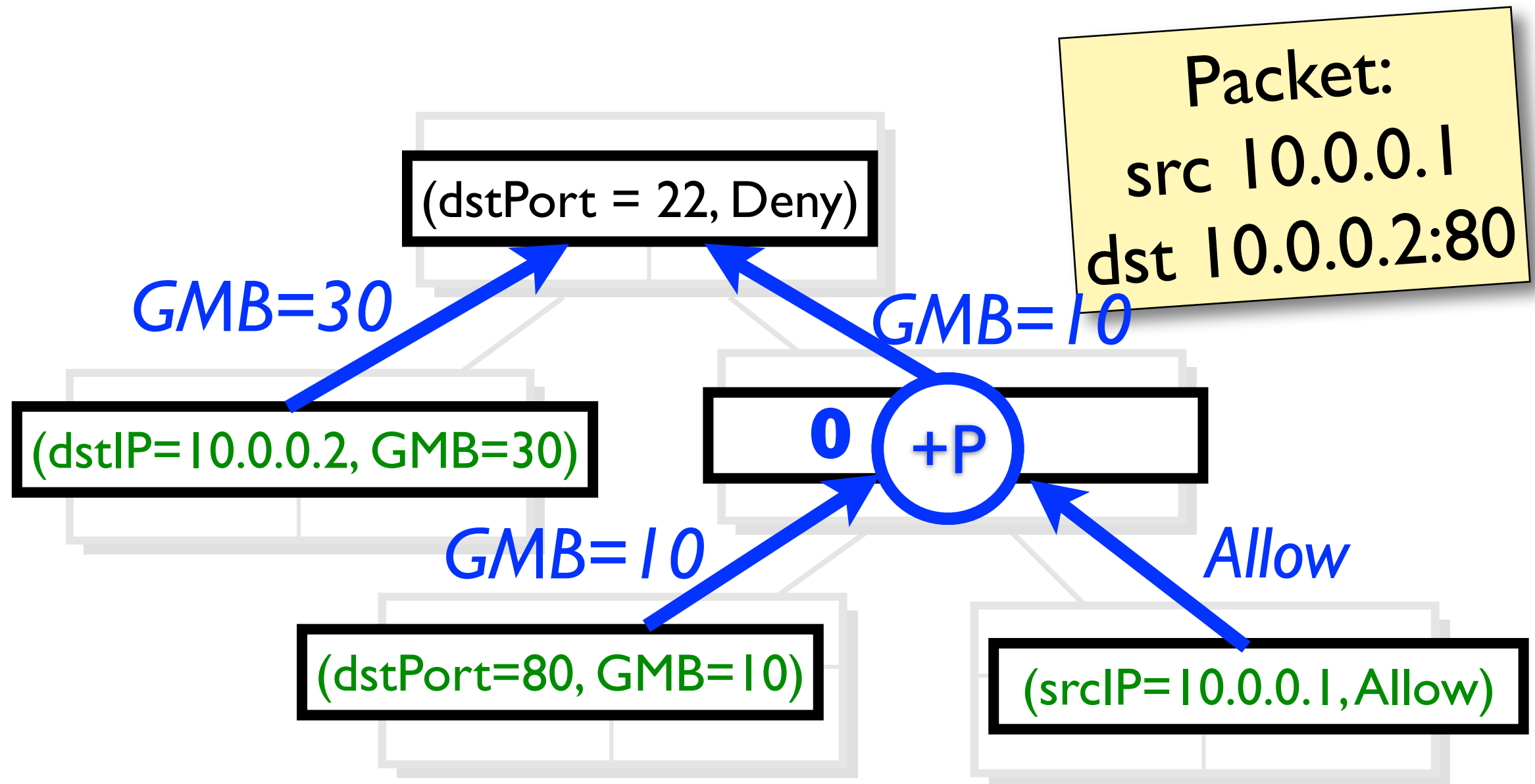
Hierarchical Flow Table



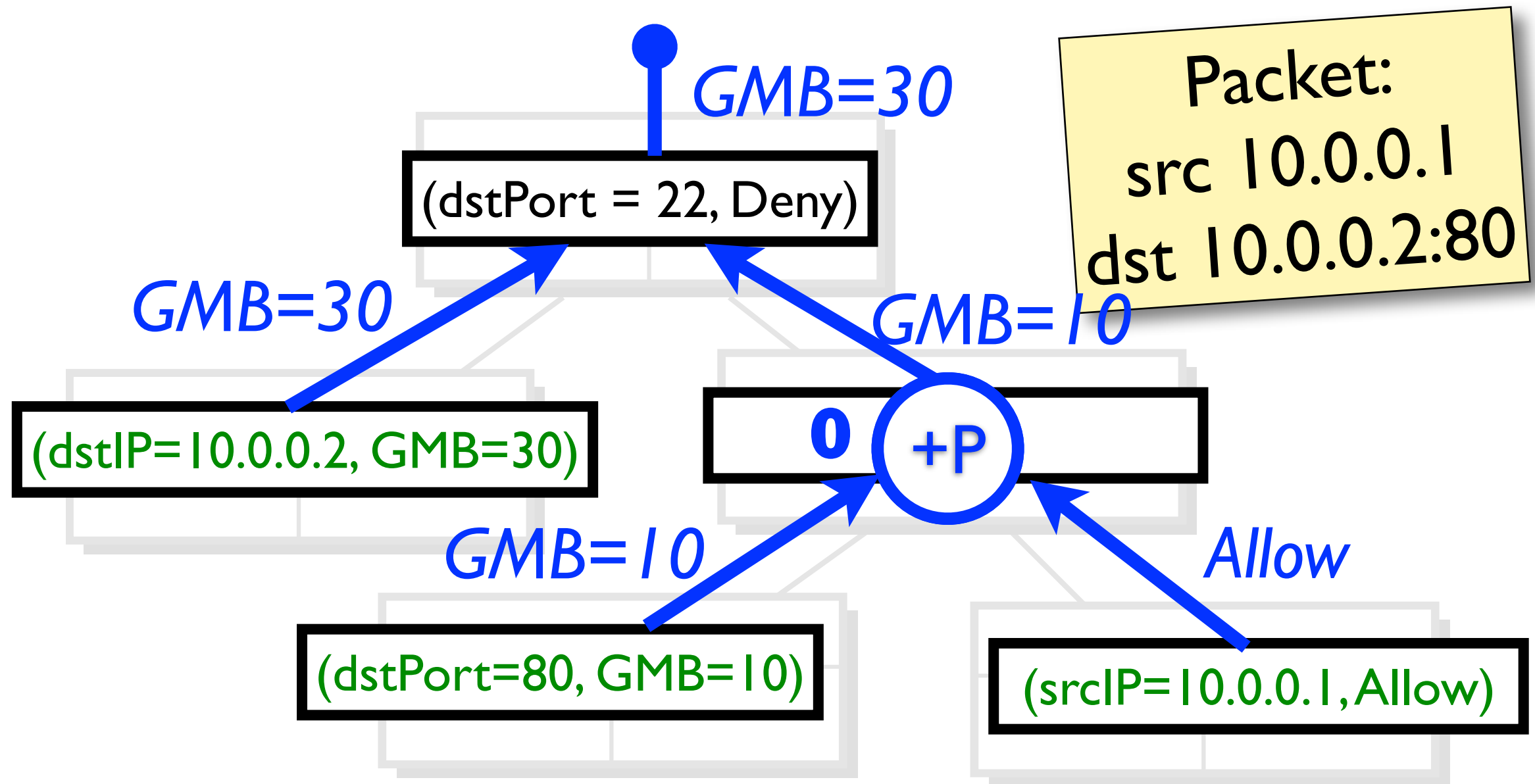
Hierarchical Flow Table



Hierarchical Flow Table



Hierarchical Flow Table

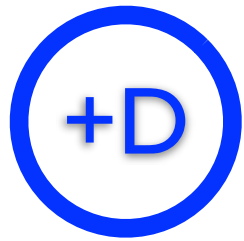


Hierarchical Flow Table

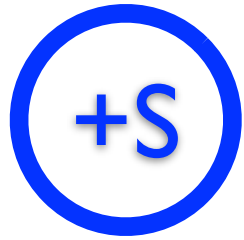
Requirements

Associative, **0-identity**

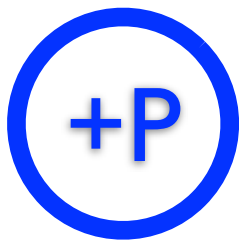
Commutative



In node



Sibling



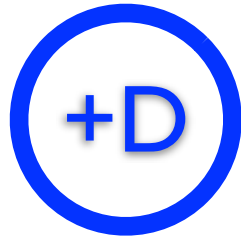
Parent-Sibling

HFT operators

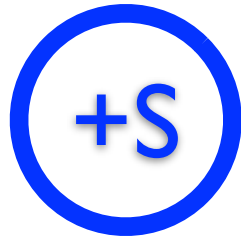
Associative, 0-identity

Requirements

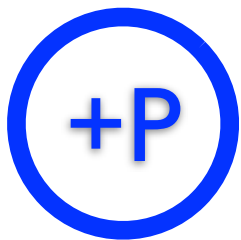
Commutative



In node



Sibling



Parent-Sibling

In PANE

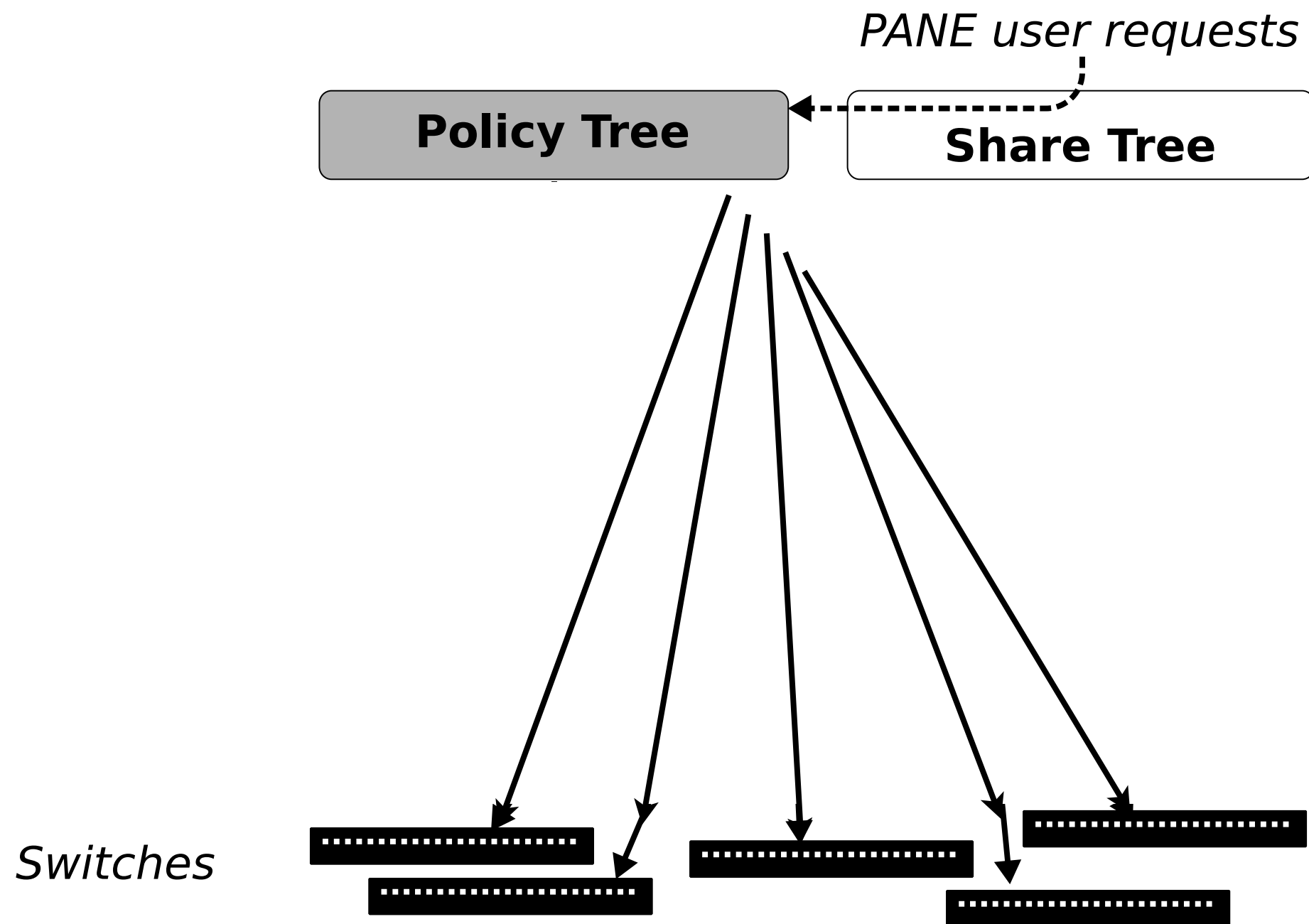
D and S identical.

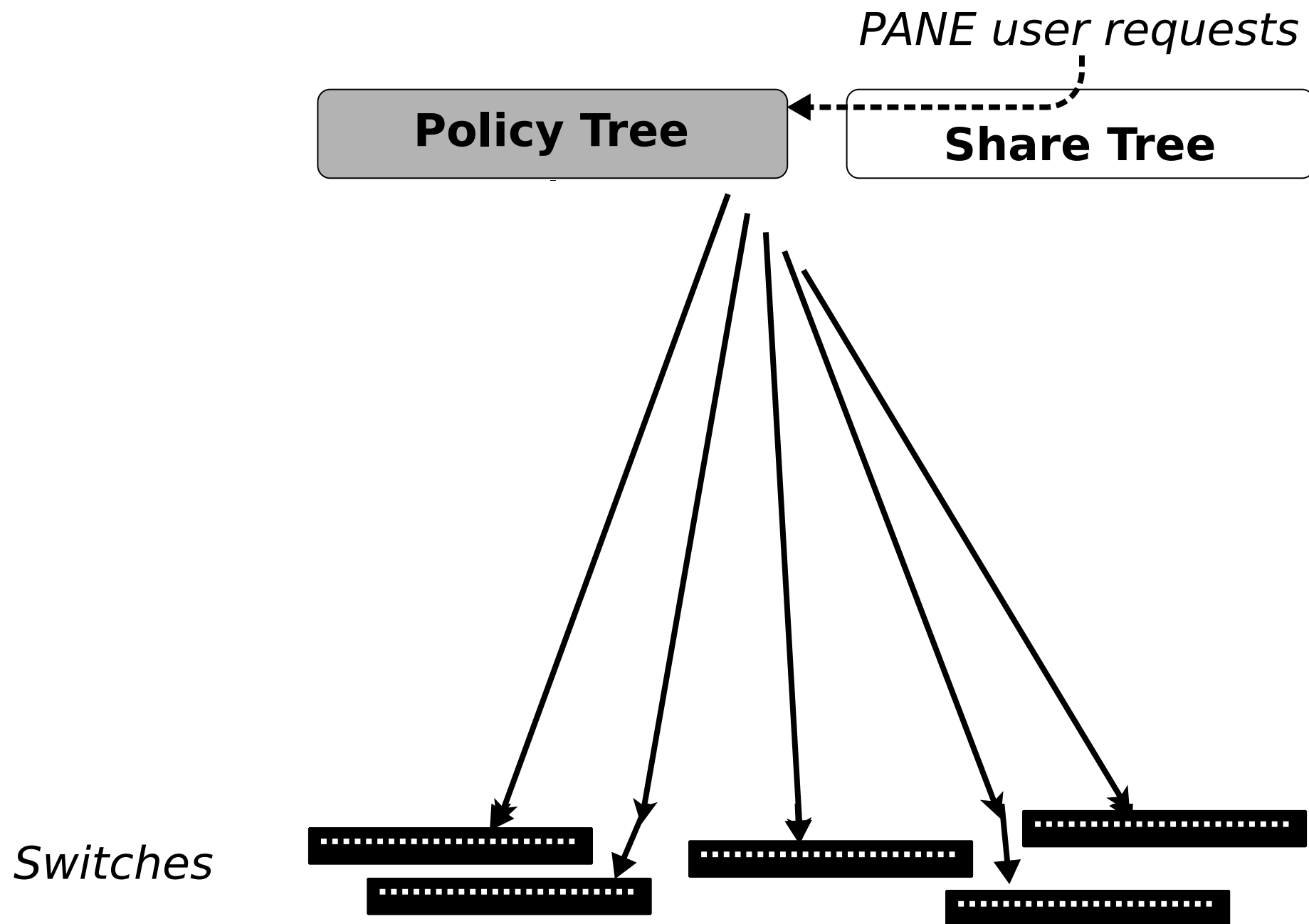
Deny overrides Allow.
GMB combines as **max**

Child overrides Parent
for Access Control
GMB combines as **max**

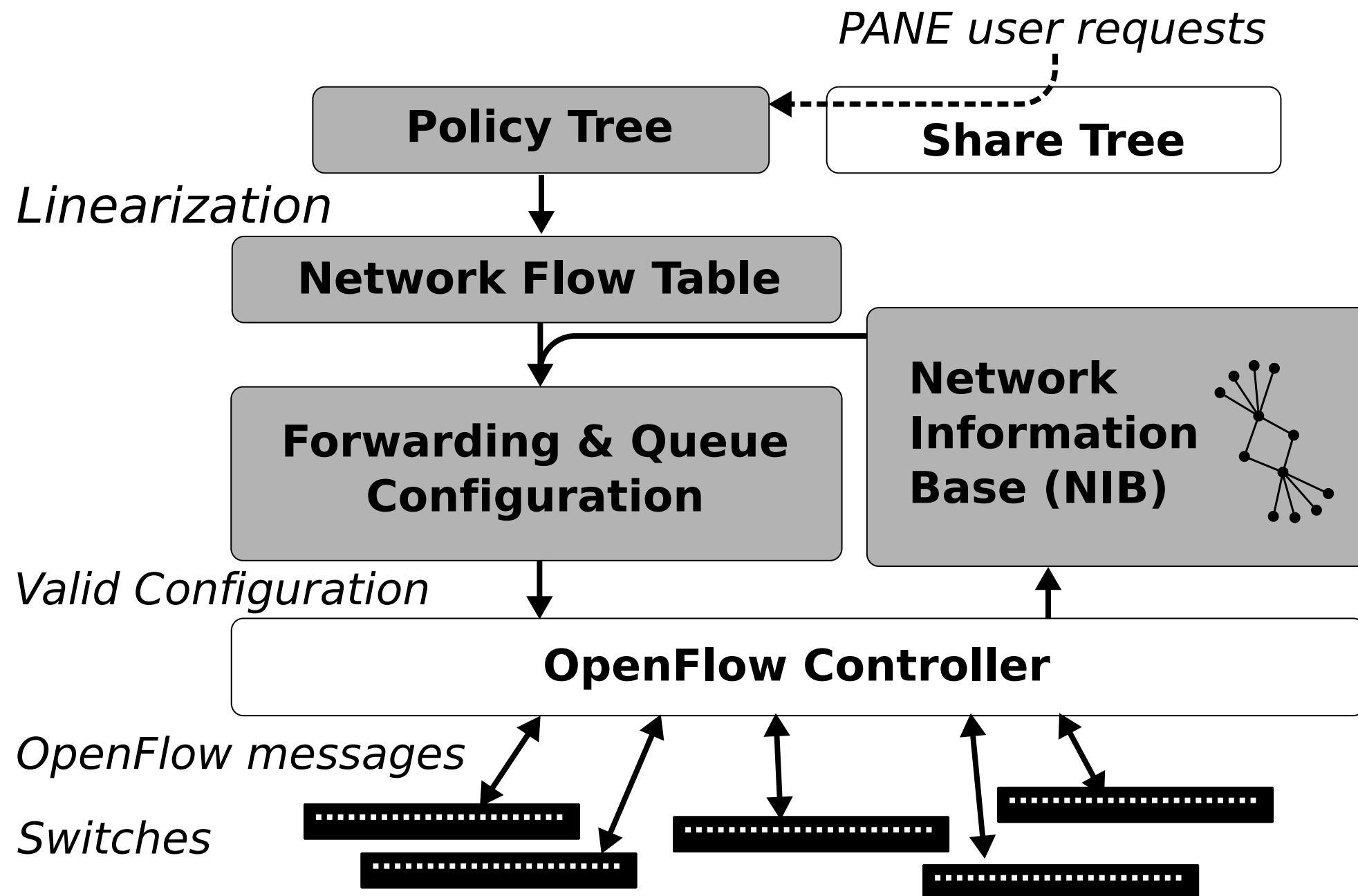
HFT operators

There's a slight problem...



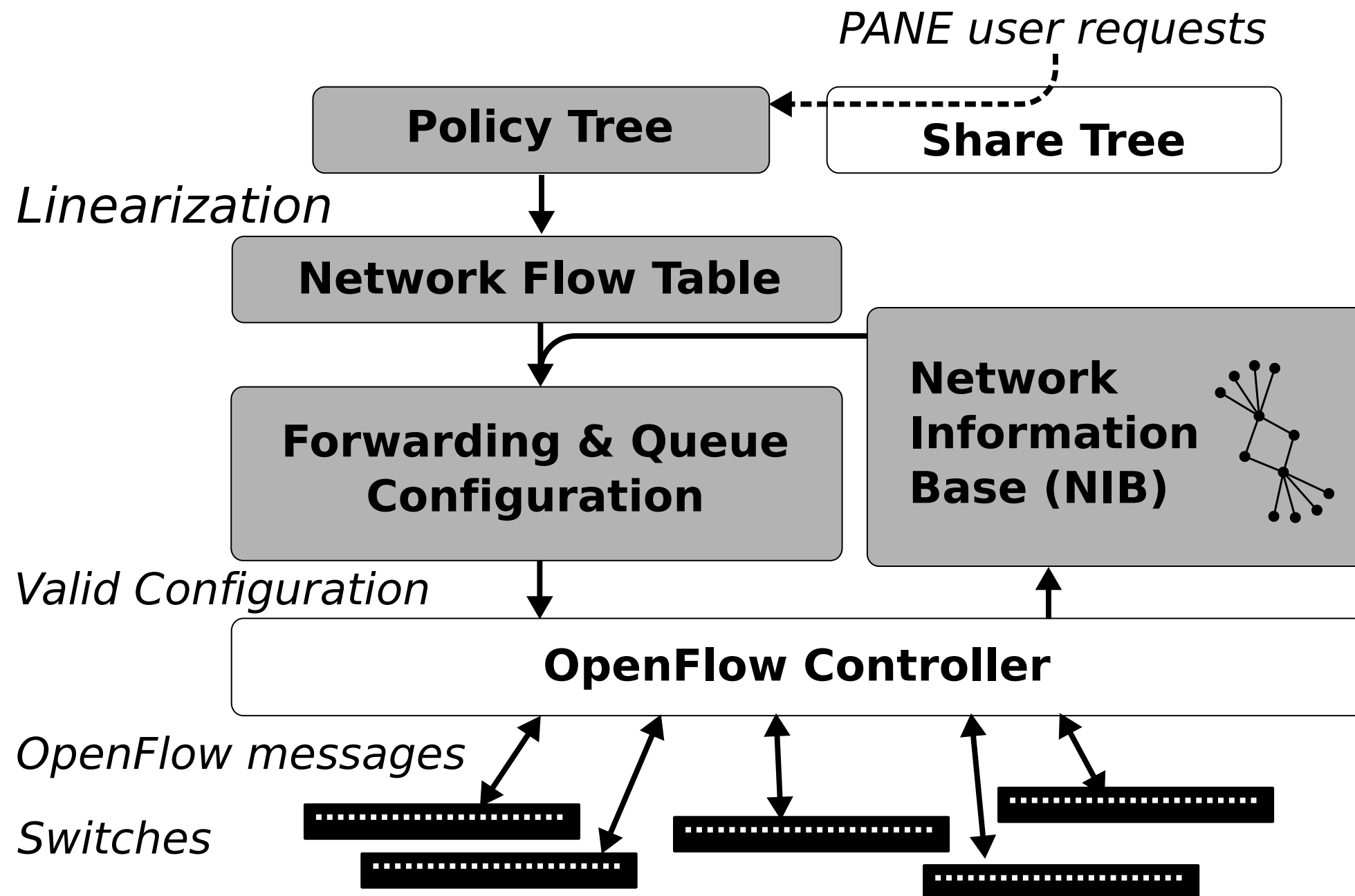


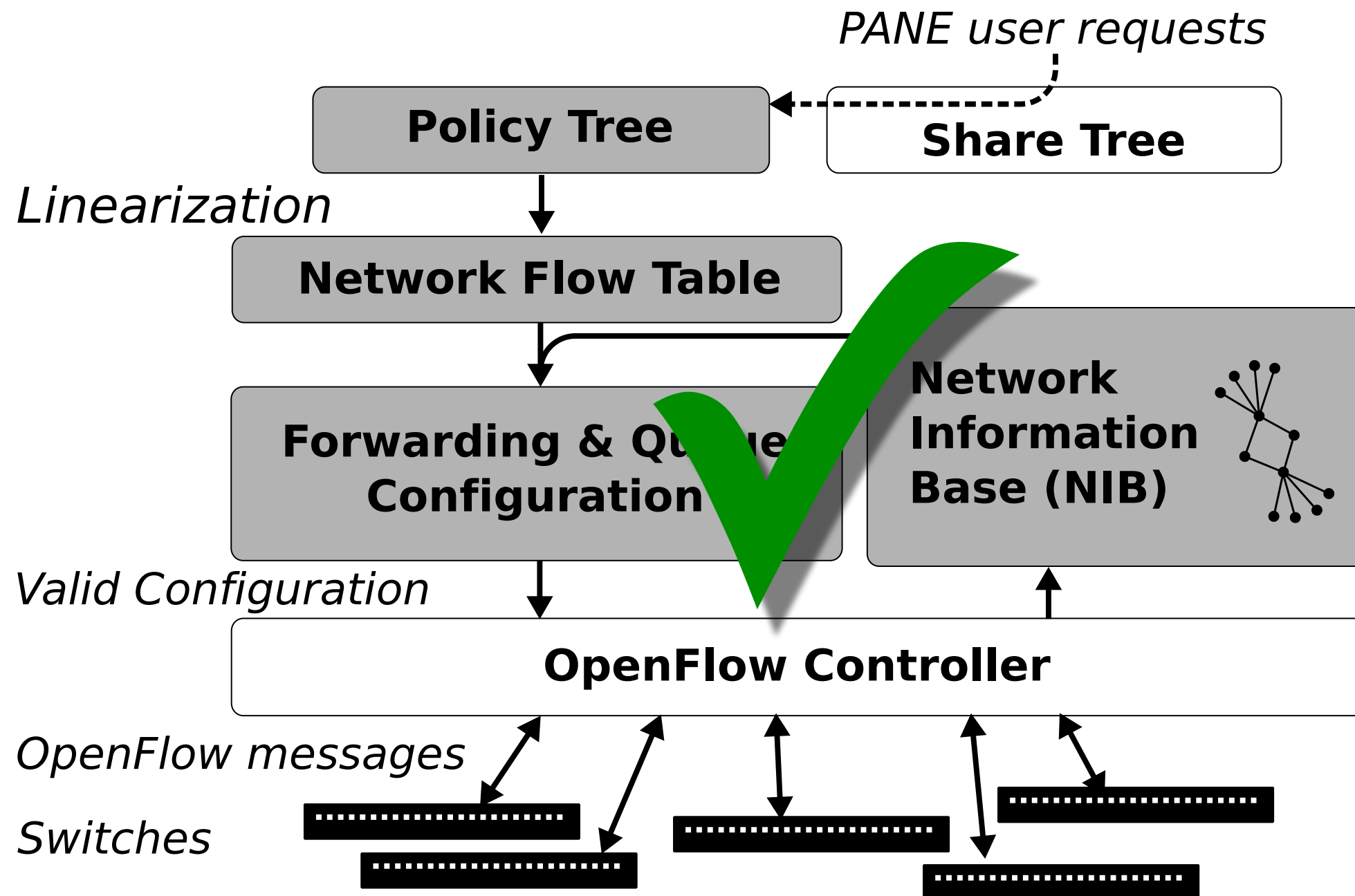
Switches don't grow trees

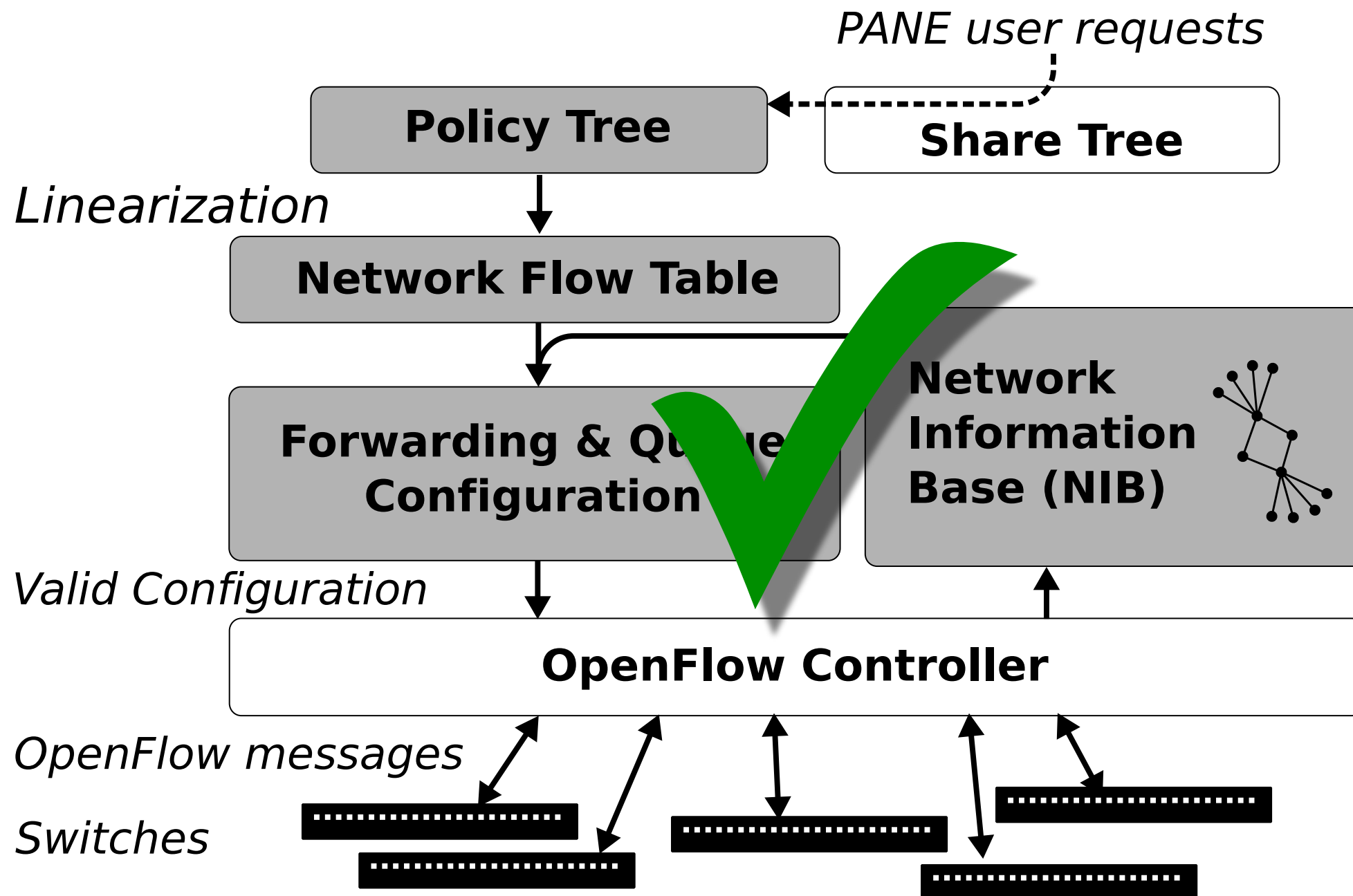


Switches don't grow trees

Current Status







We've built all these components
Currently have Access Control and GMB



Toy Evaluation

Protecting Zookeeper

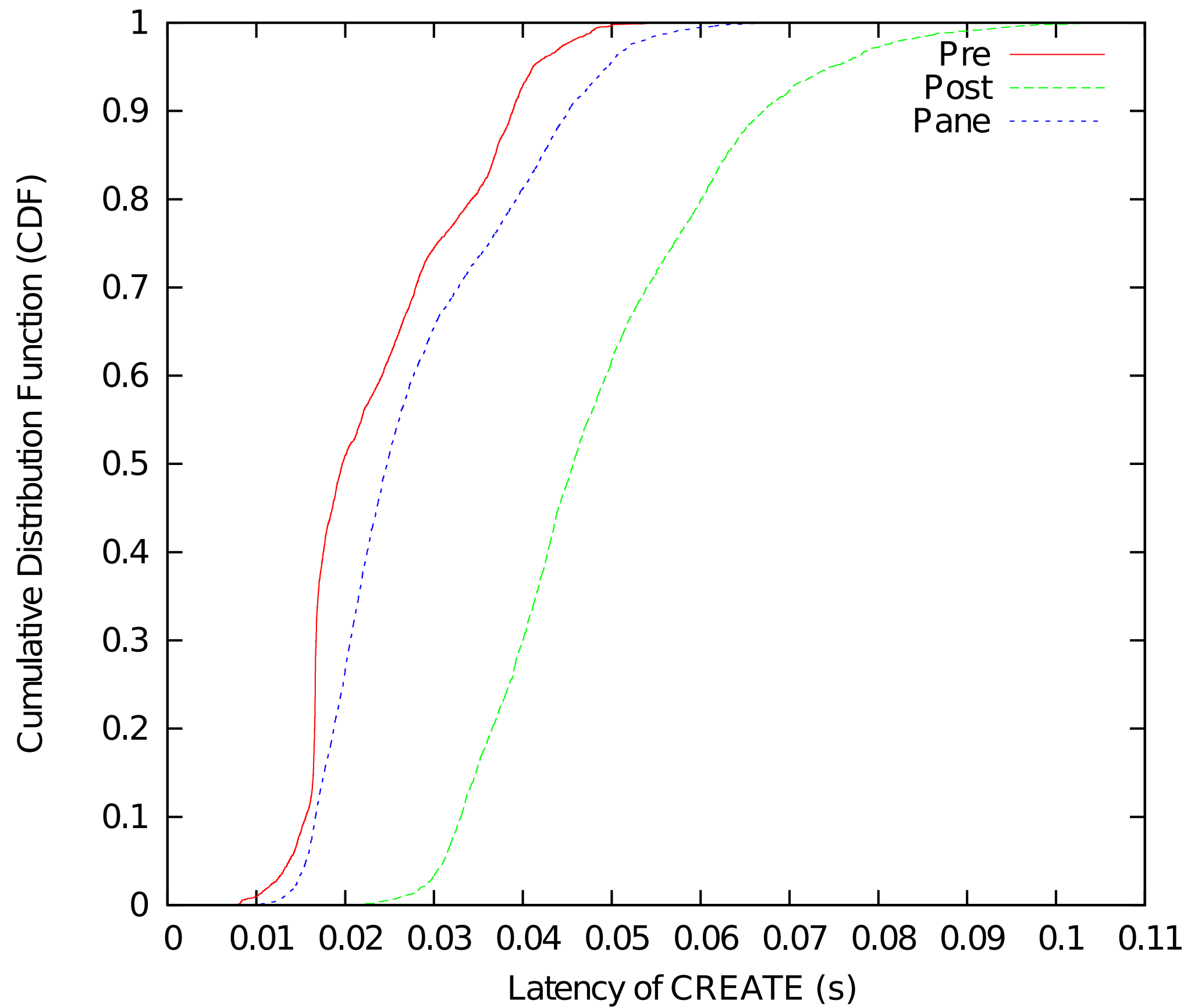
5 PANE-enabled Zookeeper servers

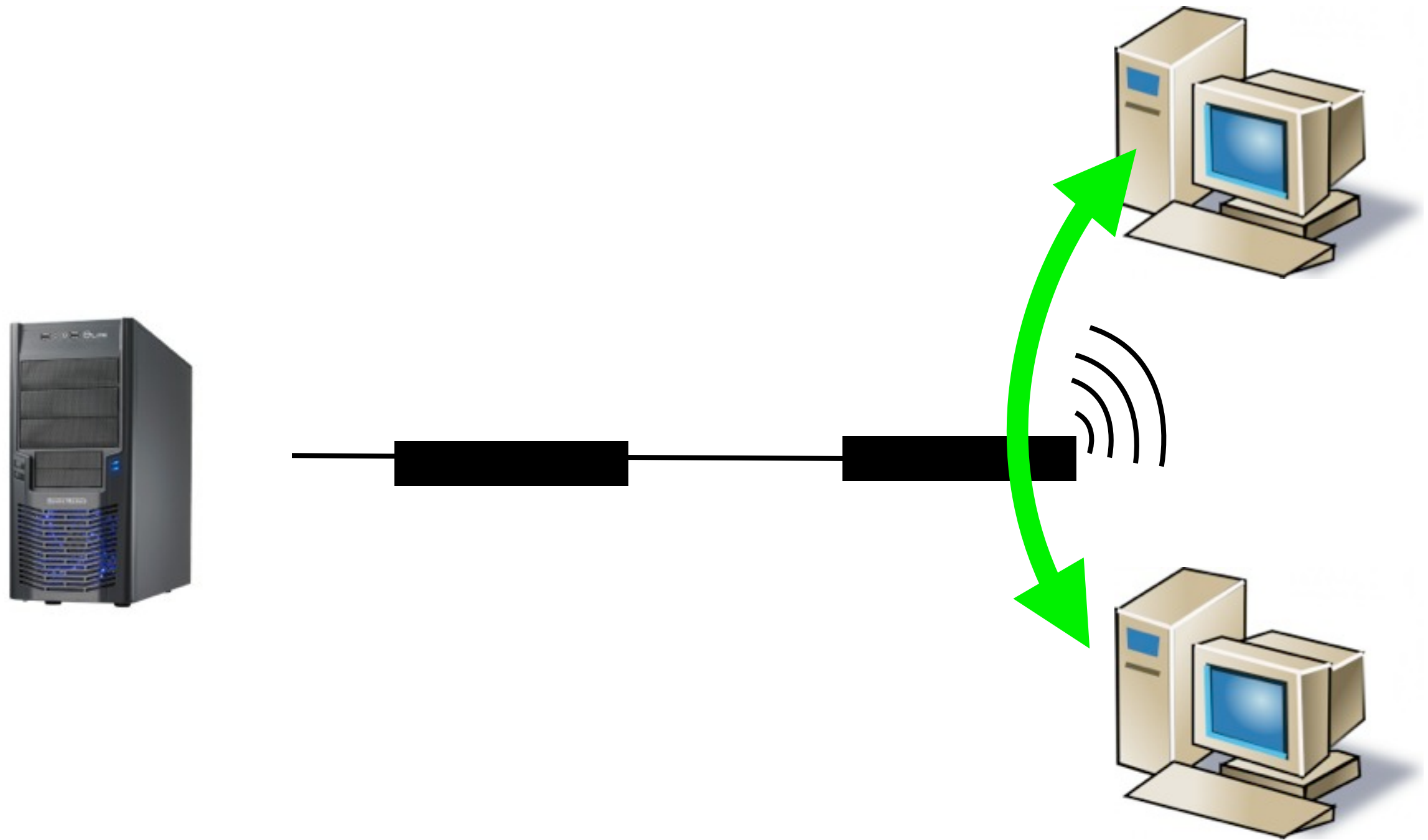
1 client

Connected via single OpenVSwitch (3.3Gbps)

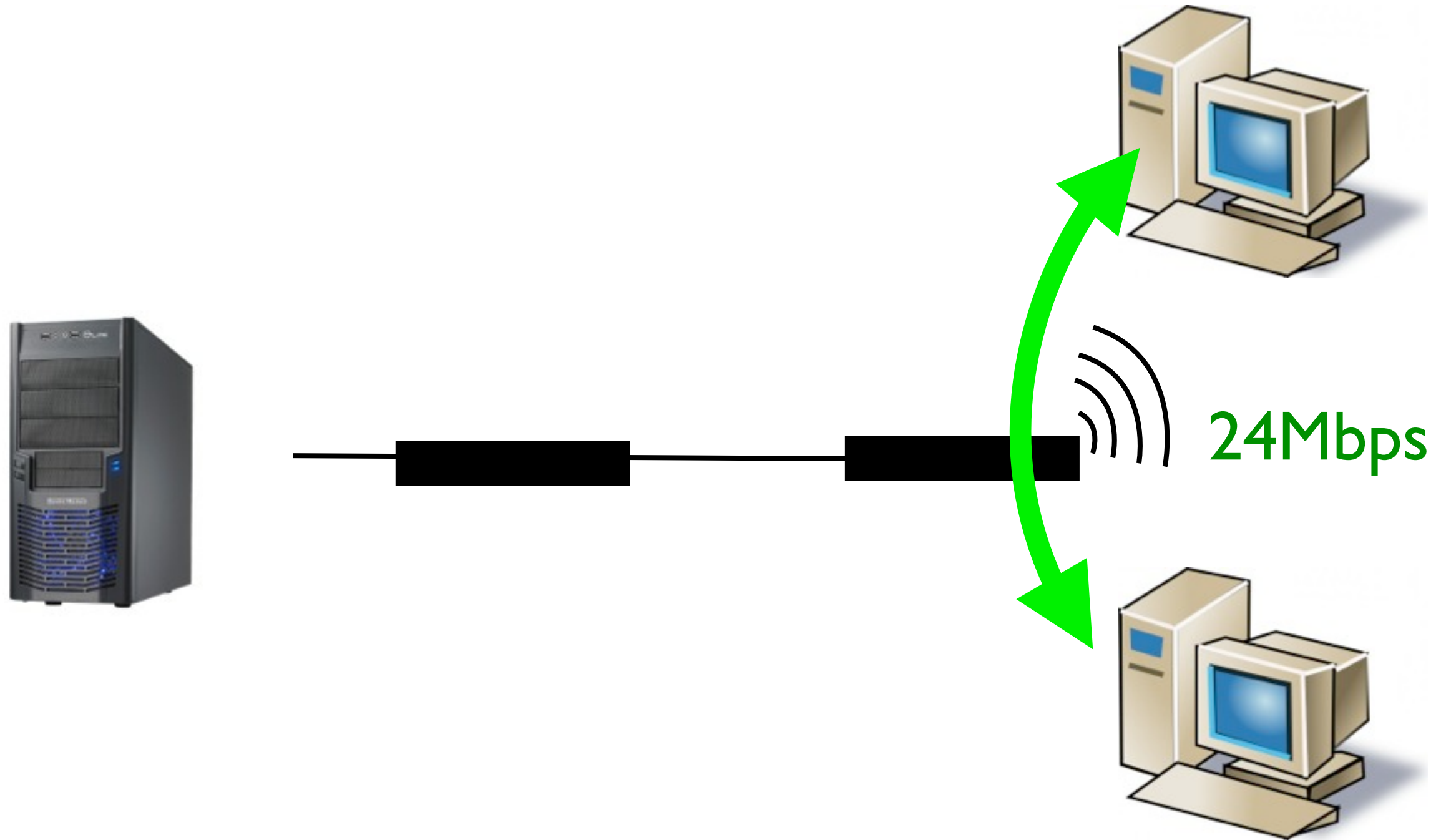
iPerf generating load on all links

Protecting Zookeeper

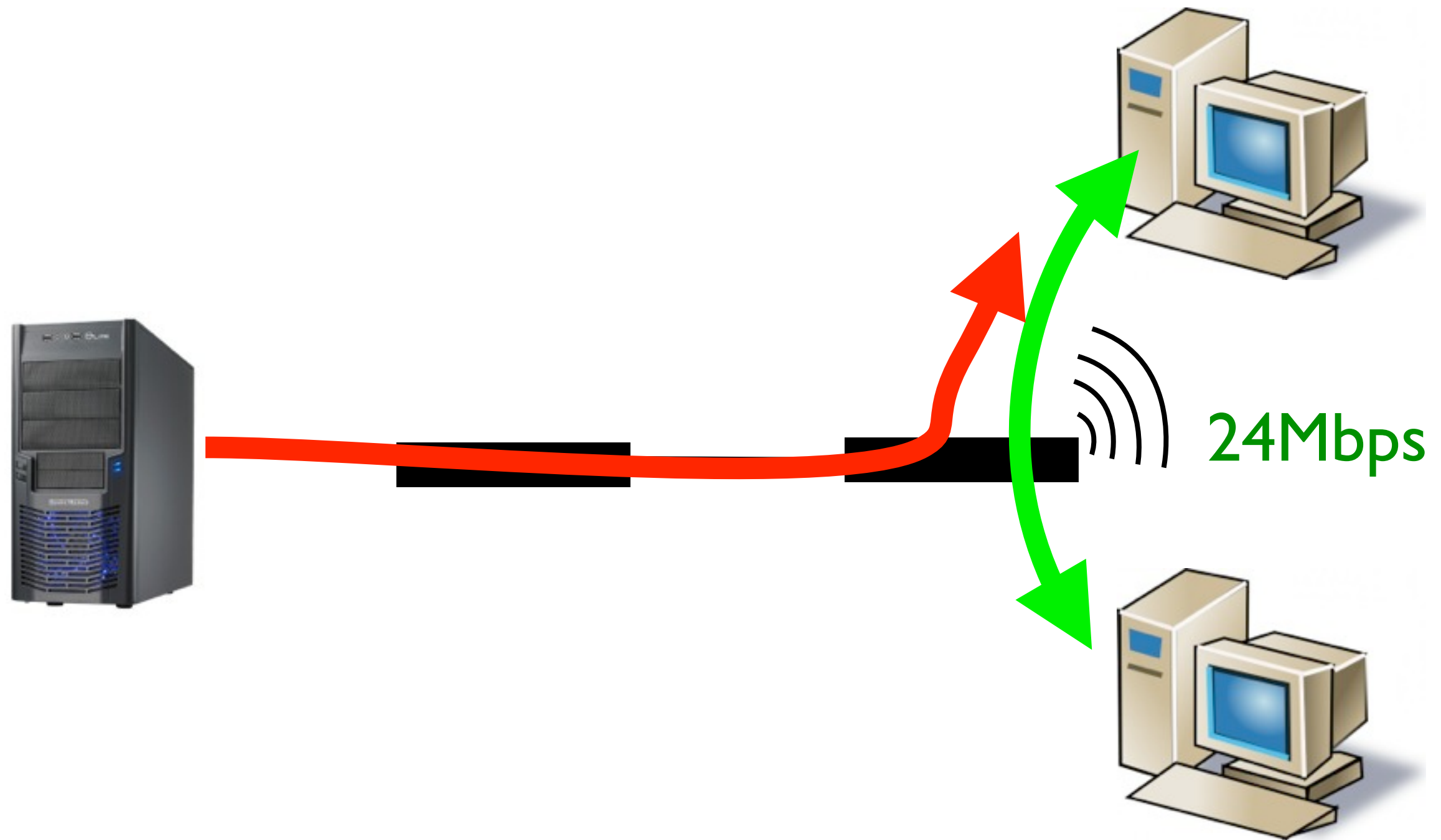




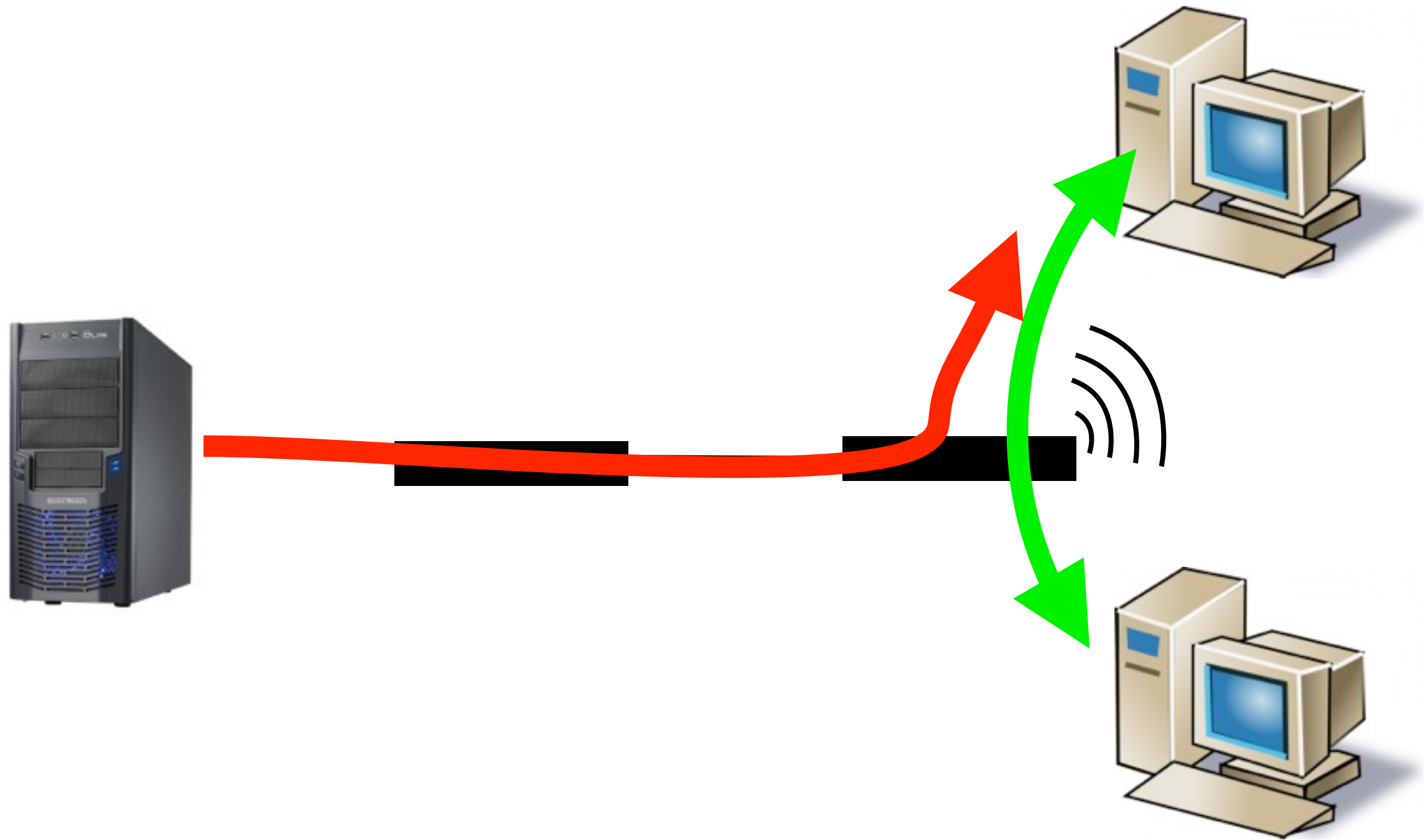
Denial-of-service



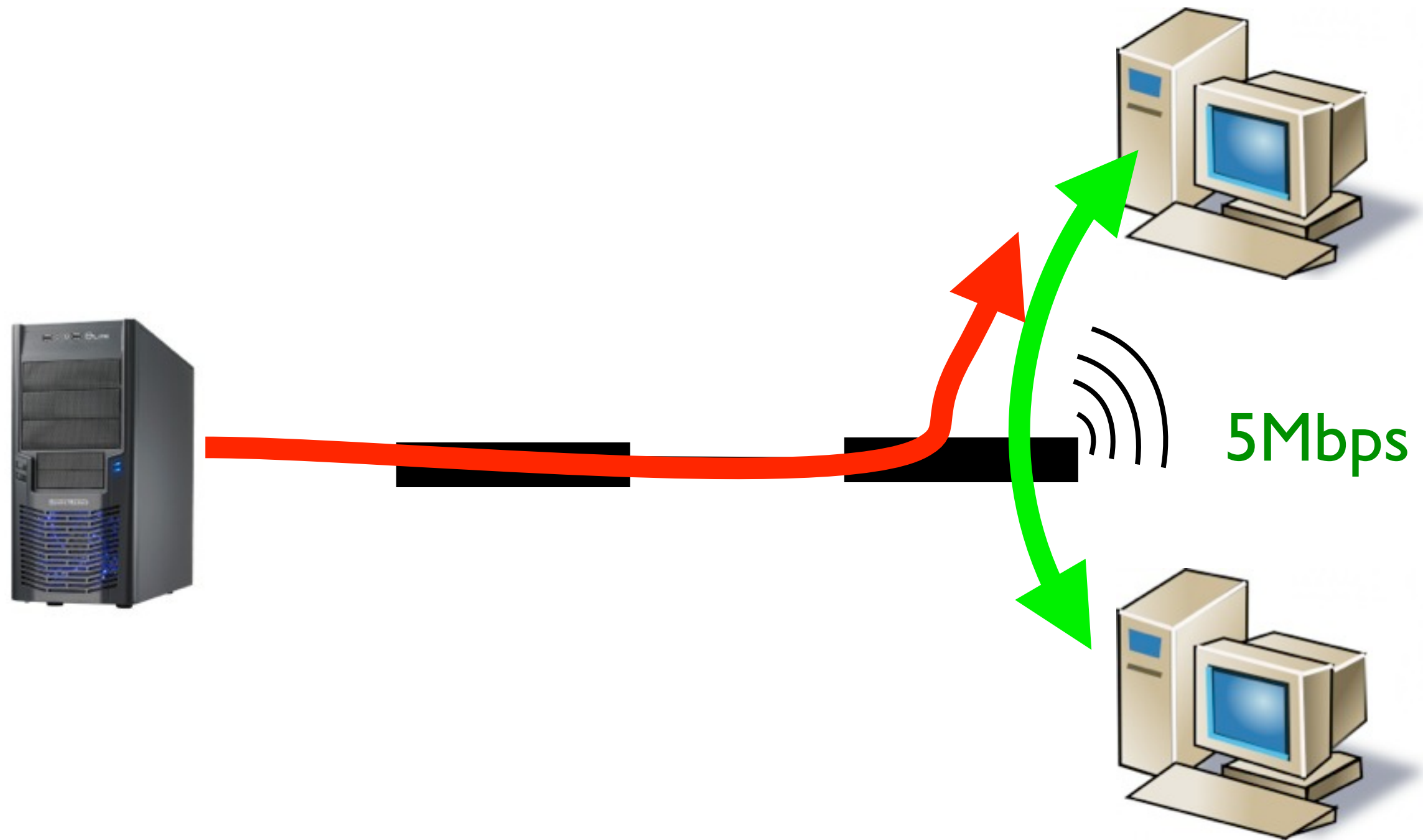
Denial-of-service



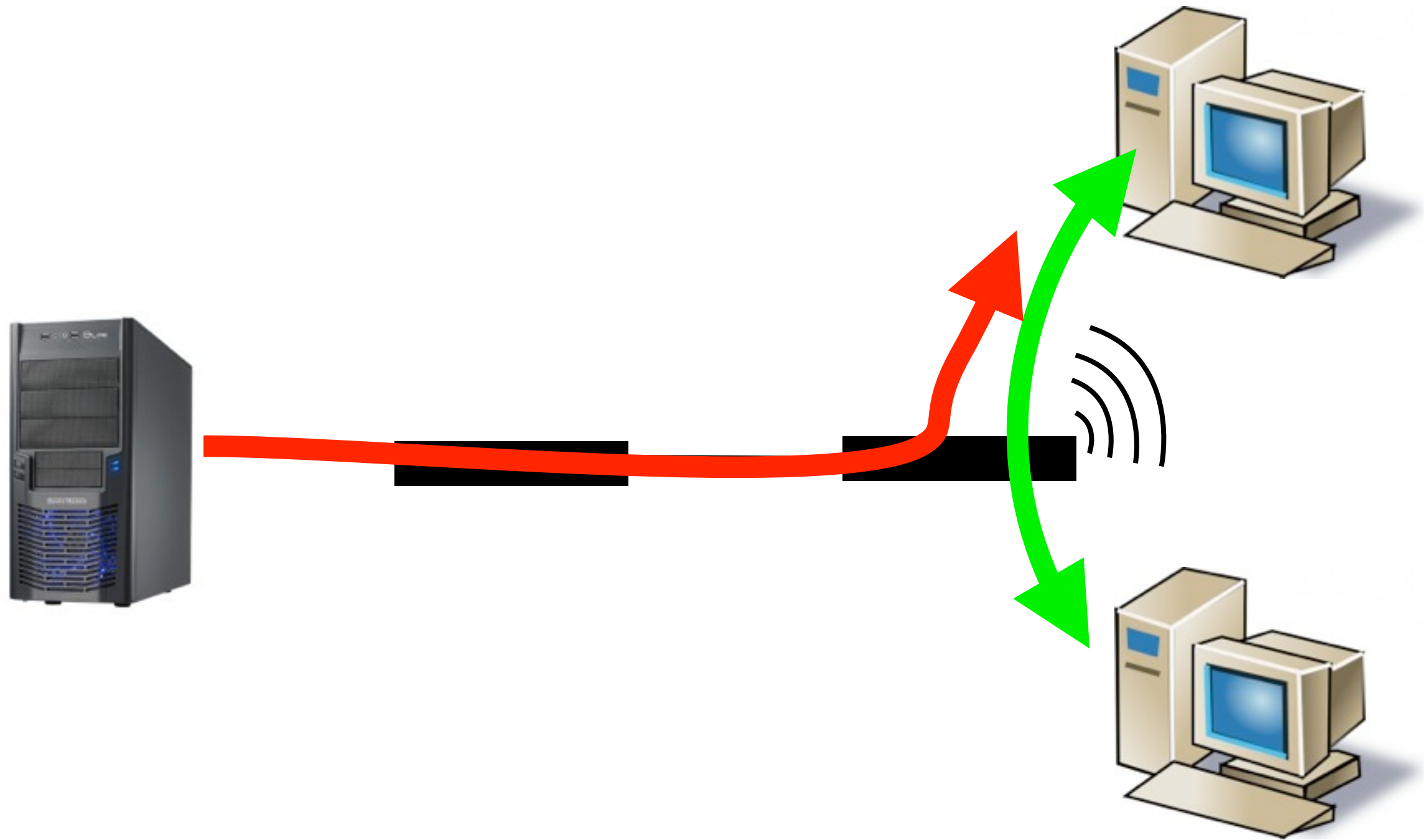
Denial-of-service



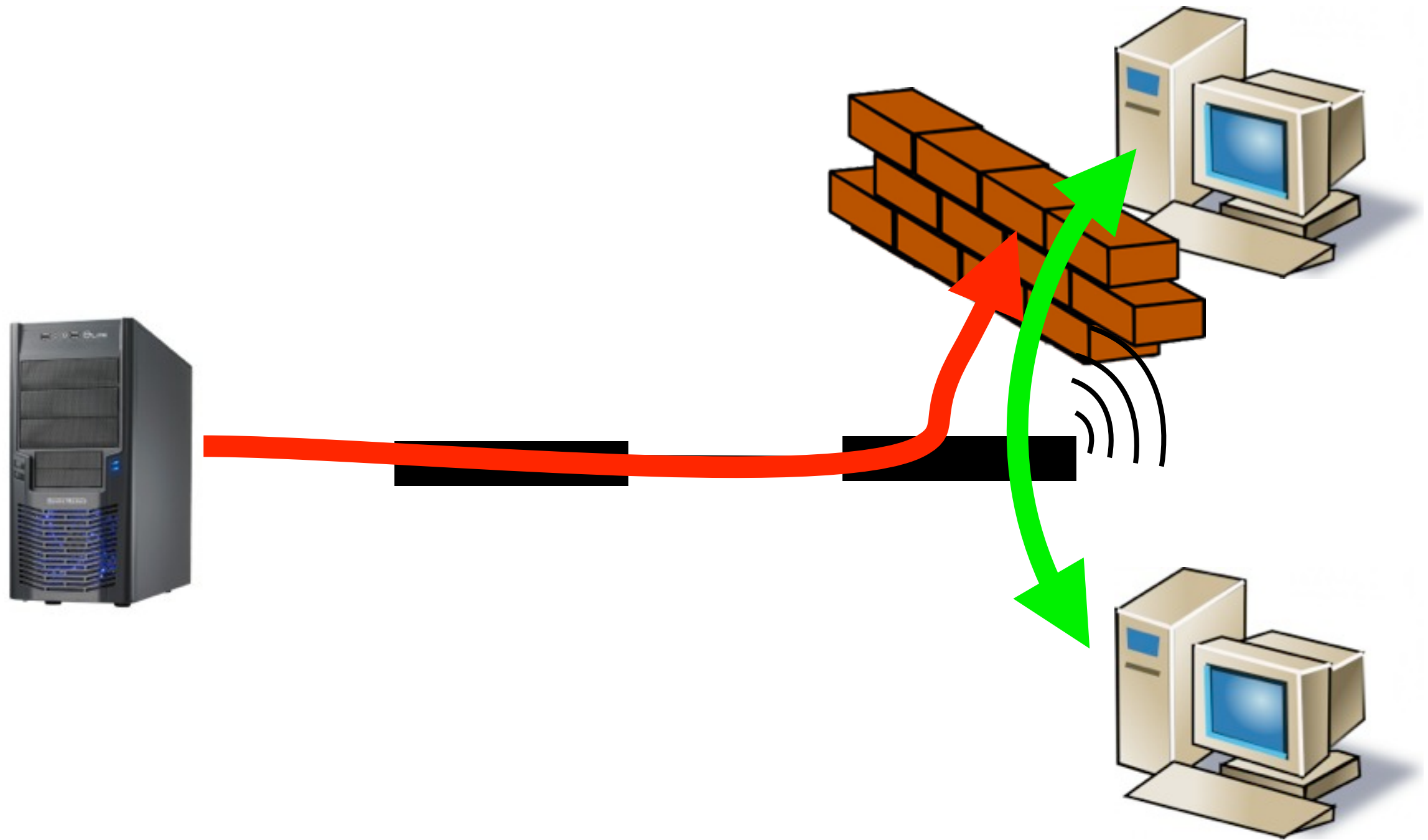
Denial-of-service



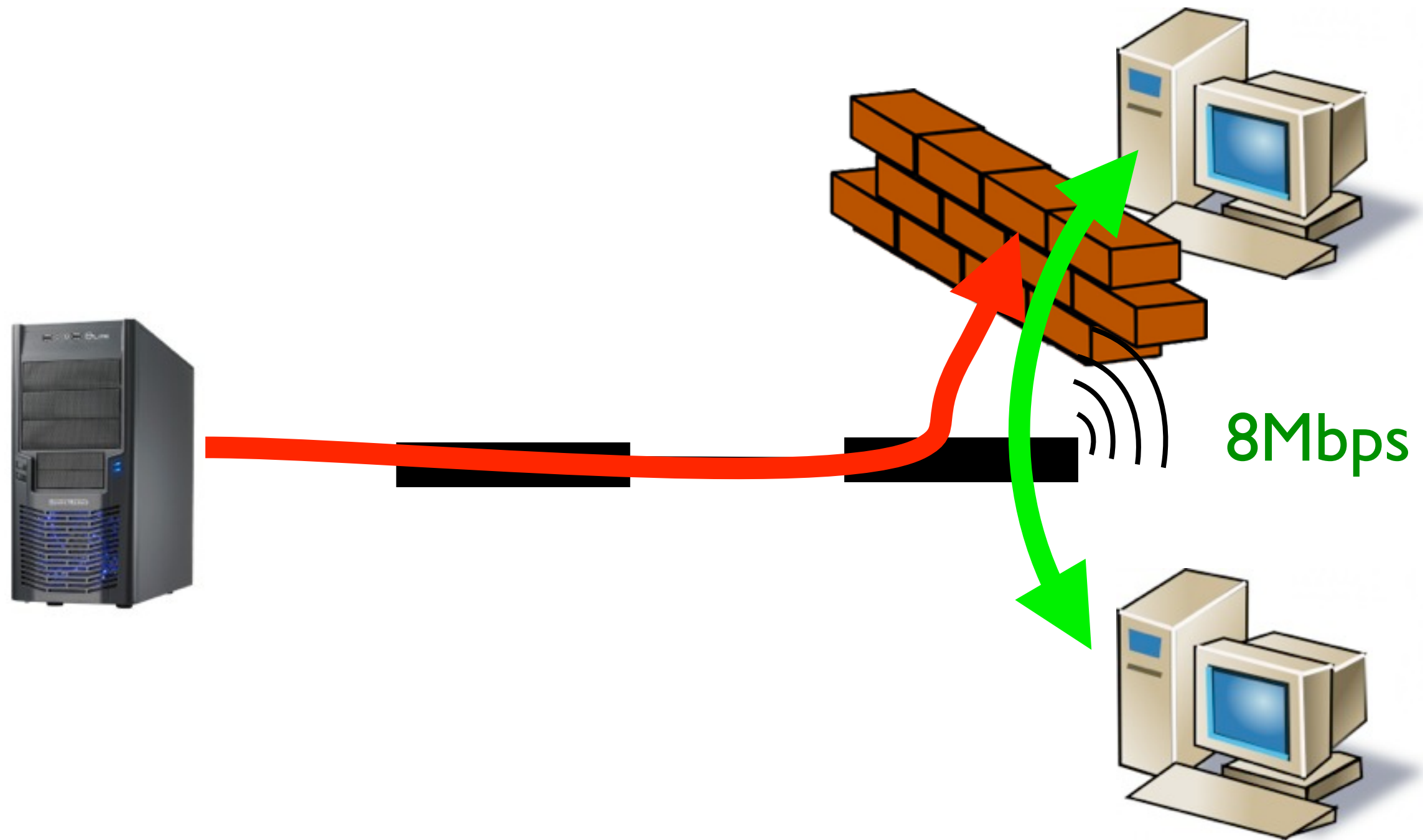
Denial-of-service



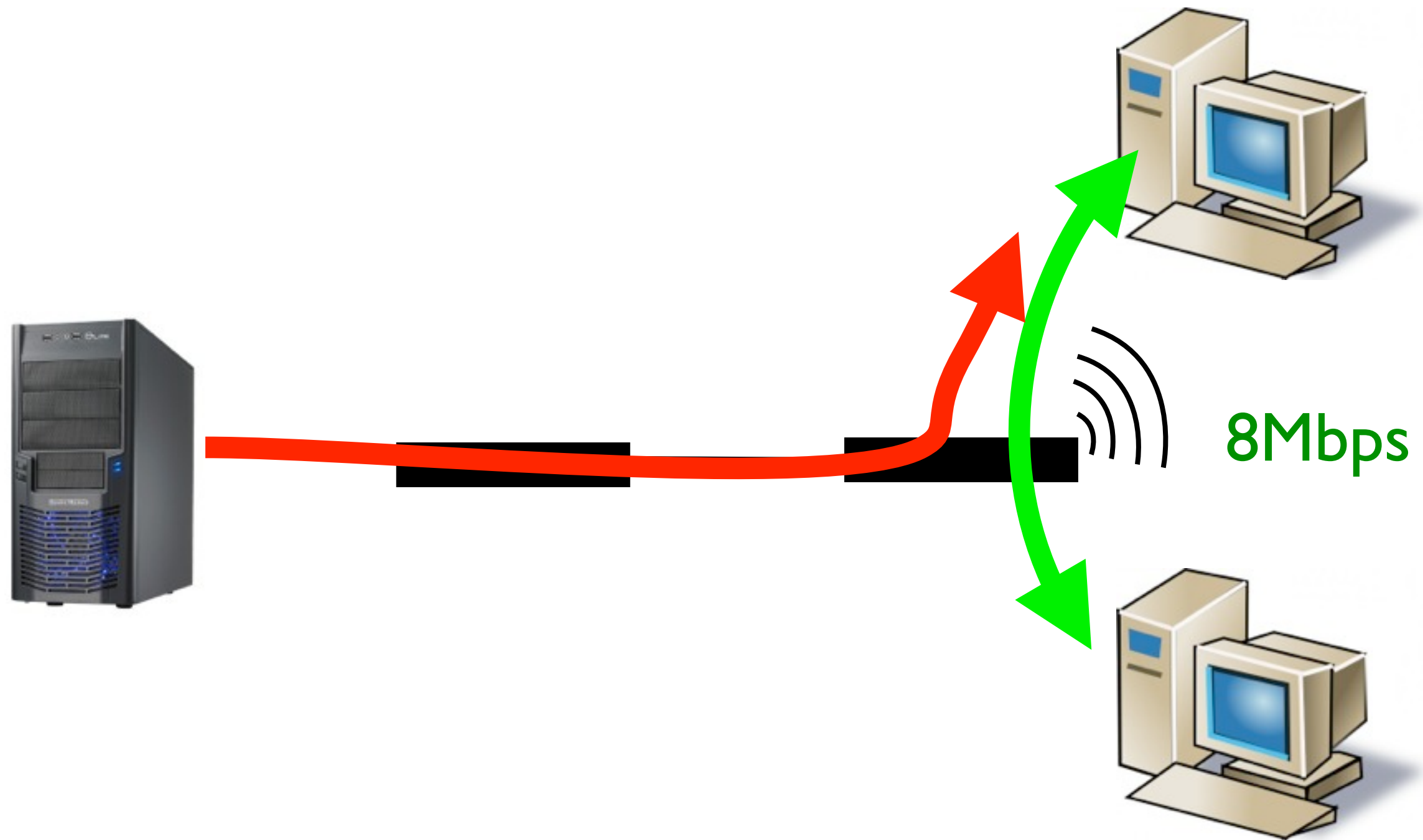
Denial-of-service



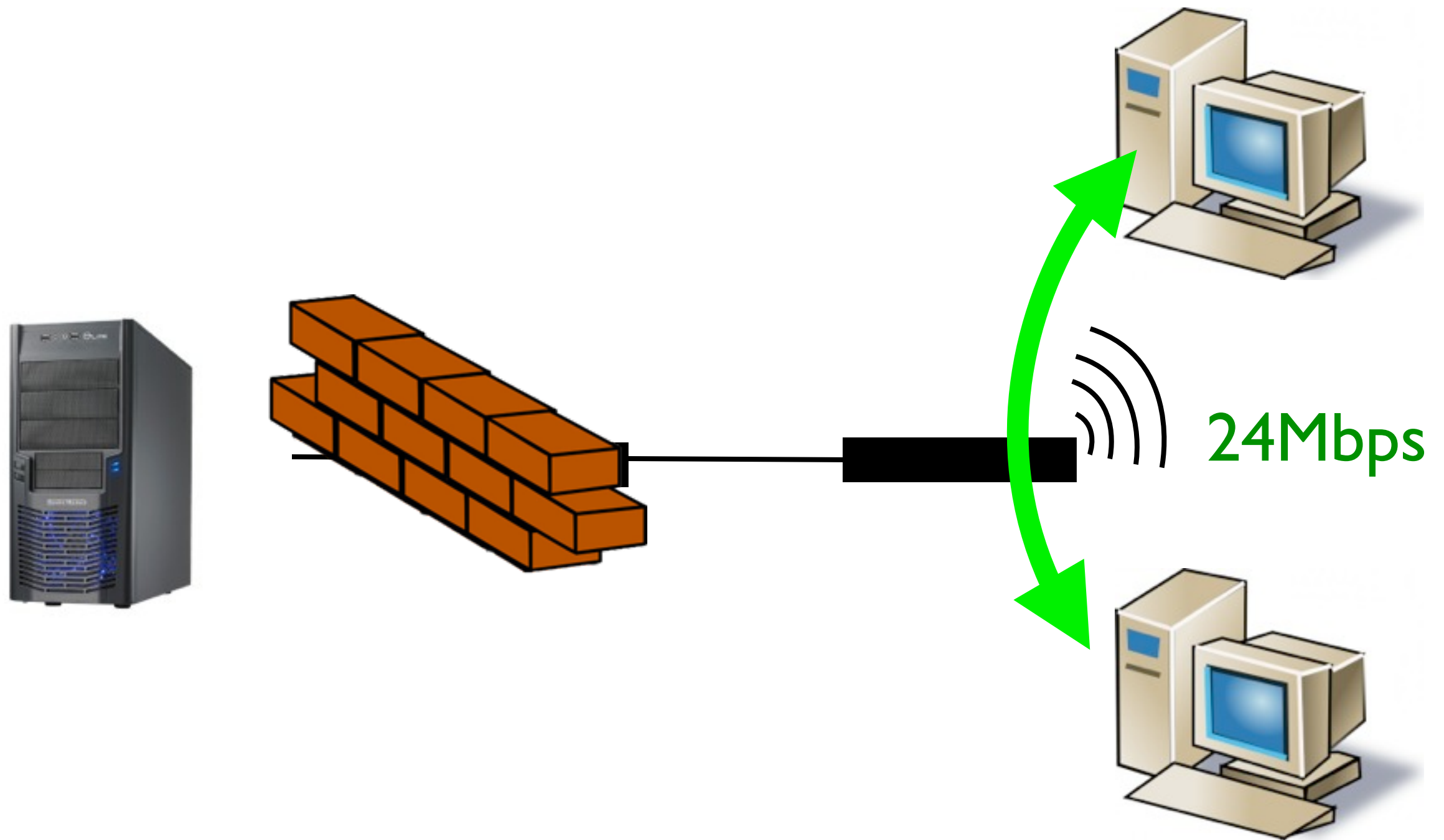
Denial-of-service



Denial-of-service



Denial-of-service



Denial-of-service

Where to go?

Implement more operators!

Where to go?

Implement more operators!

Guaranteed latency

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Path properties

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Path properties

Hints

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Path properties

Hints

Queries

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Path properties

Hints

Queries

Your application?

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Path properties

Hints

Queries

Your application?

Build a market

Where to go?

Implement more operators!

Guaranteed latency

Rate limiting

Path properties

Hints

Queries

Your application?

Build a market

Scale

Where to go?



Conclusion

Application knowledge can help cloud provider
PANE is our first step in realizing this

Conclusion

3640-123#show running-config Building configuration... Current configuration : 1432 bytes version 12.3 service config service timestamps debug datetime msec service timestamps log datetime msec service password- encryption hostname 3640-123 boot-start-marker boot-end-marker enable password 7 02050D4808095E731F no aaa new-model resource policy voice-card 3 ip subnet-zero ip cef no ip dhcp use vrf connected !--- This is the Cisco IOS Firewall configuration. !--- IN-OUT is the inspection rule for traffic that flows !--- from the inside interface of the router to the outside interface. ip inspect name IN- OUT tcp ip inspect name IN- OUT udp ip inspect name IN- OUT ftp ip inspect name IN- OUT http ip inspect name IN- OUT icmp !--- OUT-IN is the inspection rule for traffic that flows !--- from the outside interface of the router to the inside interface. !--- This rule is where SMTP/ESMTP inspection is specified. ip inspect name OUT-IN smtp no ip ips deny- action ips-	no ftp-server write-enable controller T1 3/0 framing sf linecode ami !--- The outside interface. interface Ethernet2/0 ip address 172.22.1.16 255.255.255.0 !--- Apply the access list to permit SMTP/ESMTP connections !--- to the mail server. This also allows Cisco IOS Firewall !--- to inspect SMTP or ESMTP commands. ip access-group 101 in ip nat outside !--- Apply the inspection rule OUT-IN inbound on this interface. This is !--- the rule that defines SMTP/ESMTP inspection. ip inspect OUT-IN in ip virtual- reassembly half-duplex interface Serial2/0 no ip address shutdown !--- The inside interface. interface Ethernet2/1 ip address 10.10.10.1 255.255.255.0 ip nat inside !--- Apply the inspection rule IN- OUT inbound on this interface. ip inspect IN-OUT in ip virtual- reassembly half-duplex ip http server no ip http secure- server ip classless	!--- The static translation for the mail server. ip nat inside source static 10.10.10.2 172.22.1.110 ip nat inside source static 10.10.10.5 172.22.1.111 !--- The access list to permit SMTP and ESMTP to the mail server. !--- Cisco IOS Firewall inspects permitted traffic. access-list 101 permit tcp any host 172.22.1.110 eq smtp control-plane voice-port 1/0/0 voice-port 1/0/1 voice-port 1/1/0 voice-port 1/1/1 line con 0 line aux 0 line vty 0 4 password 7 121A0C0411045D5679 login end Current configuration: version 12.2 service timestamps debug uptime service timestamps log uptime no service password-encryption hostname sec-3640 aaa new-model aaa group server tacacs+ RTP server 171.68.120.214 aaa authentication login default group RTP none aaa authorization exec default group RTP none aaa authorization auth-proxy default group RTP enable secret 5 \$1\$pqRI \$3TDNFT9FdYT8Sd/ q3S0VU1 enable password ww	ip inspect name myfw cuseeme timeout 3600 ip inspect name myfw ftp timeout 3600 ip inspect name myfw http timeout 3600 ip inspect name myfw rcmd timeout 3600 ip inspect name myfw realaudio timeout 3600 ip inspect name myfw smtp timeout 3600 ip inspect name myfw sqlnet timeout 3600 ip inspect name myfw streamworks timeout 3600 ip inspect name myfw tftp timeout 30 ip inspect name myfw udp timeout 15 ip inspect name myfw tcp timeout 3600 ip inspect name myfw vdolive ip auth-proxy auth- proxy-banner ip auth-proxy auth- cache-time 10 ip auth-proxy name list_a http ip audit notify log ip audit po max- events 100 interface Ethernet0/0 ip address 40.31.1.144 255.255.255.0 ip access-group 116 in ip nat outside ip auth-proxy list_a no ip route-cache no ip mroute-cache speed auto half-duplex no mop enabled interface Ethernet1/0 ip address 10.14.14.14 255.255.255.0	the syslog server or router. ip urlfilter audit- trail !--- use the ip urlfilter urlf- server-log command in global configuration mode to enable the logging of system messages on the URL filtering server. ip urlfilter urlf- server-log !--- use the ip urlfilter server vendor command in global configuration mode to configure a vendor server for URL filtering. Here we have configured a websense server for URL filtering ip urlfilter server vendor websense 192.168.15.15 no ftp-server write-enable !--- Below is the basic interface configuration on the router interface FastEthernet0 ip address 192.168.5.10 255.255.255.0 ip virtual- reassembly !--- use the ip inspect command in interface configuration mode to apply a set of inspection rules to an interface. Here the inspection name TEST is applied to the interface FastEthernet0. ip inspect test in duplex auto speed auto interface	ip virtual- reassembly duplex auto speed auto interface FastEthernet2 ip address 10.77.241.109 255.255.255.192 ip virtual- reassembly duplex auto speed auto interface FastEthernet2 no ip address interface Vlan1 ip address 10.77.241.111 255.255.255.192 ip virtual- reassembly ip classless ip route 10.10.10.0 255.255.255.0 172.17.1.2 ip route 10.77.0.0 255.255.0.0 10.77.241.65 !--- Configure the below commands to enable SDM access to the cisco routers ip http server ip http authentication local no ip http secure- server line con 0 line aux 0 line vty 0 4 privilege level 15 transport input telnet ssh end version 12.4 service timestamps debug datetime msec service timestamps log datetime msec no service password-encryption hostname 2851-cme2 logging message- counter syslog logging buffered 51200 warnings no aaa new-model clock timezone mst -7 clock summer-time	no ip dhcp use vrf connected ip dhcp pool pub-112-net network 172.17.112.0 255.255.255.0 default-router 172.17.112.1 dns-server 172.16.1.22 option 150 ip 172.16.1.43 domain-name bldrtme.com ip dhcp pool priv-112-net network 192.168.112.0 255.255.255.0 default-router 192.168.112.1 dns-server 172.16.1.22 domain-name bldrtme.com option 150 ip 192.168.112.1 ip domain name yourdomain.com no ipv6 cef multilink bundle- name authenticated voice translation- rule 1 rule 1 // /1001/ voice translation- profile default translate called 1 voice-card 0 no dspfarm interface GigabitEthernet0/0 description \$ETH- LAN\$ETH-SW-LAUNCH\$ \$INTF-INFO-GE 0/0\$ ip address 172.16.112.10 255.255.255.0 ip nat outside ip virtual- reassembly duplex auto speed auto interface GigabitEthernet0/1 no ip address duplex auto speed auto interface GigabitEthernet0/1. 132 encapsulation	interface GigabitEthernet0/1. 152 encapsulation dot1Q 152 ip address 192.168.112.1 255.255.255.0 ip nat inside ip virtual- reassembly interface FastEthernet0/2/0 interface FastEthernet0/2/1 interface FastEthernet0/2/2 interface FastEthernet0/2/3 interface Vlan1 ip address 198.41.9.15 255.255.255.0 router eigrp 1 network 172.16.112.0 0.0.0.255 network 172.17.112.0 0.0.0.255 no auto-summary ip forward-protocol nd ip http server ip http access- class 23 ip http authentication local ip http secure- server ip http timeout- policy idle 60 life 86400 requests 10000 ip http path flash:/gui ip nat inside source list 111 interface GigabitEthernet0/0 overload access-list 23 permit 10.10.10.0 0.0.0.7 access-list 111 deny ip 192.168.112.0 0.0.0.255 192.168.0.0 0.0.255.255 access-list 111 permit ip 192.168.112.0
---	---	--	---	--	---	---	--

```

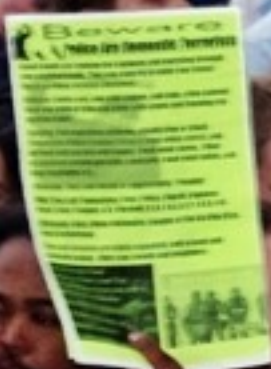
interface
GigabitEthernet0/1.
152
    encapsulation
dot1Q 152
    ip address
192.168.112.1
255.255.255.0
    ip nat inside
    ip virtual-
reassembly
interface
FastEthernet0/2/0
interface
FastEthernet0/2/1
interface
FastEthernet0/2/2
interface
FastEthernet0/2/3
interface Vlan1
    ip address
198.41.9.15
255.255.255.0
router eigrp 1
    network
172.16.112.0
0.0.0.255
    network
172.17.112.0
0.0.0.255
    no auto-summary
ip forward-protocol
nd
ip http server
ip http access-
class 23
ip http
authentication
local
ip http secure-
server
ip http timeout-
policy idle 60 life
86400 requests
10000
ip http path
flash:/gui
ip nat inside
source list 111
interface
GigabitEthernet0/0
overload
access-list 23
permit 10.10.10.0
0.0.0.7
access-list 111
deny ip
192.168.112.0
0.0.0.255
192.168.0.0
0.0.255.255
access-list 111
permit ip
192.168.112.0
0.0.0.255

```

OCCUPY EVERYTHING

#OCCUPYWALLST

WE ALREADY KNOW THAT WE OWN EVERYTHING—THE TASK IS TO EXCLUDE THE INTRUSIONS OF CAPITAL AND POWER



BANKS
BAILED
WE GO

Participatory Networking

Participatory Networking

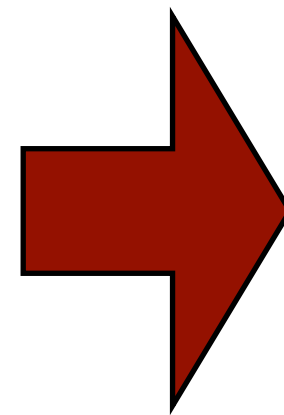
1. management API

Participatory Networking

1. management API
2. network controller

Participatory Networking

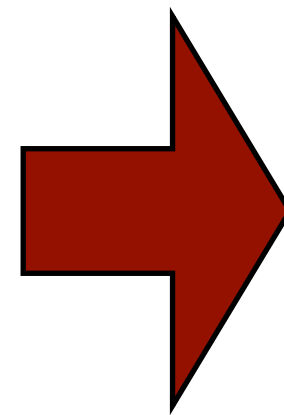
1. management API
2. network controller



Safe

Participatory Networking

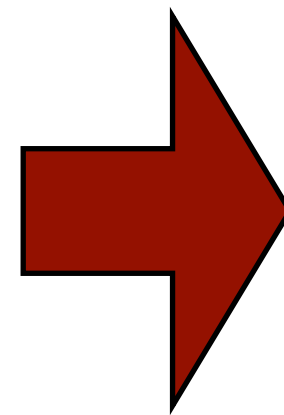
- 1. management API
- 2. network controller



Safe
Secure

Participatory Networking

1. management API
2. network controller



Safe
Secure
Fair

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Questions?

Shriram Krishnamurthi
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2012 USENIX Federated Conferences Week · June 12–15, 2012

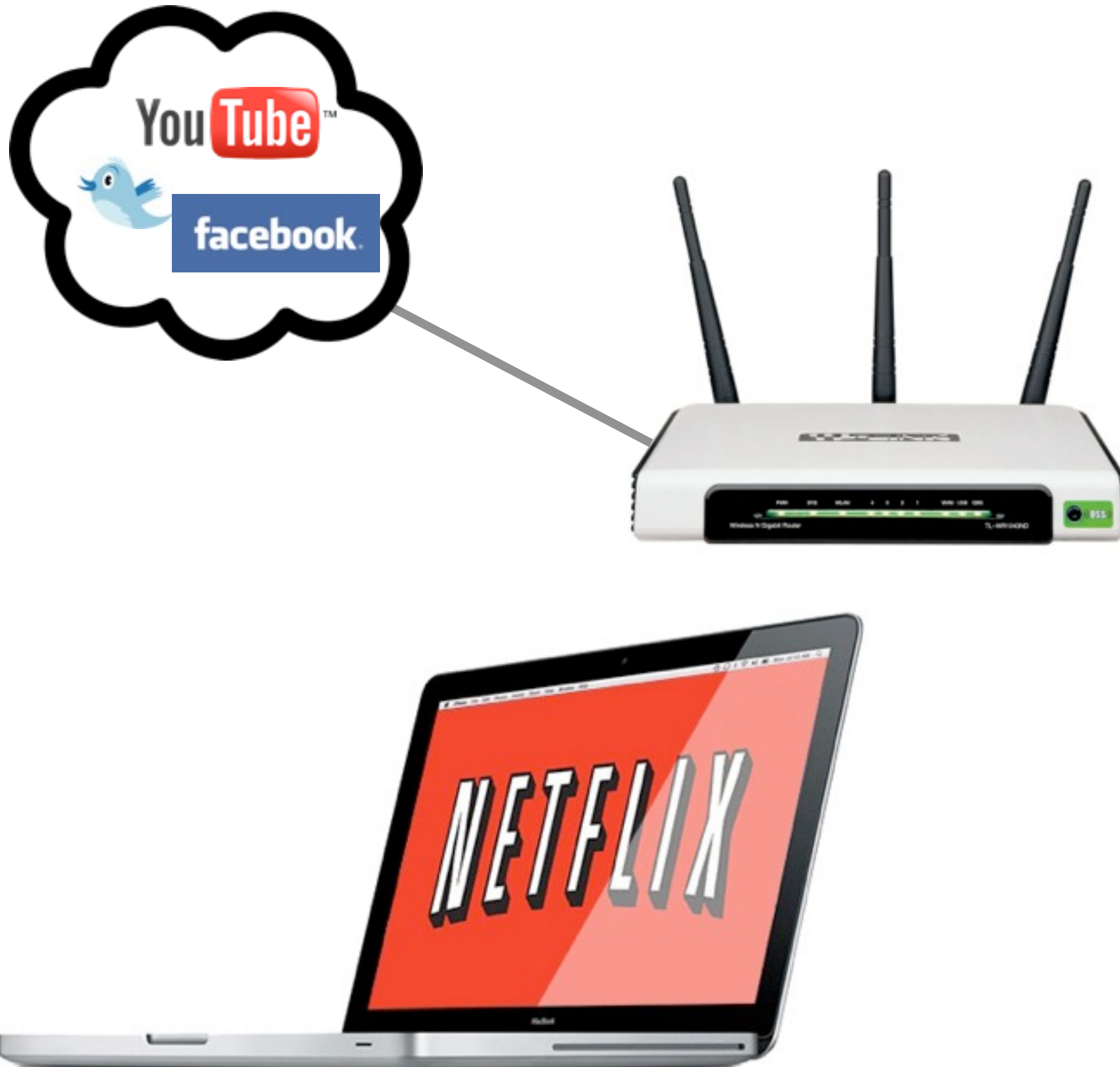
4th USENIX Workshop on Hot Topics in Cloud Computing

HotCloud '12

JUNE 12–13, 2012
BOSTON, MA



A problem in the home



NETFLIX

89%

Buffering

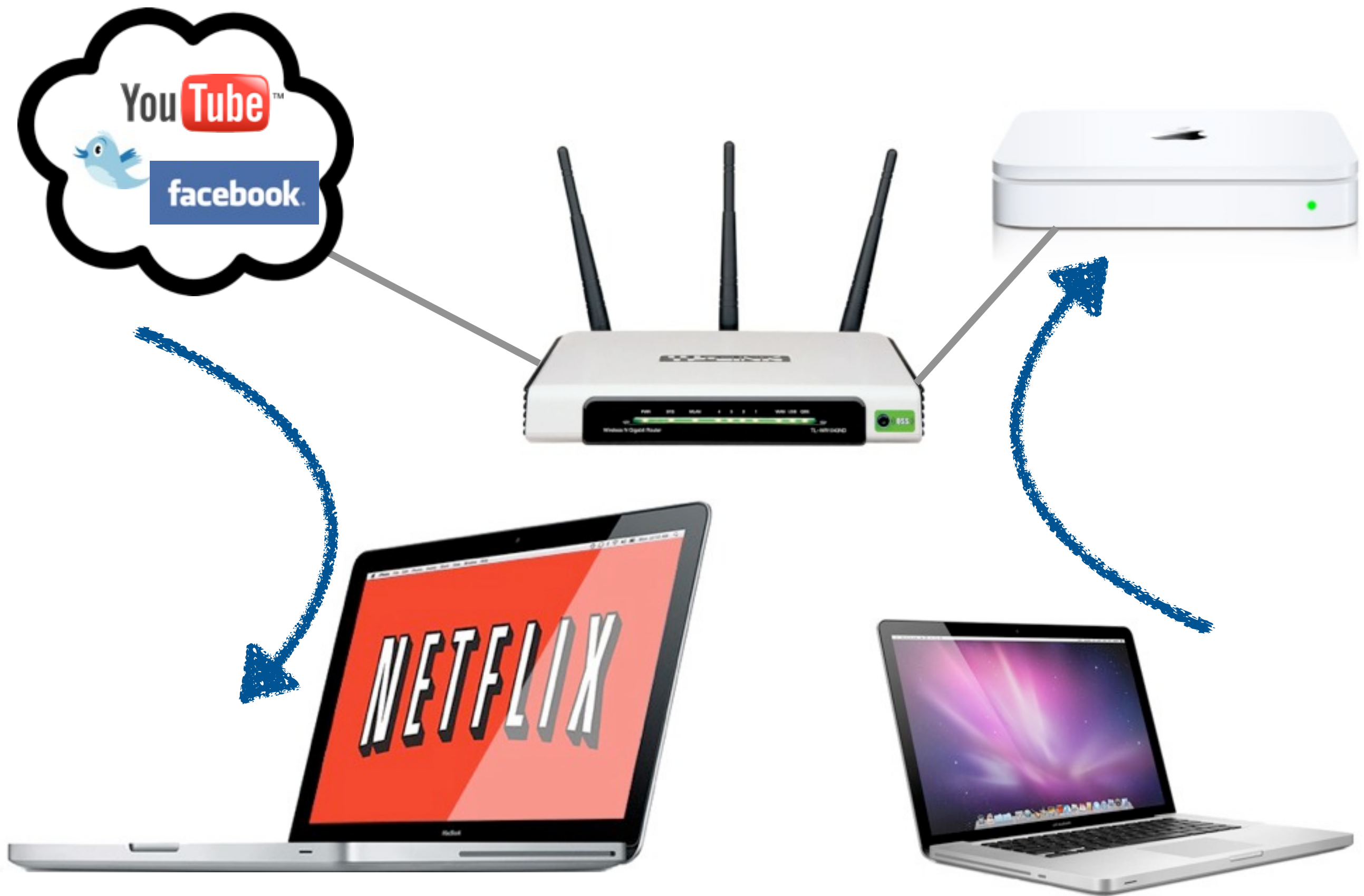


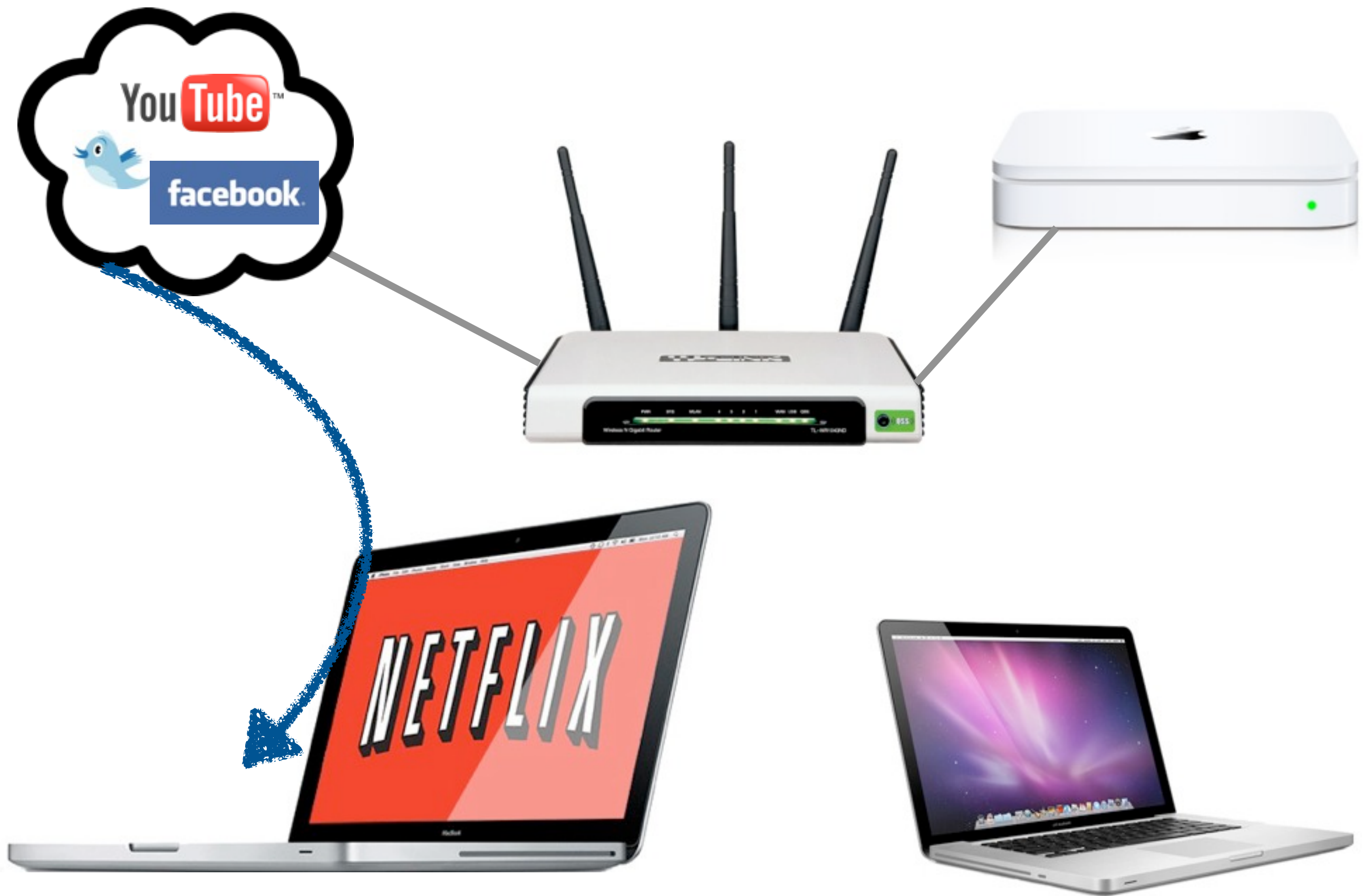
Full Screen

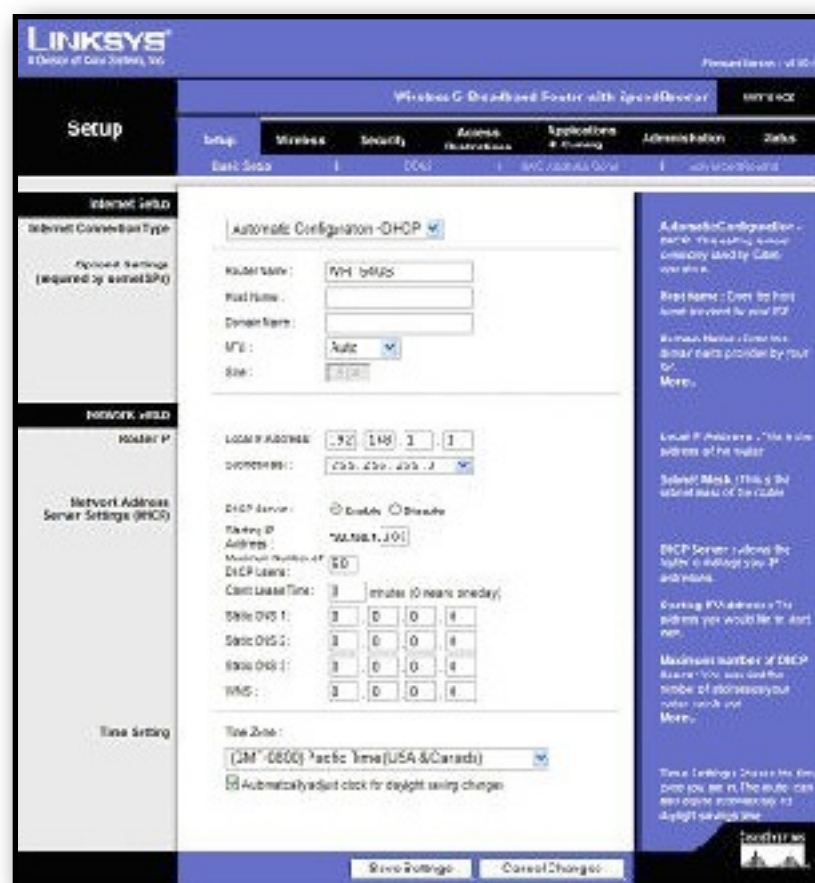


More Episodes

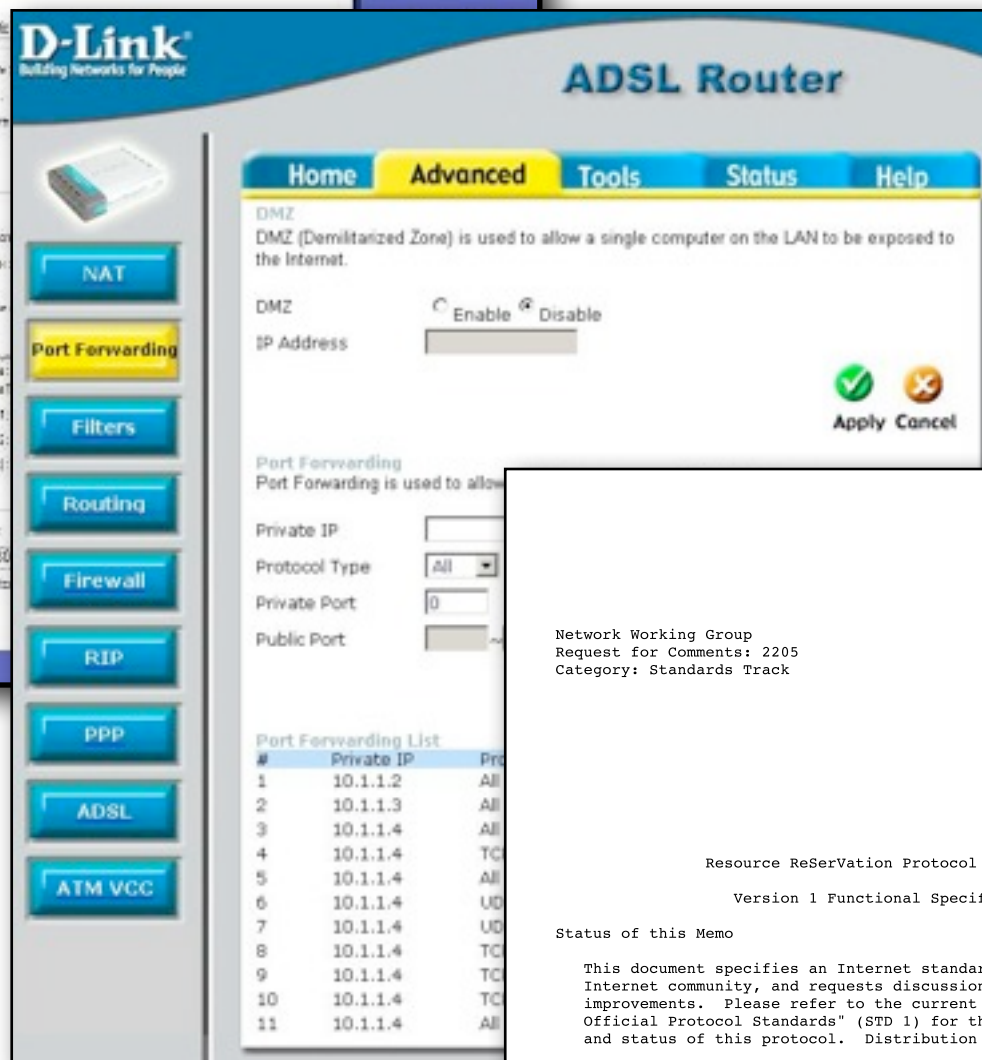
Back to Browsing











Network Working Group
Request for Comments: 2205
Category: Standards Track

R. Braden, Ed.
ISI
L. Zhang
UCLA
S. Berson
ISI
S. Herzog
IBM Research
S. Jamin
Univ. of Michigan
September 1997

Resource ReSerVation Protocol (RSVP) --

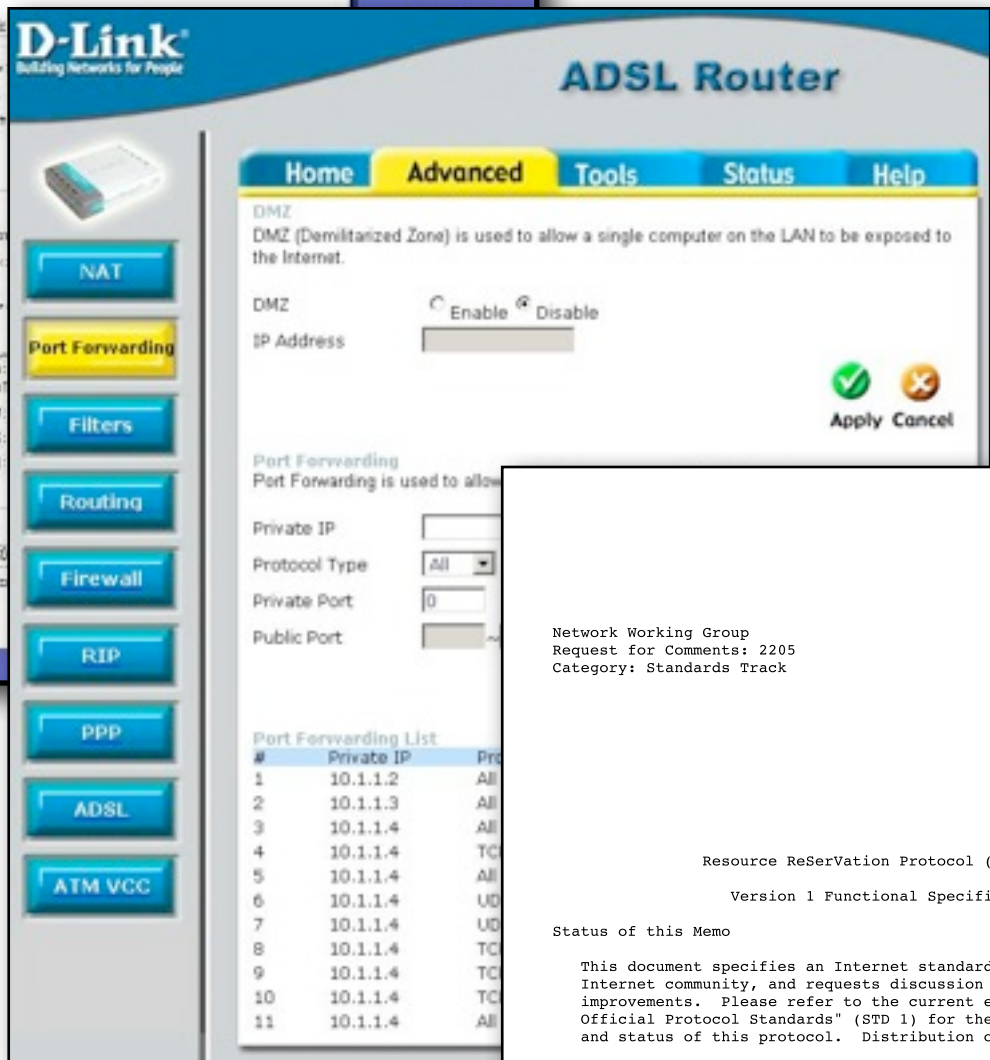
Version 1 Functional Specification

Status of this Memo

This document specifies an Internet standards track protocol for the Internet community, and requests discussion and suggestions for improvements. Please refer to the current edition of the "Internet Official Protocol Standards" (STD 1) for the standardization state and status of this protocol. Distribution of this memo is unlimited.

Abstract

This memo describes version 1 of RSVP, a resource reservation setup protocol designed for an integrated services Internet. RSVP provides receiver-initiated setup of resource reservations for multicast or unicast data flows, with good scaling and robustness properties.



Network Working Group
Request for Comments: 2205
Category: Standards Track

R. Braden, Ed., et. al.
L. Zhang
S. Braden
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September 1997

Resource ReSerVation Protocol (RSVP) -- Version 1 Functional Specification

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Abstract

This memo describes version 1 of RSVP, a resource reservation setpoint protocol designed for an integrated services Internet. RSVP provides receiver-initiated setup of resource reservations for multicast or unicast data flows, with good scaling and robustness properties.

TCP Nice: A Mechanism for Background Transfers

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Abstract

Many distributed applications can make use of large *background transfers* — transfers of data that humans are not waiting for — to improve availability, reliability, latency or consistency. However, given the rapid fluctuations of available network bandwidth and changing resource costs due to technology trends, hand tuning the aggressiveness of background transfers risks (1) complicating applications, (2) being too aggressive and interfering with other applications, and (3) being too timid and not gaining the benefits of background transfers. Our goal is for the operating system to manage network resources in order to provide a simple abstraction of near zero-cost background transfers. Our system, TCP Nice, can provably bound the interference inflicted by background flows on foreground flows in a restricted network model. And our microbenchmarks and case study applications suggest that in practice it interferes little with foreground flows, reaps a large fraction of spare network bandwidth, and simplifies application construction and deployment. For example, in our prefetching case study application, aggressive prefetching improves demand performance by a factor of three when Nice manages resources; but the same prefetching hurts demand performance by a factor of six under standard network congestion control.

1 Introduction

Many distributed applications can make use of large *background transfers* — transfers of data that humans are not waiting for — to improve service quality. For example, a broad range of applications and services such as data backup [29], prefetching [50], enterprise data distribution [20], Internet content distribution [2], and peer-to-peer storage [16, 43] can trade increased network

*This work was supported in part by an NSF CISE grant (CDA-9624082), the Texas Advanced Technology Program, the Texas Advanced Research Program, and Tivoli. Dahlin was also supported by an NSF CAREER award (CCR-9733842) and an Alfred P. Sloan Research Fellowship.

bandwidth consumption and possibly disk space for improved service latency [15, 18, 26, 32, 38, 50], improved availability [11, 53], increased scalability [2], stronger consistency [53], or support for mobility [28, 41, 47]. Many of these services have potentially unlimited bandwidth demands where incrementally more bandwidth consumption provides incrementally better service. For example, a web prefetching system can improve its hit rate by fetching objects from a virtually unlimited collection of objects that have non-zero probability of access [8, 10] or by updating cached copies more frequently as data change [13, 50, 48]; Technology trends suggest that “wasting” bandwidth and storage to improve latency and availability will become increasingly attractive in the future: per-byte network transport costs and disk storage costs are low and have been improving at 80-100% per year [9, 17, 37]; conversely network availability [11, 40, 54] and network latencies improve slowly, and long latencies and failures waste human time.

Current operating systems and networks do not provide good support for aggressive background transfers. In particular, because background transfers compete with foreground requests, they can hurt overall performance and availability by increasing network congestion. Applications must therefore carefully balance the benefits of background transfers against the risk of both *self-interference*, where applications hurt their own performance, and *cross-interference*, where applications hurt other applications' performance. Often, applications attempt to achieve this balance by setting “magic numbers” (e.g., the prefetch threshold in prefetching algorithms [18, 26]) that have little obvious relationship to system goals (e.g., availability or latency) or constraints (e.g., current spare network bandwidth).

Our goal is for the operating system to manage network resources in order to provide a simple abstraction of zero-cost background transfers. A self-tuning background transport layer will enable new classes of applications by (1) simplifying applications, (2) reducing the risk of being too aggressive, and (3) making